

CS313K: Logic, Sets, and Functions

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Lecture 9 – Chap 3 (3.4 – 3.8)

Homework Clarification

On Question 130 you're asked to "list the addresses," where the addresses are things like π_1 , π_2 , π_3 , etc. Just list the *indices* of the addresses, e.g., 1, 2, 3, etc.

On Implication

```
(defun implies (p q)
  (if p
      (if q t nil)
      t))
```

versus

```
(defun implies (p q)
  (if p
      (if q t nil)
      nil))
```

Let's prove " $x = 3$ implies that $x + 2 = 5$."

First, is it even true? What if x is 7? What if x is -23?

We'll formalize it as:

```
(implies (equal x 3)
          (equal (+ x 2) 5))
```

Is it true for all x ? Is it true when x is 7? -23?

(implies (equal x 3) (equal (+ x 2) 5))

```
(defun implies (p q)
  (if p
      (if q t nil)
      t))
```

versus

```
(defun implies (p q)
  (if p
      (if q t nil)
      nil))
```

Proof

```
(implies (equal x 3)
          (equal (+ x 2) 5))
```

```
{hyp}
```

```
(implies (equal x 3)
          (equal (+ 3 2) 5))
```

```
{computation}
```

```
(implies (equal x 3)
          t)
```

Why can we “prove” an implication by assuming the hypothesis is true?

Proof

```
(implies (equal x 3)
          (equal (+ x 2) 5))
{hyp}
(implies (equal x 3)
          (equal (+ 3 2) 5))
{computation}
(implies (equal x 3)
          t)
```

Why can we “prove” an implication by assuming the hypothesis is true? **Because the implication is T if the hypothesis is false!**

About “Addresses”

In the book I talk about the *address* of an occurrence of one term in another. For example, I might talk about π being the address of the 3rd A in

```
(if (zp a) 1 (* a (fact (- a 1))))
```

Some of you want me to tell you what an address *is*. I don't want to, because we don't really care. We just need to have a way to talk about *occurrences*.

```
(if (zp a) 1 (* a (fact (- a 1))))
0123456789012345678901234567890123456789
          10          20          30
```

But I could say “an address is the string index of the first character of the occurrence” If I said that, the 3rd a above would have address 27.

But I could, instead, say “an address is $(\alpha \ n)$, where α is a term and n is a number” meaning the n^{th} occurrence of the term α .

There are many other ways I could define *address*.

But how I define address doesn't matter, as long as we all understand that they identify a particular occurrence of a term.

About Propositional Occurrences

$(\text{IF } x \ y \ z)$ – x is used propositionally but whether y and z are being used propositionally depends on the context.

In $(\text{IF } (\text{IF } x \ y \ z) \ 1 \ 2)$, y and z are being used propositionally.

In $(+ (\text{IF } x \ y \ z) \ 0)$, y and z they are not.

The Propositional Functions

The arguments to AND, OR, NOT, IMPLIES, and IFF are all used propositionally if the term itself is used propositionally.

So all the variables in

(IMPLIES (AND p q) (OR s (NOT r)))

are being used propositionally.

I will never use these functions in non-propositional

ways. I.e., I won't write $(+ (\text{AND } p \ q) \ 3)$ even though it is legal.

About Maintaining Equivalence

IFF

(implies (and (natp e)

(mem e x))

(natp (if (and (f e) (g x))

(h (f x))

x)))

About Maintaining Equivalence

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About Maintaining Equivalence

IFF IFF IFF EQ

(implies (and (natp e)

IFF

(mem e x))

IFF

(natp (if (and (f e) (g x))

(h (f x))

x)))

About Maintaining Equivalence

IFF IFF IFF EQ

(implies (and (natp e)

IFF EQ EQ

(mem e x))

IFF

(natp (if (and (f e) (g x))

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About Maintaining Equivalence

IFF IFF IFF EQ

(implies (and (natp e)

IFF EQ EQ

(mem e x))

IFF EQ

(natp (if (and (f e) (g x))

(h (f x))

x)))

About Maintaining Equivalence

```
IFF      IFF  IFF  EQ
(implies (and (natp e)
              IFF EQ EQ
              (mem e x))
          IFF  EQ  IFF
(natp (if (and (f e) (g x))
          EQ
          (h (f x))
          EQ
          x)))
```

About Maintaining Equivalence

IFF IFF IFF EQ

(implies (and (natp e)

IFF EQ EQ

(mem e x))

IFF EQ IFF IFF IFF

(natp (if (and (f e) (g x))

EQ

(h (f x))

EQ

x)))

About Maintaining Equivalence

```
IFF      IFF  IFF  EQ
(implies (and (natp e)
              IFF EQ EQ
              (mem e x))
          IFF  EQ  IFF  IFFEQ  IFFEQ
(natp (if (and (f e) (g x))
          EQ EQ EQ
          (h (f x))
          EQ
          x)))
```

About Maintaining Equivalence

IFF IFF IFF EQ

(implies (and (natp e)

IFF EQ EQ

(mem e x))

IFF

(and (if (and (f e) (g x))

(h (f x))

x)

...))

About Maintaining Equivalence

```
IFF      IFF  IFF  EQ
(implies (and (natp e)
              IFF EQ EQ
              (mem e x))
          IFF IFF
          (and (if (and (f e) (g x))
                   (h (f x))
                 x)
              ...))
```

About Maintaining Equivalence

```
IFF      IFF  IFF  EQ
(implies (and (natp e)
              IFF EQ EQ
              (mem e x))
          IFF  IFF  IFF
          (and (if (and (f e) (g x))
                  IFF
                  (h (f x))
                  IFF
                  x)
              ...))
```