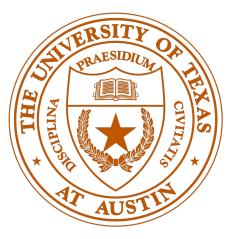
From Crash Consistency to Transactions



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Persistent data is structured; crash consistency hard

- Structured data abstractions built on file system
 - SQLite, BerkeleyDB... -- Embedded DB
 - LevelDB, Redis, MongoDB... -- Key-value store
 - Images, binary blobs...
 Files
- Applications manage storage themselves
 - ...and poorly!
 - The POSIX interface is no longer sufficient

Easy to use & deploy

Data safe on crash

ACID across abstractions

High performance

A transactional file system is the answer

- Structured data uses file system storage
 - Easy management often outweighs high performance
- File system transactions provides API and mechanisms

Easy to use & deploy

Data safe on crash

• Transactions preserve consistency

High performance

- Transactions reduce work & syncs
- Concurrent transactions scalable

ACID across abstractions

Unify different types of updates

We need transactions across storage abstractions

- The Android mail client receives an email with attachment
 - Stores attachment as a regular file
 - File name of attachment stored in SQLite
 - Stores email text in SQLite
- Great work when you can get it, but what can go wrong?
 - Crashes can orphan attachment files
 - Crashes can leave incomplete attachments
 - And this level of crash consistency costs dearly in performance!

How many syncs do you need?

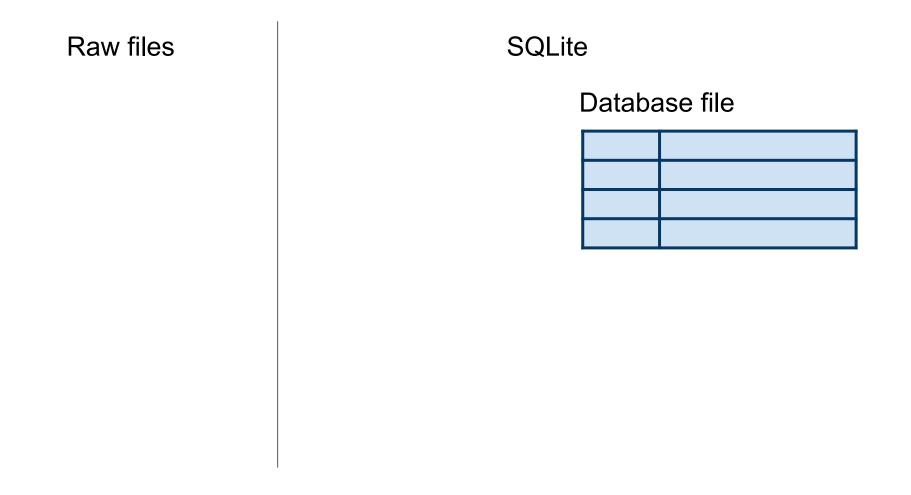
- The Android mail client receives an email with attachment
 - Stores attachment as a regular file (maybe 1 sync?)
 - File name of attachment stored in SQLite
 - Stores email text in SQLite (maybe 1 sync for db? 2 total?)

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How many syncs do you need?

- The Android mail client receives an email with attachment
 - Stores attachment as a regular file (maybe 1 sync?)
 - File name of attachment stored in SQLite
 - Stores email text in SQLite (maybe 1 sync for db? 2?)
- Requires 6 syncs!
 - If you create/delete a file, sync the parent directory

Atomically inserting a message with attachment.



Atomically inserting a message with attachment.

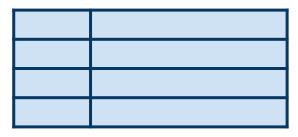
Raw files

Attachment file



1.create(/dir/attachment) write(/dir/attachment) fsync(/dir/attachment) fsync(/dir/) SQLite

Database file



Atomically inserting a message with attachment.

Raw files Attachment file 1.create(/dir/attachment)

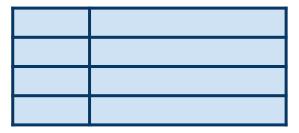
write(/dir/attachment) **fsync**(/dir/attachment) fsync(/dir/)

Roll-back log Rollback info 2.create(/dir/journal)

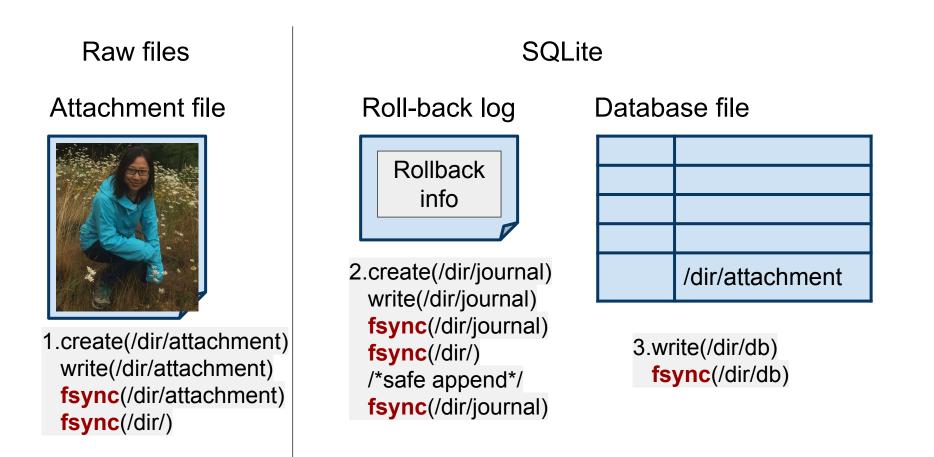
write(/dir/journal) **fsync**(/dir/journal) fsync(/dir/) /*safe append*/ **fsync**(/dir/journal)

SQLite

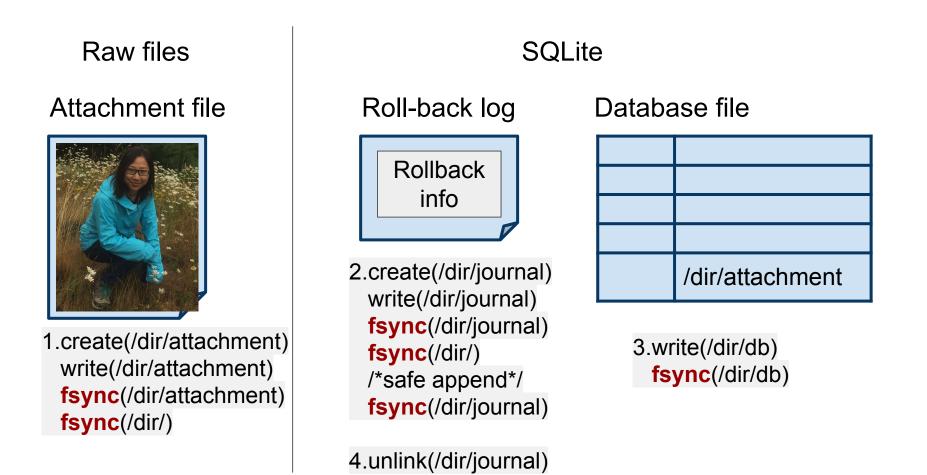
Database file



Atomically inserting a message with attachment.



Atomically inserting a message with attachment.



Application consistency using POSIX is slow

• SQLite on ext4: fsync() per transaction (1kB/tx), with FULL synchronization level.

	fsync/tx			
Journal mode	Insert	Update		
Rollback (default)	4	4		
Write ahead log (WAL)	5	5		
No journal (unsafe)	1	1		

SQLite

Database file

3.write(/dir/db) fsync(/dir/db)

4.unlink(/dir/journal)

/*safe append*/ fsync(/dir/journal)

Roll-back log

Rollback info

2.create(/dir/journal)

write(/dir/journal) fsync(/dir/journal)

fsync(/dir/)

Raw files

Attachment file

1.create(/dir/attachment)

write(/dir/attachment)

fsync(/dir/attachment)

fsync(/dir/)

Roll-back log

Rollback info

2.create(/dir/jou write(/dir/jour

fsync(/dir/)

4.unlink(/dir/journal)

Database file

urnal)	/dir/attachment
nal)	

fsync(/dir/journal) /*safe append*/ fsvnc(/dir/iournal)

SQLite

3.write(/dir/db) fsync(/dir/db)

System support for crash consistent updates

- Application needs consistent, persistent updates
 - Complicated and ad hoc implementation Ο
 - Crashes can orphan attachment files Ο
 - Crashes can create incomplete attachment files. \bigcirc
- Sync and redundant writes lead to poor performance.
- Need mechanism for cross-abstraction commit

The file system should provide transactional services!

But haven't we tried this before?

Haven't we seen this movie before?

- Complex implementation
 - Transactional OS: QuickSilver [TOCS 88], TxOS [SOSP 09] (**10k LOC**)
 - In-kernel transactional file systems: Valor [FAST 09]
- Hardware dependent
 - CFS [ATC 15], MARS [SOSP 13], TxFLash [OSDI 08], Isotope [FAST 16]
- Performance overhead
 - Valor [FAST 09] (**35% overhead**).
- Hard to use
 - Windows NTFS (TxF), released 2006 (deprecated 2012)

Modify the following code to use Windows NTFS (TxF) transactions.

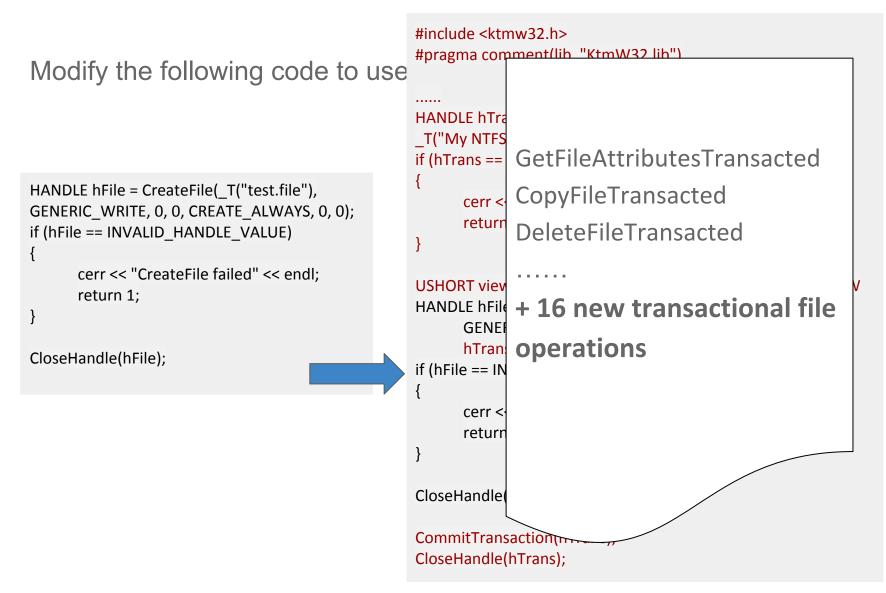
```
HANDLE hFile = CreateFile(_T("test.file"),
GENERIC_WRITE, 0, 0, CREATE_ALWAYS, 0, 0);
if (hFile == INVALID_HANDLE_VALUE)
{
cerr << "CreateFile failed" << endl;
return 1;
}
CloseHandle(hFile);
```

```
#pragma comment(lib, "KtmW32.lib")
Modify the following code to use
                                                   HANDLE hTrans = CreateTransaction(NULL,0, 0, 0, 0, NULL,
                                                   T("My NTFS Transaction"));
                                                   if (hTrans == INVALID HANDLE VALUE)
HANDLE hFile = CreateFile( T("test.file"),
                                                         cerr << "CreateTransaction failed" << endl:
GENERIC WRITE, 0, 0, CREATE ALWAYS, 0, 0);
                                                         return 1;
if (hFile == INVALID HANDLE VALUE)
      cerr << "CreateFile failed" << endl;
                                                   USHORT view = 0xFFFE; // TXFS MINIVERSION DEFAULT VIEW
      return 1;
                                                   HANDLE hFile = CreateFileTransacted( T("test.file"),
}
                                                         GENERIC WRITE, 0, 0, CREATE ALWAYS, 0, 0,
                                                         hTrans, &view, NULL);
CloseHandle(hFile);
                                                   if (hFile == INVALID HANDLE VALUE)
                                                         cerr << "CreateFileTransacted failed" << endl;
                                                         return 1;
```

CloseHandle(hFile);

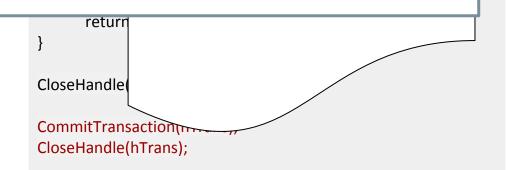
#include <ktmw32.h>

CommitTransaction(hTrans); CloseHandle(hTrans);



Microsoft deprecates TxF (2012) Modify the HANDLE hFile = Cr GENERIC WRITE, if (hFile == INVALI cerr << "Cre return 1; } CloseHandle(hFile

"While TxF is a powerful set of APIs, there has been extremely limited developer interest in this API platform since Windows Vista primarily due to its complexity and various nuances which developers need to consider as part of application development."



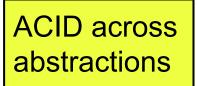
T2FS (Texas Transactional File System)

- Based on Linux ext4
 - Uses file system journal
- Simple interface
 - fs_tx_begin, fs_tx_end, fs_tx_abort
- Usable by any abstraction that stores data in the file system
 - E.g., embedded databases, key-value stores
- Improves performance for structured data
 - Fewer sync calls
- Increases scalability

19

High performance





Data safe on crash

T2FS API

Modify the following code to use T2FS transactions.

```
HANDLE hFile = CreateFile(_T("test.file"),
GENERIC_WRITE, 0, 0, CREATE_ALWAYS, 0, 0);
if (hFile == INVALID_HANDLE_VALUE)
{
cerr << "CreateFile failed" << endl;
return 1;
}
CloseHandle(hFile);
```

T2FS API

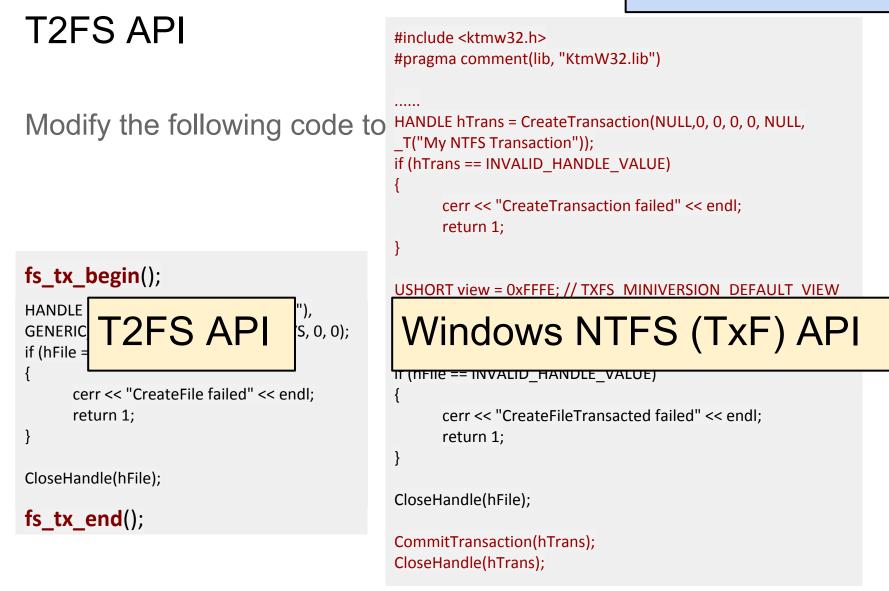
Modify the following code to use T2FS transactions.

fs_tx_begin();

```
HANDLE hFile = CreateFile(_T("test.file"),
GENERIC_WRITE, 0, 0, CREATE_ALWAYS, 0, 0);
if (hFile == INVALID_HANDLE_VALUE)
{
    cerr << "CreateFile failed" << endl;
    return 1;
}
CloseHandle(hFile);
```

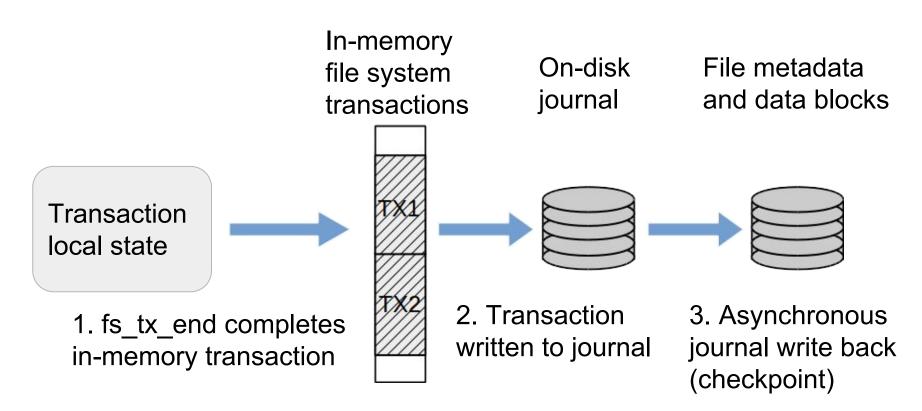
```
fs_tx_end();
```

Easy to use & deploy



T2FS managing and persisting transactions

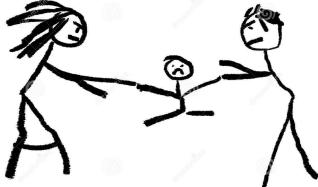
- Decreased complexity: use the file systems' crash consistency mechanism to create transactions.
 - Ext4 journal or ZFS copy-on-write



Data safe on crash

Isolation and Conflict detection

- In-progress writes are all local to kernel thread
- Eager conflict detection on inodes, directory entries
 - Enables flexible contention management
- Fine-grained page locks
 - More scalable than reader/writer lock



Modify the Android mail application to use T2FS transactions.

Raw files SQLite Roll-back log Database file Attachment file Rollback info 2.create(/dir/journal) /dir/attachment write(/dir/journal) **fsync**(/dir/journal) 1.create(/dir/attachment) 3.write(/dir/db) fsync(/dir/) write(/dir/attachment) fsync(/dir/db) /*safe append*/ **fsync**(/dir/attachment) **fsync**(/dir/journal) fsync(/dir/) 4.unlink(/dir/journal)

Modify the Android mail application to use T2FS transactions.

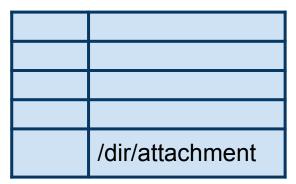
Raw files

Attachment file



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Modify the Android mail application to use T2FS transactions.

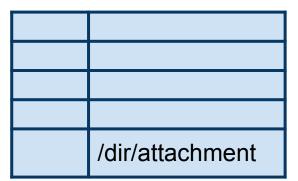
Raw files

Attachment file



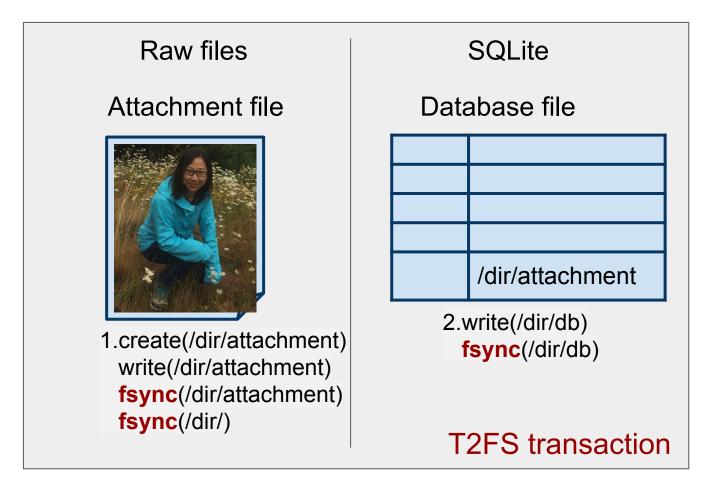
1.create(/dir/attachment) write(/dir/attachment) fsync(/dir/attachment) fsync(/dir/) SQLite

Database file

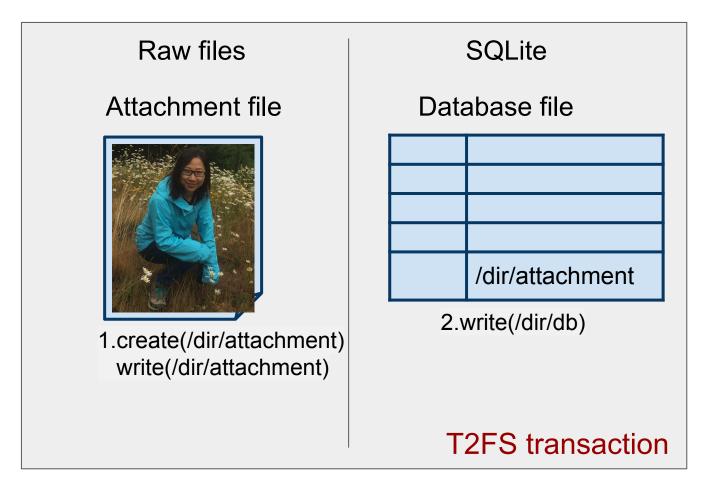


2.write(/dir/db) fsync(/dir/db)

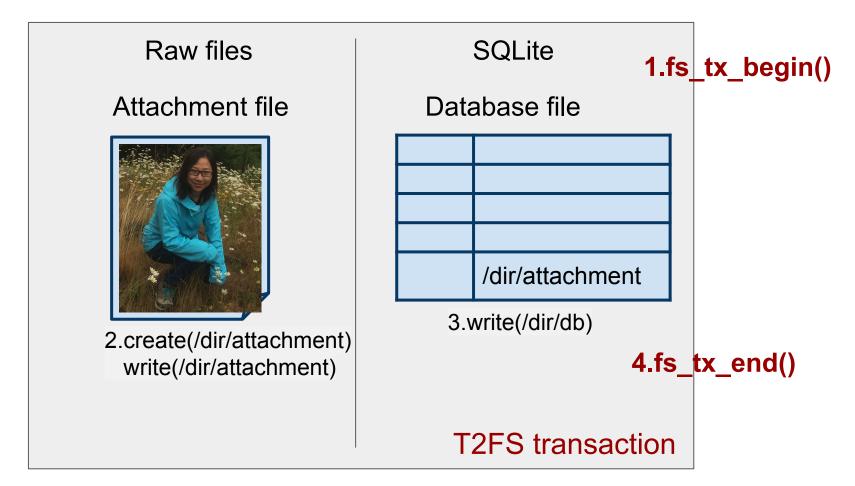
Modify the Android mail application to use T2FS transactions.



Modify the Android mail application to use T2FS transactions.



Modify the Android mail application to use T2FS transactions.



Evaluation: single-threaded SQLite



1.5M 1KB operations. 10K operations grouped in a transaction. Database prepopulated with 15M rows.

Transactions as a foundation for other optimizations

- Enable automatic file system optimizations
 - Eliminate temporary durable files.
 - e.g. SQLite delete mode, directly wrapped by T2FS transaction
 - Consolidate IO across transactions.
 - Delay persistence during commit
- Use transactional mechanism to implement unrelated file system optimizations
 - Separate ordering from durability (osync [SOSP 13]).

Summary

Easy to use & deploy

Data safe on crash

ACID across abstractions

High performance

- Persistent data is structured; tough to make crash consistent
 - All data stored in the file system
- A transactional file system has the right API and mechanisms
- The file system journal makes implementing transactions easier
- Need transactions across storage abstractions

Thank you!