

"Less than" in terms of "at most"

If, a in the case of comparing reals, \leq is a total order, $<$ is most simply defined by

$$(0) \quad x < y \equiv \neg(y \leq x)$$

The "transitivity" rule

$$(1) \quad a \leq b \wedge b < c \Rightarrow a < c$$

is then a direct consequence of the transitivity of \leq :

$$\begin{aligned} & a \leq b \wedge b < c \Rightarrow a < c \\ = & \{ (0), \text{twice} \} \\ & a \leq b \wedge \neg(c \leq b) \Rightarrow \neg(c \leq a) \\ = & \{ \text{shunting, twice} \} \\ & c \leq a \wedge a \leq b \Rightarrow c \leq b \\ = & \{ \leq \text{transitive} \} \\ & \text{true} \end{aligned}$$

If \leq is not a total order, this does not work, and I used to use the definition

$$(2) \quad x < y \equiv x \leq y \wedge x \neq y$$

The proof of (1) is then as follows:

$$\begin{aligned} & a \leq b \wedge b < c \Rightarrow a < c \\ = & \{ (2), \text{twice} \} \\ & a \leq b \wedge b \leq c \wedge b \neq c \Rightarrow a \leq c \wedge a \neq c \\ = & \{ \leq \text{is transitive} \} \end{aligned}$$

$$\begin{aligned}
& a \leq b \wedge b \leq c \wedge b \neq c \Rightarrow a \neq c \\
= & \quad \{ \text{shunting, twice} \} \\
& a \leq b \wedge b \leq c \wedge a = c \Rightarrow b = c \\
= & \quad \{ \text{Leibniz} \} \\
& c \leq b \wedge b \leq c \wedge a = c \Rightarrow b = c \\
= & \quad \{ \leq \text{ reflexive} \} \\
& \text{true}
\end{aligned}$$

In [GS94] - "A logical approach to discrete math" by Gries & Schneider - I found the much nicer definition:

$$(3) \quad x < y \equiv x \leq y \wedge \neg(y \leq x)$$

The proof of (1) is now very clean:

$$\begin{aligned}
& a \leq b \wedge b < c \Rightarrow a < c \\
= & \quad \{ (3), \text{ twice} \} \\
& a \leq b \wedge b \leq c \wedge \neg(c \leq b) \Rightarrow a \leq c \wedge \neg(c \leq a) \\
\Leftarrow & \quad \{ \text{pred. calc.} \} \\
& (a \leq b \wedge b \leq c \Rightarrow a \leq c) \wedge \\
& (c \leq a \wedge a \leq b \Rightarrow c \leq b) \\
= & \quad \{ \leq \text{ transitive, twice} \} \\
& \text{true.}
\end{aligned}$$

The moral of the story is that the conjunct $x \neq y$ in (2) is awkward: it forces in general, in one way or another, an appeal to Leibniz, and in this special case also an appeal to the reflexivity of \leq . Definition (3) is much nicer because it

uses \Leftarrow only.

In retrospect I am amazed that I did not come up with (3) myself, for it is a straightforward strengthening of (0), with which I grew up.

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