

CS429: Computer Organization and Architecture

Instruction Set Architecture V

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Basic Data Types

Integral

- Stored and operated on in general registers.
- Signed vs. unsigned depends on instructions used.

Intel	GAS	Bytes	C
byte	b	1	[unsigned] char
word	w	2	[unsigned] short
double word	l	4	[unsigned] int
quad word	q	8	[unsigned] long int

Floating Point

Stored and operated on in floating point registers.

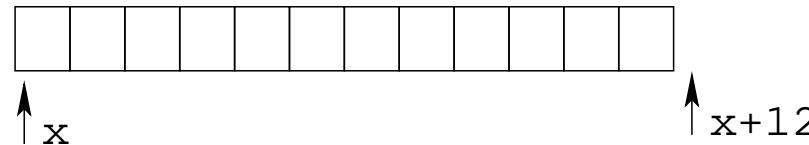
Intel	GAS	Bytes	C
Single	s	4	float
Double	l	8	double
Extended	t	10/12	long double

Array Allocation

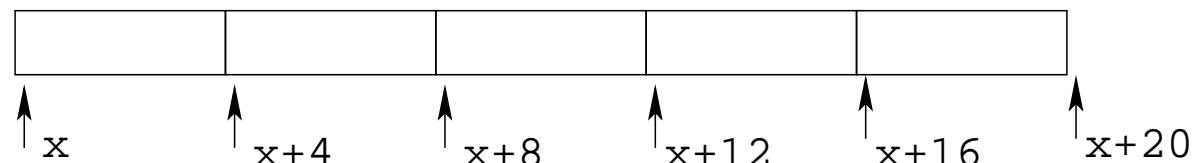
Basic Principle: $T \ A[L]$

- Array (named A) of data type T and length L.
- Contiguously allocated region of $L * \text{sizeof}(T)$ bytes.

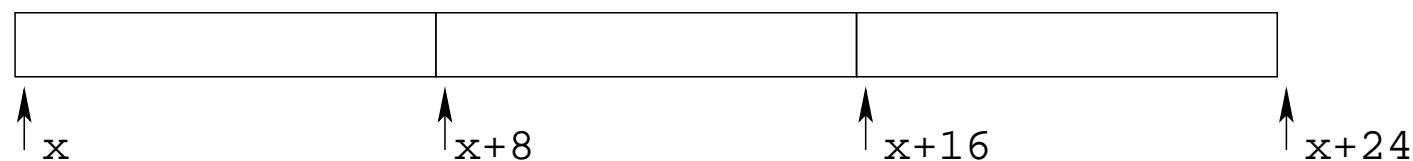
```
char string[12];
```



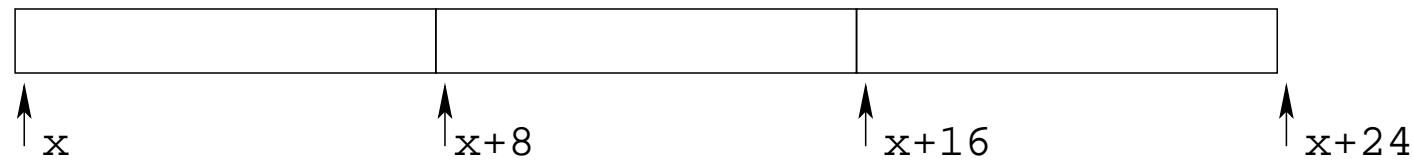
```
int val[5];
```



```
double a[3];
```

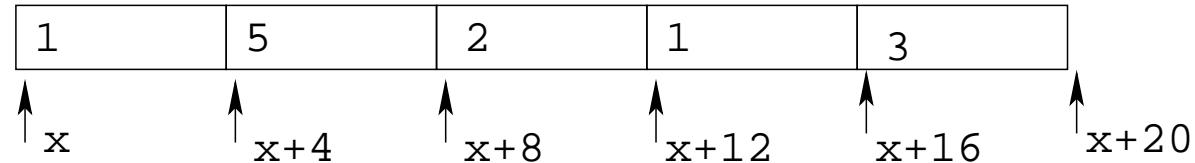


```
char *p[3];
```



Array Access

```
int val[5];
```

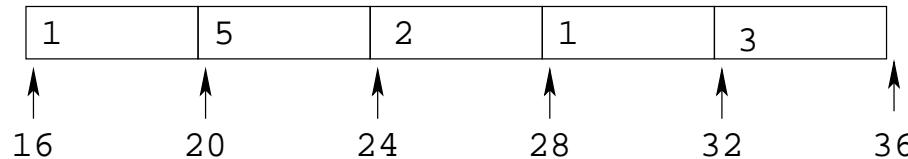


Reference	Type	Value
<code>val[4]</code>	<code>int</code>	3
<code>val</code>	<code>int *</code>	x
<code>val+1</code>	<code>int *</code>	$x + 4$
<code>&val[2]</code>	<code>int *</code>	$x + 8$
<code>val[5]</code>	<code>int</code>	??
<code>*(val+1)</code>	<code>int</code>	5
<code>val+j</code>	<code>int *</code>	$x + 4j$

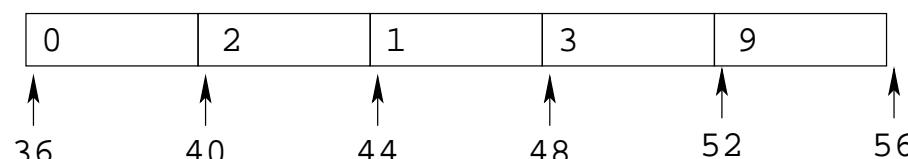
Note the use of pointer arithmetic.

Array Example

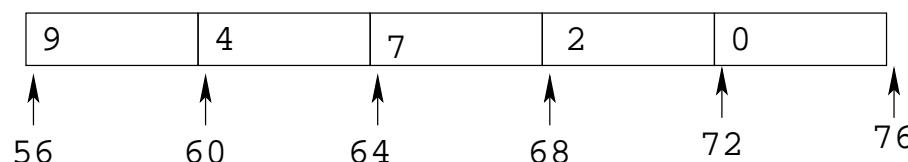
zip_dig cmu;



zip_dig mit;



zip_dig ucb;



```
#define ZLEN 5
typedef int zip_dig [ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

Declaration `zip_dig cmu` is equivalent to `int cmu[5]`.

Example arrays were allocated in successive 20 byte block.

That's not guaranteed to happen in general.

Array Accessing Example

```
int get_digit
    ( zip_dig z, int dig )
{
    return z[dig];
}
```

Memory Reference Code

```
# %rdi = z
# %rsi = dig
movl (%rdi,%rsi,4),%eax # z[dig]
```

Computation

- Register %rdi contains the starting address of the array.
- Register %rsi contains the array index.
- The desired digit is at $\%rdi + (4 * \%rsi)$.
- User memory reference $(\%rdi, \%rsi, 4)$.

Array Loop Example

```
void zincr( zip_dig z ) {
    size_t i;
    for (i = 0; i < ZLEN; i++)
        z[i]++;
}
```

```
# %rdi = z
movl $0, %eax          #     i = 0
jmp .L3                #     goto middle
.L4:                  # loop:
    addl $1, (%rdi, %rax, 4) #     z[i]++
    addq $1, %rax           #     i++
.L3:                  # middle:
    cmpq $4, %rax          #     i:4
    jbe .L4                #     if <=, goto loop
    ret                    #     return
```

Multidimensional (Nested) Arrays

Declaration: $T \ A[R][C];$

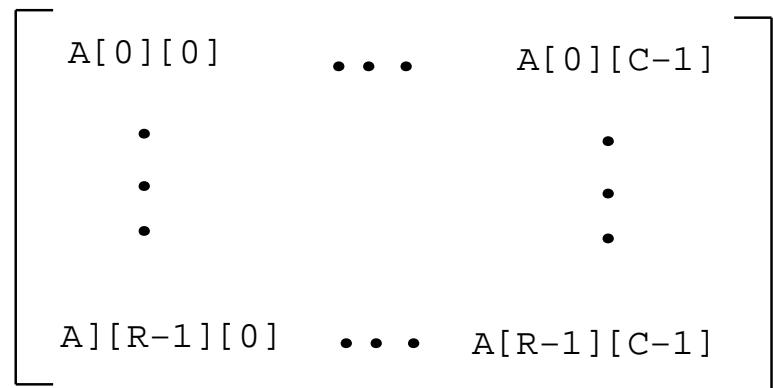
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

Array Size: $R * C * K$ bytes

Arrangement: Row-Major ordering (guaranteed)

Row major order means the elements are stored in the following order:

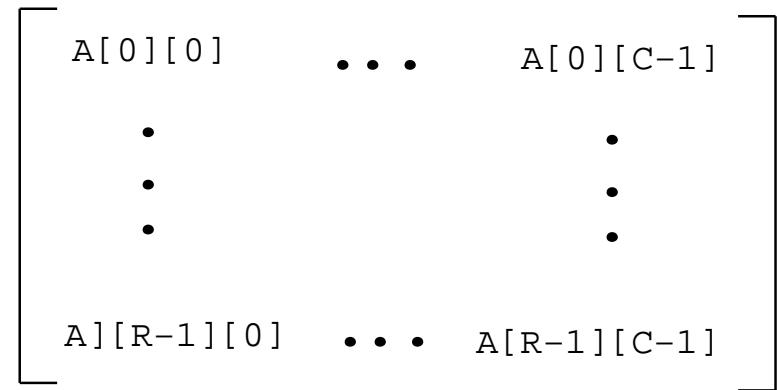
$$[A_{0,0}, \dots, A_{0,C-1}, A_{1,0}, \dots, A_{1,C-1}, \dots, A_{R-1,0}, \dots, A_{R-1,C-1}].$$



Multidimensional Array Access

Declaration: `T A [R] [C] ;`

- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

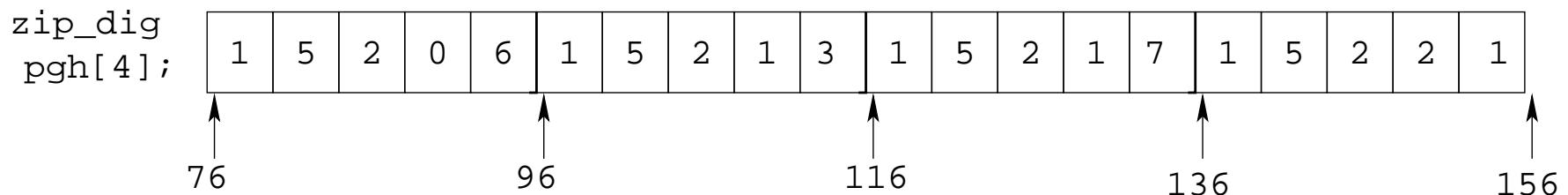


To access element $A[i][j]$, perform the following computation:

$$A + i * C * K + j * K$$

Nested Array Example

```
#define PCOUNT 4
zip_dig pgm[PCOUNT] =
{{1, 5, 2, 0, 6},
 {1, 5, 2, 1, 3},
 {1, 5, 2, 1, 7},
 {1, 5, 2, 2, 1}};
```



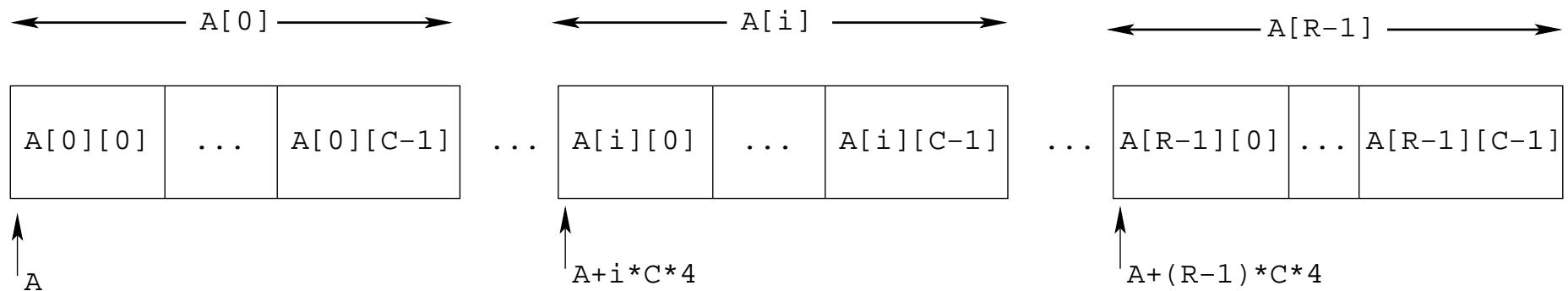
- Declaration “`zip_dig pgm[4]`” is equivalent to “`int pgm[4][5]`.”
- Variable pgm denotes an array of 4 elements allocated contiguously.
- Each element is an array of 5 ints, which are allocated contiguously.
- This is “row-major” ordering of all elements, guaranteed.

Nested Array Row Access

Row Vectors:

Given a nested array declaration $T \ A[R][C]$, you can think of this as an array of arrays.

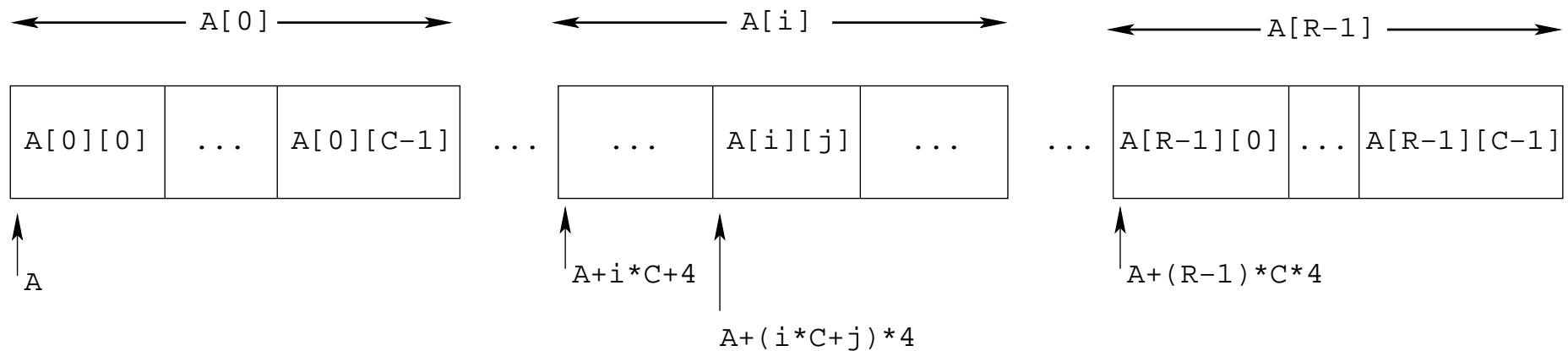
- $A[i]$ is an array of C elements.
- Each element of $A[i]$ has type T , and requires K bytes.
- The starting address of $A[i]$ is $A + i * C * K$.



Nested Array Element Access

Array Elements

- $A[i][j]$ is an element of type T , which requires K bytes.
- The address is $A + (i * C + j) * K$.

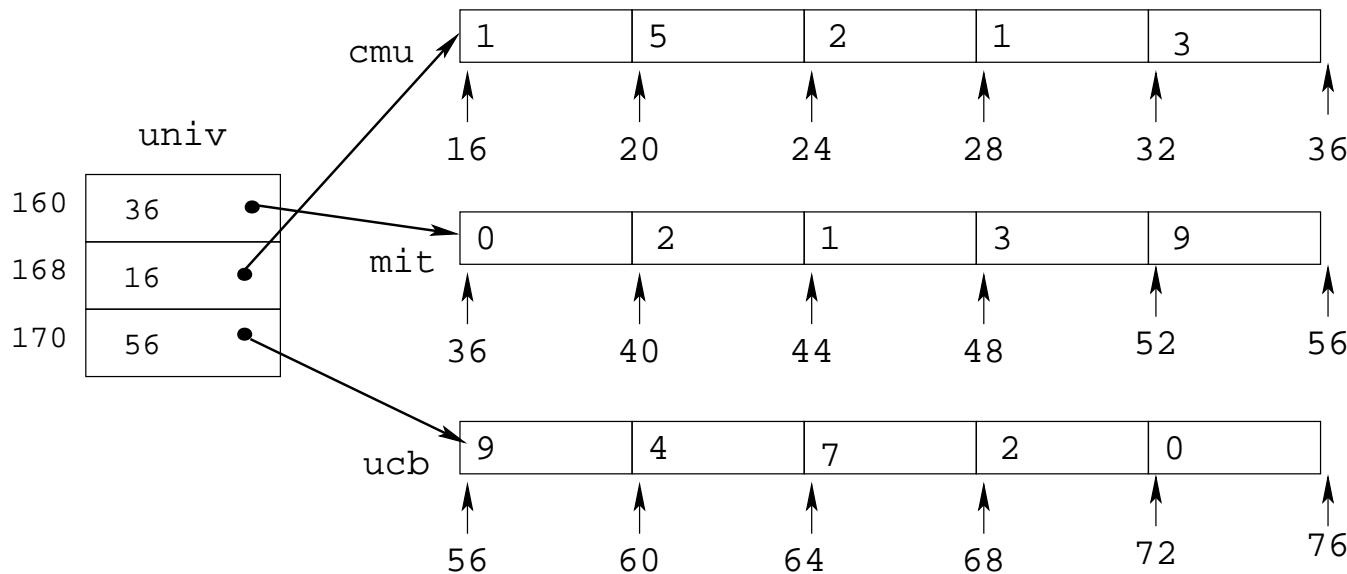


Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };

#define UCOUNT 3
int *univ [UCOUNT]
    = {mit, cmu, ucb};
```

- Variable `univ` denotes an array of 3 elements.
- Each element is a pointer (8 bytes).
- Each pointer points to an array of ints (may vary in length; i.e., ragged array is possible).



Element Access in a Multi-Level Array

Computation

```
int get_univ_digit
    (size_t index, size_t dig
     )
{
    return univ[index][dig];
}
```

- Element access
 $\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{dig}]$
- Must do two memory reads:
 - First get pointer to row array.
 - Then access element within the row.

```
salq    $2, %rsi          # 4*dig
addq    univ(%rdi,8),%rsi # p = univ[index] + 4*dig
movl    (%rsi), %eax      # return *p
ret
```

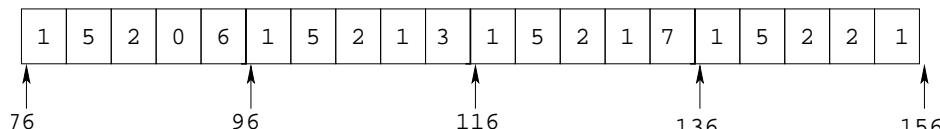
Array Element Accesses

Nested Array

```
int get_pgh_digit
    (size_t index,
     size_t dig)
{
    return pgh[index][dig];
}
```

Element at

$\text{Mem}[\text{pgh} + 20 * \text{index} + 4 * \text{dig}]$

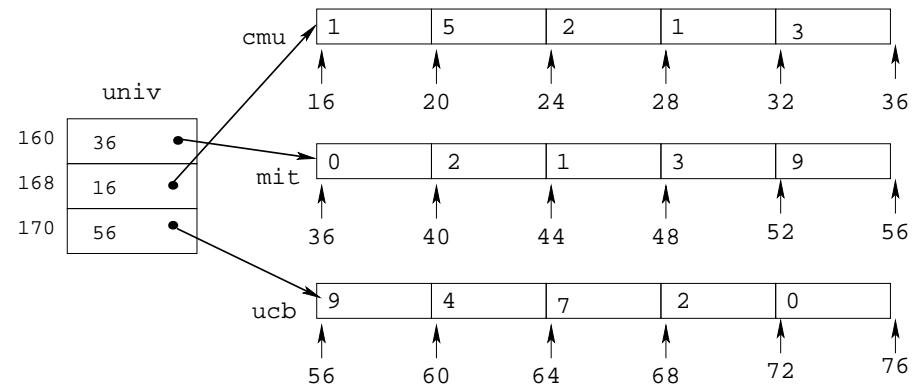


Multi-Level Array

```
int get_univ_dig
    (size_t index,
     size_t dig)
{
    return univ[index][dig];
}
```

Element at

$\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{dig}]$



Similar C references, but different address computations.

$N \times N$ Matrix Code

Fixed dimensions:

Know value of N at compile time.

Variable dimensions,

explicit indexing:

Traditional way to implement dynamic arrays

Variable dimensions,

implicit indexing:

Now supported by gcc

```
#define N 16
typedef int fix_matrix[N][N];
/* Get element a[i][j] */
int fix_ele( fix_matrix a,
             size_t i, size_t j ) {
    return a[i][j];
}
```

```
#define IDX(n, i, j) ((i) * (n) + (j))
/* Get element a[i][j] */
int vec_ele( size_t n, int *a,
             size_t i, size_t j ) {
    return a[IDX(n, i, j)];
}
```

```
/* Get element a[i][j] */
int var_ele( size_t n, int a[n][n],
             size_t i, size_t j ) {
    return a[i][j];
}
```

16 × 16 Matrix Access

Array Elements

- Address $A + i * (C * K) + j * K$
- C = 16, K = 4

```
/* Get element a[i][j] */  
int fix_ele( fix_matrix a, size_t i, size_t j ) {  
    return a[i][j];  
}
```

```
# a in %rdi, i in %rsi, j in %rdx  
salq $6, %rsi          # 64*i  
addq %rsi, %rdi        # a + 64*i  
movl (%rdi, %rdx, 4), %eax # M[a + 64*i + 4*j]
```

$n \times n$ Matrix Access

Array Elements

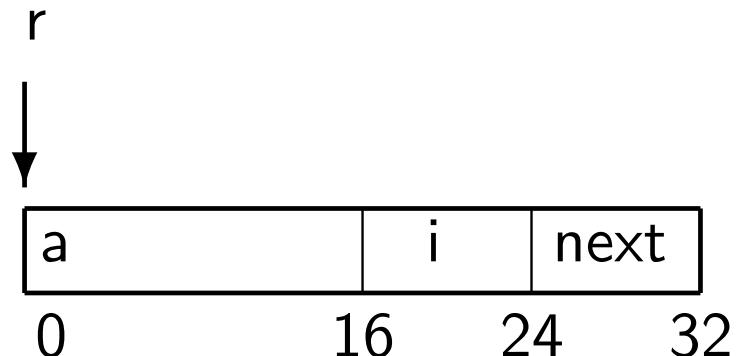
- Address $A + i * (C * K) + j * K$
- $C = n, K = 4$
- Must perform integer multiplication

```
/* Get element a[i][j] */  
int var_ele( size_t n, int a[n][n], size_t i, size_t j  
)  
{  
    return a[i][j];  
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx  
imulq    %rdx, %rdi          # n*i  
leaq     (%rsi, %rdi, 4), %rax   # a + 4*n*i  
movl     (%rax, %rcx, 4), %eax   # a + 4*n*i + 4*j  
ret
```

Structure Representation

```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



Structure represented as block of memory

- Big enough to hold all the fields

Fields ordered according to declaration

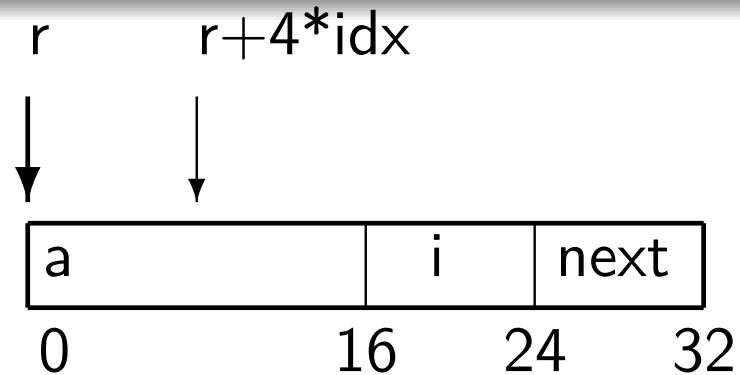
- Even if another ordering could yield a more compact representation

Compiler determines overall size and position of fields

- Machine-level program has no understanding of the structures in the source code.

Generating Pointer to Structure Member

```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



```
int *get_ap  
(struct rec *r, size_t idx)  
{  
    return &r->a[idx];  
}
```

Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as `r + 4*idx`

```
# r in %rdi, idx in %rsi  
leaq (%rdi, %rsi, 4), %rax  
ret
```

BTW: why does `r->i` need 8 bytes? Alignment. (Next slide set)

Aside on Structures: Arrow vs. Dot



If you have a *pointer r* to a structure, use $r \rightarrow x$ to access component x.

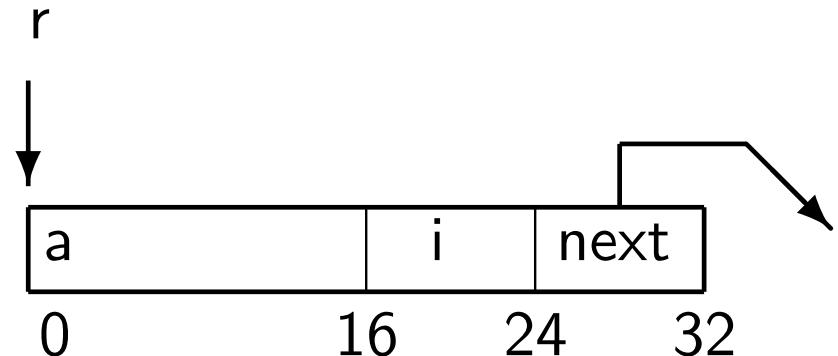
If you have the structure s itself, use $s.x$.

$r \rightarrow x$ is just syntactic sugar for $(*r).x$

Following Linked List

```
void set_val
    (struct rec *r, int val)
{
    while (r) {
        int i = r->i;
        r->a[i] = val;
        r = r->next;
    }
}
```

```
struct rec {
    int a[4];
    size_t i;
    struct rec *next;
};
```



```
.L11:                                # loop:
    testq   %rdi, %rdi                 # Test r
    je      .L12                      # if = 0, goto done
    movq    16(%rdi), %rax            # i = M[r+16]
    movl    %esi, (%rdi, %rax, 4)     # M[r+4*i] = val
    movq    24(%rdi), %rdi            # r = M[r+24]
    jmp     .L11                      # goto loop
.L12:                                # done:
```