

Large Scale Distributivity

We have two little distributivity tautologies:

$$(\mathbf{j} \vee (\mathbf{y} \wedge \mathbf{c})) \Leftrightarrow ((\mathbf{j} \vee \mathbf{y}) \wedge (\mathbf{j} \vee \mathbf{c}))$$

and

$$(\mathbf{j} \wedge (\mathbf{y} \vee \mathbf{c})) \Leftrightarrow ((\mathbf{j} \wedge \mathbf{y}) \vee (\mathbf{j} \wedge \mathbf{c})).$$

What happens when we have more than three sentences?

We simply need to use the rules several times.

Let's deal with the "or-over-and" case from above first. In particular, this applies in the reduction to CNF if, on the way to a conjunction of disjunctive clauses, you actually have a disjunction of conjunctive clauses (that we could call Disjunctive Normal Form). So we get

$$\begin{aligned} ((\mathbf{j} \wedge \mathbf{q}) \vee (\mathbf{y} \wedge \mathbf{c})) &\Leftrightarrow (((\mathbf{j} \wedge \mathbf{q}) \vee \mathbf{y}) \wedge ((\mathbf{j} \wedge \mathbf{q}) \vee \mathbf{c})) \\ &\Leftrightarrow (((\mathbf{j} \vee \mathbf{y}) \wedge (\mathbf{q} \vee \mathbf{y})) \wedge ((\mathbf{j} \vee \mathbf{c}) \wedge (\mathbf{q} \vee \mathbf{c}))) \\ &\Leftrightarrow (\mathbf{j} \vee \mathbf{y}) \wedge (\mathbf{q} \vee \mathbf{y}) \wedge (\mathbf{j} \vee \mathbf{c}) \wedge (\mathbf{q} \vee \mathbf{c}) \end{aligned}$$

(Notice in the final line I shamelessly removed the parenthesis to make it clearer.)

Notice what we have in the end: a conjunction of every disjunction of an element of $\{\mathbf{j}, \mathbf{q}\}$ with an element of $\{\mathbf{y}, \mathbf{c}\}$. So what do you think would result with even more – for example $(\mathbf{j}_1 \wedge \mathbf{j}_2 \wedge \dots \wedge \mathbf{j}_m) \vee (\mathbf{y}_1 \wedge \mathbf{y}_2 \wedge \dots \wedge \mathbf{y}_n)$? To help you remember these things, you might want to think of an analogous arithmetic identity. You know that

$$a \cdot (b + c) = (a \cdot b) + (a \cdot c)$$

And you also know that

$$(a + d) \cdot (b + c) = (a \cdot b) + (a \cdot c) + (d \cdot b) + (d \cdot c).$$

Notice you have the sum of every product of an element of $\{a, d\}$ with an element of $\{b, c\}$. (Recall in the logical version you had a conjunction of every disjunction of an element of $\{\mathbf{j}, \mathbf{q}\}$ with an element of $\{\mathbf{y}, \mathbf{c}\}$.) That ought to be a memory clue and also a major suggest of what you should get for the large version $(\mathbf{j}_1 \wedge \mathbf{j}_2 \wedge \dots \wedge \mathbf{j}_m) \vee (\mathbf{y}_1 \wedge \mathbf{y}_2 \wedge \dots \wedge \mathbf{y}_n)$.

So now push yourself:

1. What is the CNF of $(\mathbf{j}_1 \wedge \mathbf{j}_2 \wedge \dots \wedge \mathbf{j}_m) \vee (\mathbf{y}_1 \wedge \mathbf{y}_2 \wedge \dots \wedge \mathbf{y}_n)$?

$$\begin{aligned} &(\mathbf{j}_1 \vee \mathbf{y}_1) \wedge (\mathbf{j}_1 \vee \mathbf{y}_2) \wedge \dots \wedge (\mathbf{j}_1 \vee \mathbf{y}_n) \\ &\wedge (\mathbf{j}_2 \vee \mathbf{y}_1) \wedge (\mathbf{j}_2 \vee \mathbf{y}_2) \wedge \dots \wedge (\mathbf{j}_2 \vee \mathbf{y}_n) \\ &\wedge \dots \wedge (\mathbf{j}_m \vee \mathbf{y}_1) \wedge (\mathbf{j}_m \vee \mathbf{y}_2) \wedge \dots \wedge (\mathbf{j}_m \vee \mathbf{y}_n) \end{aligned}$$

2. What should you get if you reverse the conjunctions and disjunctions? That is, what is equivalent to $(\mathbf{j} \vee \mathbf{q}) \wedge (\mathbf{y} \vee \mathbf{c})$? (Does the arithmetic analog still hold?)

$$\begin{aligned} ((\mathbf{j} \vee \mathbf{q}) \wedge (\mathbf{y} \vee \mathbf{c})) &\Leftrightarrow (((\mathbf{j} \vee \mathbf{q}) \wedge \mathbf{y}) \vee ((\mathbf{j} \vee \mathbf{q}) \wedge \mathbf{c})) \\ &\Leftrightarrow (((\mathbf{j} \wedge \mathbf{y}) \vee (\mathbf{q} \wedge \mathbf{y})) \vee ((\mathbf{j} \wedge \mathbf{c}) \vee (\mathbf{q} \wedge \mathbf{c}))) \\ &\Leftrightarrow (\mathbf{j} \wedge \mathbf{y}) \vee (\mathbf{q} \wedge \mathbf{y}) \vee (\mathbf{j} \wedge \mathbf{c}) \vee (\mathbf{q} \wedge \mathbf{c}) \end{aligned}$$

The arithmetic analog holds but now addition is analogous to disjunction and multiplication is analogous to conjunction.

3. What is the DNF of $(\mathbf{j}_1 \vee \mathbf{j}_2 \vee \dots \vee \mathbf{j}_m) \wedge (\mathbf{y}_1 \vee \mathbf{y}_2 \vee \dots \vee \mathbf{y}_n)$?

$$\begin{aligned} &(\mathbf{j}_1 \wedge \mathbf{y}_1) \vee (\mathbf{j}_1 \wedge \mathbf{y}_2) \vee \dots \vee (\mathbf{j}_1 \wedge \mathbf{y}_n) \\ &\vee (\mathbf{j}_2 \wedge \mathbf{y}_1) \wedge (\mathbf{j}_2 \wedge \mathbf{y}_2) \vee \dots \vee (\mathbf{j}_2 \wedge \mathbf{y}_n) \\ &\vee \dots \vee (\mathbf{j}_m \wedge \mathbf{y}_1) \vee (\mathbf{j}_m \wedge \mathbf{y}_2) \vee \dots \vee (\mathbf{j}_m \wedge \mathbf{y}_n) \end{aligned}$$

4. What is the CNF of $(\mathbf{j}_1^1 \wedge \mathbf{j}_2^1 \wedge \dots \wedge \mathbf{j}_{m_1}^1) \vee (\mathbf{j}_1^2 \wedge \mathbf{j}_2^2 \wedge \dots \wedge \mathbf{j}_{m_2}^2) \vee \dots \vee (\mathbf{j}_1^n \wedge \mathbf{j}_2^n \wedge \dots \wedge \mathbf{j}_{m_n}^n)$?

It is the conjunction of $m_1 m_2 \dots m_n$ disjunctive clauses. Each of those clauses has the disjunction of one element of $\{\mathbf{j}_1^1 \mathbf{j}_2^1, \dots, \mathbf{j}_{m_1}^1\}$, one element of $\{\mathbf{j}_1^2 \mathbf{j}_2^2, \dots, \mathbf{j}_{m_2}^2\}$, ..., and one element of $\{\mathbf{j}_1^n \mathbf{j}_2^n, \dots, \mathbf{j}_{m_n}^n\}$.

5. What is the DNF of $(\mathbf{j}_1^1 \vee \mathbf{j}_2^1 \vee \dots \vee \mathbf{j}_{m_1}^1) \wedge (\mathbf{j}_1^2 \vee \mathbf{j}_2^2 \vee \dots \vee \mathbf{j}_{m_2}^2) \wedge \dots \wedge (\mathbf{j}_1^n \vee \mathbf{j}_2^n \vee \dots \vee \mathbf{j}_{m_n}^n)$?

It is the disjunction of $m_1 m_2 \dots m_n$ conjunctive clauses. Each of those clauses has the conjunction of one element of $\{\mathbf{j}_1^1 \mathbf{j}_2^1, \dots, \mathbf{j}_{m_1}^1\}$, one element of $\{\mathbf{j}_1^2 \mathbf{j}_2^2, \dots, \mathbf{j}_{m_2}^2\}$, ..., and one element of $\{\mathbf{j}_1^n \mathbf{j}_2^n, \dots, \mathbf{j}_{m_n}^n\}$.