

Recap Sum-Check Protocol

Given a func. $\Phi: \mathbb{F}^m \rightarrow \mathbb{F}$ of $\deg \leq d$

HCF, Verify that

$$\sum_{\vec{x} \in H^m} \Phi(\vec{x}) = c \quad \text{where } c \in \mathbb{F}$$

Define $g: [\Phi]: \mathbb{F}^i \rightarrow \mathbb{F}$

$$g: [\Phi](x_1, \dots, x_i) = \sum_{a_{i+1}, \dots, a_m \in H} \Phi(x_1, \dots, x_i, a_{i+1}, \dots, a_m)$$

Observe:

* $\deg g: [\Phi] \leq \deg \Phi$

* $g: [\Phi](x_1, \dots, x_i) = \sum_{a \in H} g_{i+1}: [\Phi](x_1, \dots, x_i, a)$

Sum-Check Lemma

Given a $\deg \leq d$ poly $\Phi: \mathbb{F}^m \rightarrow \mathbb{F}$, a set $H \subset \mathbb{F}$, $c \in \mathbb{F}$,
(denote $G_0 = c$, $G_m = \Phi$)

I. If $\sum_{\vec{x} \in H^m} \Phi(\vec{x}) = c \Rightarrow \exists G_1, \dots, G_{m-1} \quad G_i: \mathbb{F}^i \rightarrow \mathbb{F}$

s.t. $\forall i, r_1, \dots, r_i \in H \quad G_i(r_1, \dots, r_i) = \sum_{a \in H} G_{i+1}(r_1, \dots, r_i, a)$

II If $\sum_{\vec{x} \in H^m} \Phi(\vec{x}) \neq c \Rightarrow \forall G_1, \dots, G_{m-1} \deg \leq d$

$$P(\forall i, r_1, \dots, r_i \in H \quad G_i(r_1, \dots, r_i) \neq \sum_{a \in H} G_{i+1}(r_1, \dots, r_i, a)) \leq \frac{d}{|\mathbb{F}|}$$

Proof I - Obvious.

II - there must be i for which $G_i \neq g_i[\Phi]$.

Let i be the maximal such i . So,

$$G_{i+1} \equiv g_{i+1}[\Phi], \text{ i.e., } \sum_{a \in H} G_{i+1}(x_1, \dots, x_i, a) \equiv g_i(x_1, \dots, x_i),$$

so we're testing whether $G_i(\mathbb{F}) \stackrel{?}{=} g_i(\mathbb{F})$,

and the lemma follows from Schwartz-Zippel. \square

Observe that the number of queries into $\Phi, G_1, \dots, G_{m-1}$ is $\leq m \cdot (|H|+2)$ (rather than $|H|^m$!).

Low Degree Extension

Given $H \subset \mathbb{F}$, $f: H^m \rightarrow \mathbb{F}$

There exists a unique $\deg < |H|$ in each var extension $\hat{f}: \mathbb{F}^m \rightarrow \mathbb{F}$.

Thm $NPC \subseteq P[P(O(\log n), \text{poly} \log n)]$

Proof We start with the NPhard gap-ZT. \rightarrow

Assume an instance $P_1, \dots, P_s \in \mathbb{F}[x_1, \dots, x_n]$
(defined by coefficients)

Step 1 Recast problem into more dimensions.

Denote p_{ijk}^t - coeff. of $x_i x_j x_k$ in P_t .

(some of i, j, k may be equal. $x_0 \equiv 1$)

can view (p_{ijk}^t) as a func. from $\{0, \dots, n\}^3$ to \mathbb{F} ,
or as a func. from $(H^m)^3$ to \mathbb{F} .

gap-ZT Given $P_1, \dots, P_s \in \mathbb{F}[x_1, \dots, x_n]$
 $\deg P_i \leq 3$. Decide between:

- 1) $\exists a_1, \dots, a_n \in \mathbb{F}$ s.t.
 $P_t(a_1, \dots, a_n) = 0$
- 2) $\forall a_1, \dots, a_n \in \mathbb{F}$
 $P_t(a_1, \dots, a_n) = 0$ for $\leq \epsilon$

$|\mathbb{F}| = \Theta((\log n)^{10})$

Verifier

Input $P_1, \dots, P_s \in \mathbb{F}[x_1, \dots, x_n]$

1) Choose random $t \in [s]$

2) Compute new rep. HCF $(H) = \log(n+1)$
 $m = \frac{\log(n+1)}{\log \log(n+1)}$

Fix mapping $\{0, \dots, n\}$ to H^m .

P_t is now viewed as a fun. $(H^m)^3 \rightarrow \mathbb{F}$.

Compute Q_t $(\mathbb{F}^m)^3 \rightarrow \mathbb{F}$ extension of P_t
deg $< H$ in each var

Honest Prover

Let $\vec{a} = (a_1, a_2, \dots, a_n)$ denote the common zero of p_1, \dots, p_m .

Can view \vec{a} as $A: \{0, \dots, n\} \rightarrow \mathbb{F}$

also as $A: \mathbb{H}^m \rightarrow \mathbb{F}$. Let \hat{A} be its LDE

$$\sum_{\bar{u}, \bar{v}, \bar{w} \in \mathbb{H}^m} Q(\bar{u}, \bar{v}, \bar{w}) \hat{A}(\bar{u}) \hat{A}(\bar{v}) \hat{A}(\bar{w}) = p^+(a_1, \dots, a_n)$$

Denote $\Phi(\bar{u}, \bar{v}, \bar{w}) = Q(\bar{u}, \bar{v}, \bar{w}) \hat{A}(\bar{u}) \hat{A}(\bar{v}) \hat{A}(\bar{w})$

Verifier

3) (a) Select at random $\bar{u}, \bar{v}, \bar{w} \in \mathbb{F}^m$

check if $\Phi(\bar{u}, \bar{v}, \bar{w}) \stackrel{?}{=} Q(\bar{u}, \bar{v}, \bar{w}) \hat{A}(\bar{u}) \hat{A}(\bar{v}) \hat{A}(\bar{w})$

(b) Select at random $r_1, \dots, r_{3m} \in \mathbb{F}$,

check if $\forall i \in \{0, \dots, 3m-1\}$

$$G_i(r_1, \dots, r_i) \stackrel{?}{=} \sum_{a \in \mathbb{H}} G_{i+1}(r_1, \dots, r_i, a) \quad (\text{Denote } G_0 \equiv 0, G_m \equiv \Phi)$$

4) Test that $A, \Phi, G_1, \dots, G_{3m-1}$ are of low degree.

(not counting Idt)

#random bits

$$\log S + 3m \log |\mathbb{F}| + 3m \log |\mathbb{F}| = O(\log n)$$

#queries to proof

$$4 \log |\mathbb{F}| + 3m \cdot (H+1) \log |\mathbb{F}| = O(\log^2 n)$$

Completeness If there are a_1, \dots, a_n s.t. $P_+(p_+(a_1, \dots, a_n) = 0) = 1$, then there are $\{\Phi^+\}, \{\beta_i^+\}, \{\beta_{3m-1}^+\}, \hat{A}$ s.t. $P(\text{Ver accepts}) = 1$.

Soundness Lemma Assume that for all $a_1, \dots, a_n \in \mathbb{F}$, $P_+(p_+(a_1, \dots, a_n) = 0) < \frac{1}{2}$. Then: for all low degree functions \hat{A} , Φ^+ , β_i^+ s.t. $\deg \hat{A} \leq m(H-1)$, $\deg \Phi^+ \leq 3m(H-1)$, $\deg \beta_i^+ \leq 3m(H-1)$, $P(\text{Ver accepts}) < \text{soundness}$ (some constant).

Proof With prob. $> \frac{1}{2}$ over the choice of \pm ,

$$P_+(\hat{A}(\vec{u}), \dots, \hat{A}(\vec{w})) \neq 0 \quad (*)$$

↑
the point that \pm was mapped to

Assume that (*) is the case.

Let Q be the LDE of p_+ .

Observe that $\deg Q(\vec{u}, \vec{v}, \vec{w}) \hat{A}(\vec{u}) \hat{A}(\vec{v}) \hat{A}(\vec{w}) \leq \deg Q + 3 \deg \hat{A} < 3m(H) + 3m(H) = 6m(H)$

Since $\deg \Phi < 6m(H)$, if $\Phi(\vec{x}, \vec{y}, \vec{z}) \neq Q(\vec{x}, \vec{y}, \vec{z}) \hat{A}(\vec{x}) \hat{A}(\vec{y}) \hat{A}(\vec{z})$,

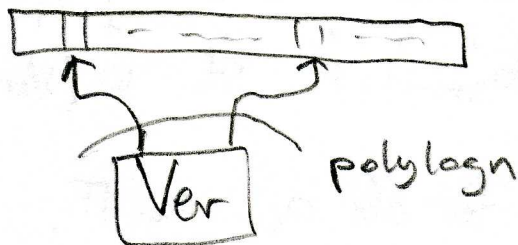
step 3a) passes w.p. $< \frac{6m(H)}{|\mathbb{F}|}$

Assume " $\Phi \equiv Q \hat{A} \hat{A} \hat{A}$ ". We know that $0 \neq P_+(\hat{A}(u) - \hat{A}(w)) = \sum_{\vec{u}, \vec{v}, \vec{w} \in \mathbb{F}^m} \Phi(\vec{u}, \vec{v}, \vec{w})$ assuming $\hat{A}(\vec{0}) = 1$, by the Sum check Lemma,

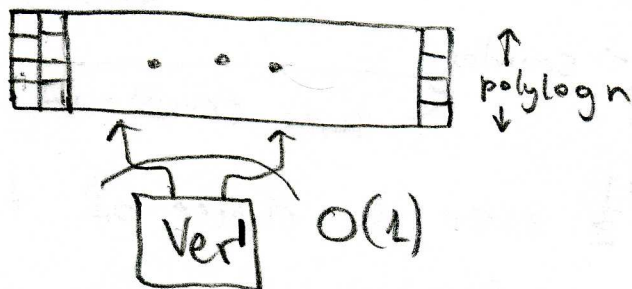
$$P(\text{Ver accepts}) < \frac{6m(H)}{|\mathbb{F}|}$$

Total acceptance prob. $\leq \frac{1}{2} + O\left(\frac{m(H)}{|\mathbb{F}|}\right) < \frac{2}{3}$ for large enough \mathbb{F} .

So far



We will next show



Note
this
direction
is trivial

$$\underline{\text{Thm}} \quad \text{NPC} \subseteq \text{PCP}[O(\log n), \text{polylogn}]_{\mathbb{F}_{0,1}} \Rightarrow \text{NPC} \subseteq \text{PCP}[O(\log n), O(1)]_{\mathbb{F}_{0,1}}^{\text{polylogn}}$$

How to prove $\text{NPC} \subseteq \text{PCP}[O(\log n), O(1)]_{\mathbb{F}_{0,1}}$?] Motivation
By composition.

Def A degree- d polynomial curve is a function $\gamma: \mathbb{F} \rightarrow \mathbb{F}^m$,
 $\gamma = (\gamma_1, \dots, \gamma_m)$, $\gamma_i: \mathbb{F} \rightarrow \mathbb{F}$, s.t. $\deg \gamma_i \leq d$.

Example $t \mapsto \gamma \rightarrow (t^3, 3t-1, 2t^2+4t-1, 7)$

γ can be written as $\gamma(t) = \bar{a}_0 + \bar{a}_1 t + \bar{a}_2 t^2 + \dots + \bar{a}_d t^d$

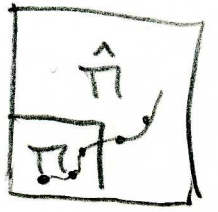
* If $f: \mathbb{F}^m \rightarrow \mathbb{F}$ is a degree $\leq r$ function, then $g: \mathbb{F} \rightarrow \mathbb{F}$ defined by $g(t) := f(\gamma(t))$ is of $\deg \leq r \cdot d$.

* Given points $\bar{x}_1, \dots, \bar{x}_k \in \mathbb{F}^m$, there is a unique $\deg < k$ curve $\gamma: \mathbb{F} \rightarrow \mathbb{F}^m$ s.t. $\gamma(i) = \bar{x}_i$. (assuming $\{1, \dots, k\} \subset \mathbb{F}$)

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Ver'

Let r' = random bits
 ① Simulate Ver on the input and compute indices i_1, \dots, i_q ($q = \text{poly}(\log n)$) to read from π .
 compute $\vec{x}_{i_1}, \dots, \vec{x}_{i_q} \in \mathbb{F}^m$ corr. points.



② compute γ - the $\text{deg} < q$ curve s.t. $\gamma(i) = \vec{x}_i$
 Ask the prover for $g: \mathbb{F} \rightarrow \mathbb{F}$, supposedly $g(t) = \hat{\pi}(\gamma(t))$

$$|H| = \log |\mathbb{F}|$$

$$m = \frac{\log |\mathbb{F}|}{\log \log |\mathbb{F}|}$$

$$|\mathbb{F}| = |H|^{\text{constant}}$$

Tests:

① Select random $r \in \mathbb{F}$ and check $g(r) = \hat{\pi}(\gamma(r))$

② Eval. $g(a), g(b), \dots$ and accept iff Ver would have accepted.

③ Test $\hat{\pi}$ for low degree.

$$\text{deg } g = \text{deg}(\hat{\pi}) \cdot \text{deg}(\gamma) \leq m \cdot |H| \cdot q$$

to describe g , we need $m \cdot |H| \cdot q \log |\mathbb{F}| = \text{poly}(\log n)$ bits.

Completeness Clearly $\hat{\pi}$ can be computed from π . Set $g := \hat{\pi} \circ \gamma$ for all γ .

Soundness Suppose that $\forall \pi \ P(\text{Ver acc}) < \frac{1}{2}$.

Lemma Then for any $\text{deg } m \cdot |H| \ \hat{\pi}: \mathbb{F}^m \rightarrow \mathbb{F}$ and any poly $\{\gamma\}_r$

$$P(\text{Ver}' \text{ acc}) < \frac{1}{2} + \frac{m \cdot |H| \cdot q}{|\mathbb{F}|}$$

Proof If r' is s.t. Ver rejects, then $\hat{\pi}(\gamma_{r'}(1), \dots, \gamma_{r'}(q))$ cause Ver to reject. So Ver' only acc. if $g \neq \hat{\pi} \circ \gamma_{r'}$, but then it acc in Test ① w. prob $\leq \frac{m \cdot |H| \cdot q}{|\mathbb{F}|}$.

□