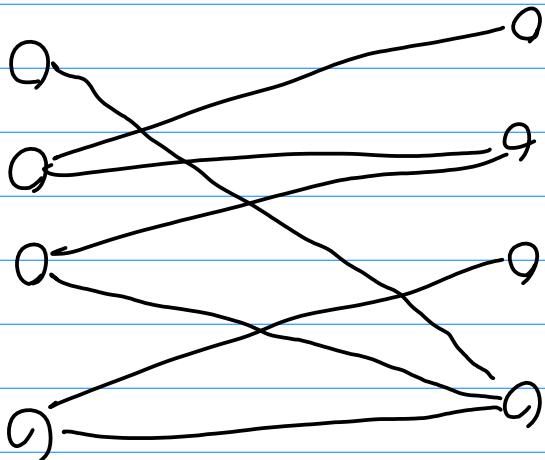


Bipartite Matching (\rightarrow other flow applications)

People

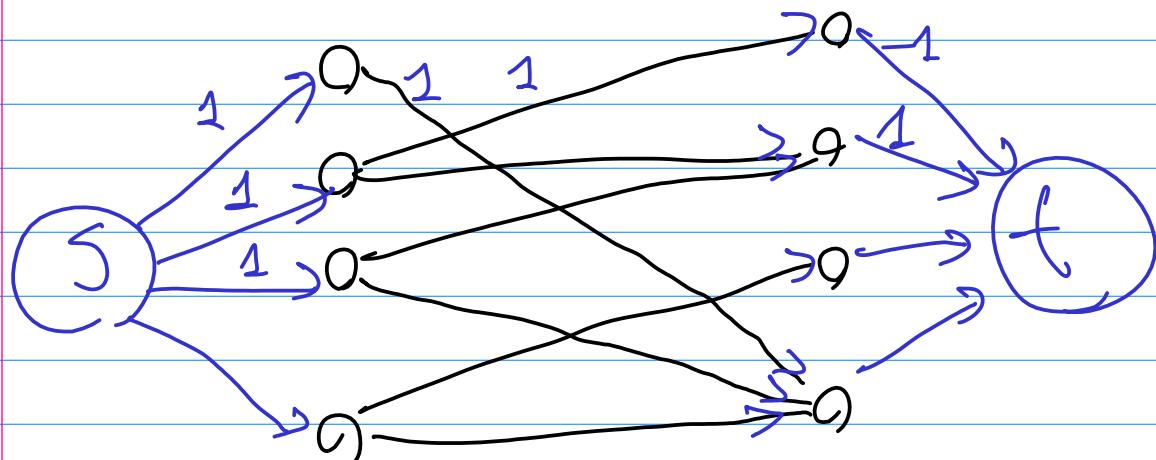
Slots



draw edge if person can make slot

Each person needs one slot
Each slot can take 1 person

Q: how many people can be accommodated?



Answer: max flow, draw S, t ,
edges $S \rightarrow \text{left } t$, $\text{right } t \rightarrow t$,

orient all edges to right.

why?

(1) any matching of size m gives size- m flow

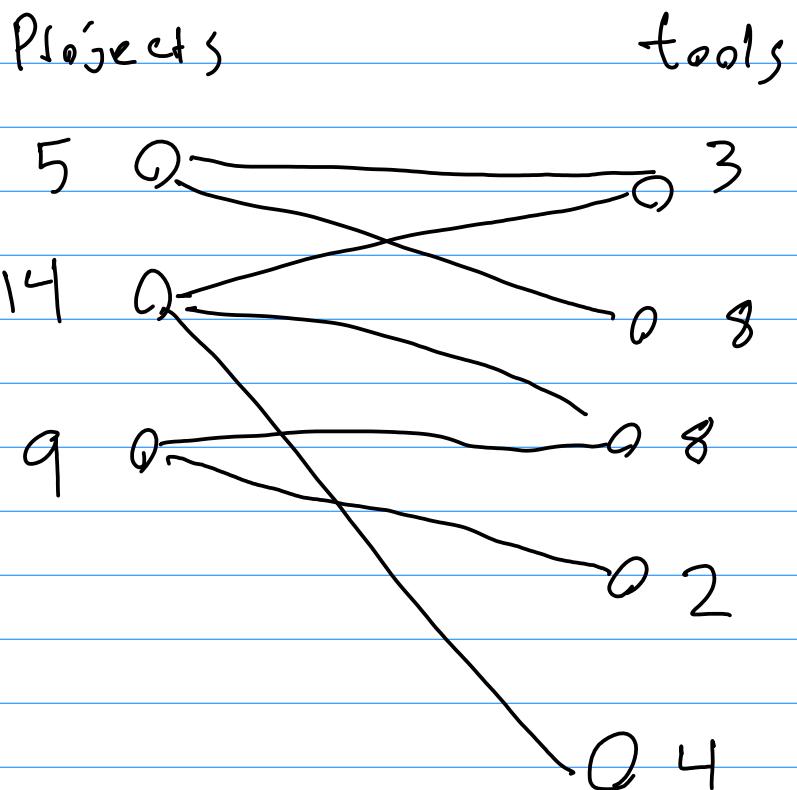
(2) any size- m (integer) flow gives size- m matching
[path decomposition]

\Rightarrow Max matching = max flow

Variants:

- 3 people per slot.
[right \rightarrow edges capacity 3]
- 4 slots per person
[s \rightarrow left capacity 4]

Project Selection



Given: Projects to do, get Paid P_i

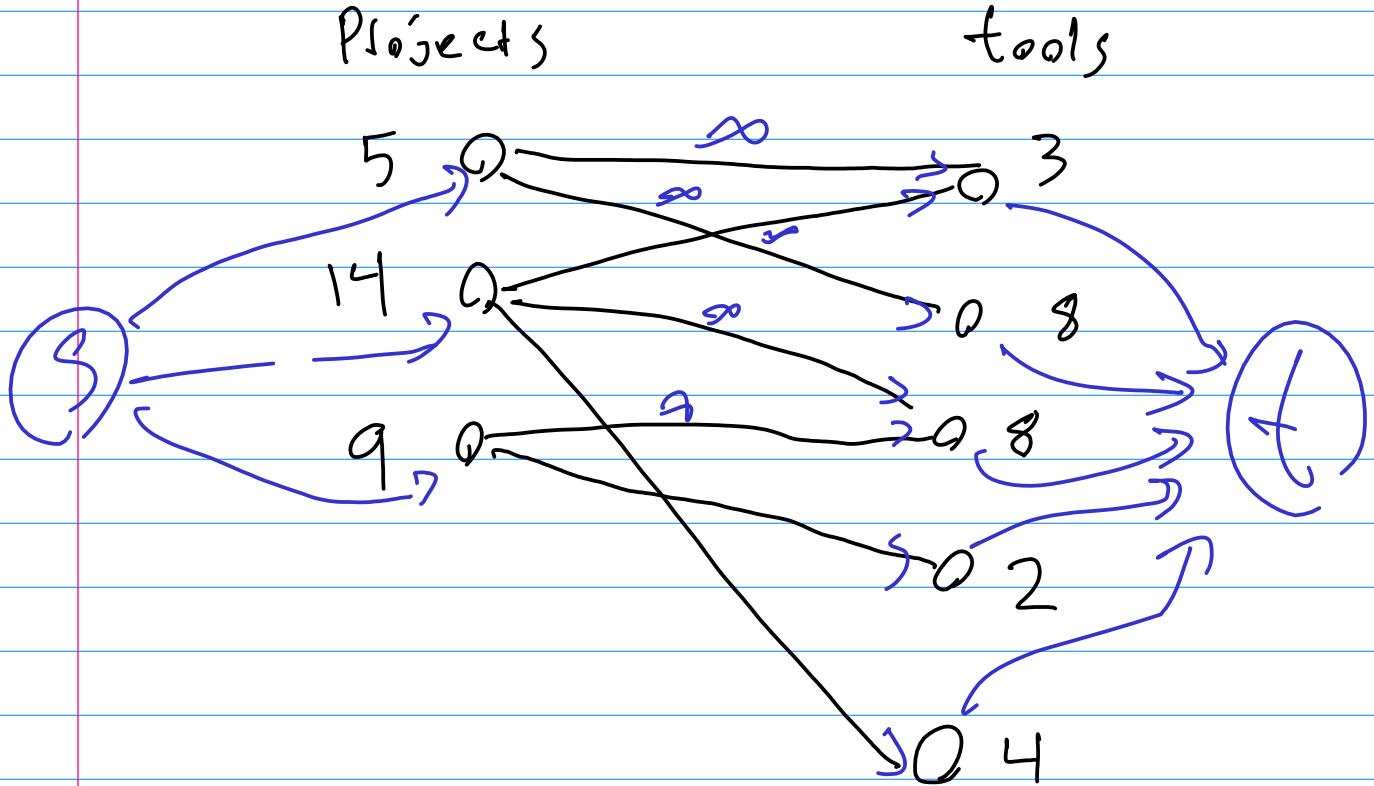
tools you need, Cost C_j

each project needs a subset of tools
tools can be reused.

Q: Choose subset of projects to do
& tools to buy

s.t. each project you do, you buy + tools
& Maximize

- (total value of projects done)
- (total cost of tools bought)



Answer: max flow (\equiv min cut)

Draw edges $S \rightarrow$ project, tools \rightarrow t
 Capacity = given value

∞ capacity on project \rightarrow tool edges

\Rightarrow min cut only cuts $S \rightarrow$ project or tool \rightarrow t edges.

any $S-f$ cut cuts
 - projects you don't complete
 - tools you do buy

Valid $s-t$ cut

\Rightarrow no $s \rightarrow \text{project} \rightarrow \text{tool} \rightarrow t$ path
over uncut edges

\Rightarrow for every (project, tool) pair (P, l)
where P needs l

either $s \rightarrow P$ cut or $l \rightarrow t$ cut

\Rightarrow either P canceled or l bought

\Rightarrow Valid choice to build/buy.

Cost of cut =

$$\text{Cost}(S) = \text{Total Value of canceled projects}$$
$$+ (\text{Total price of bought tools})$$

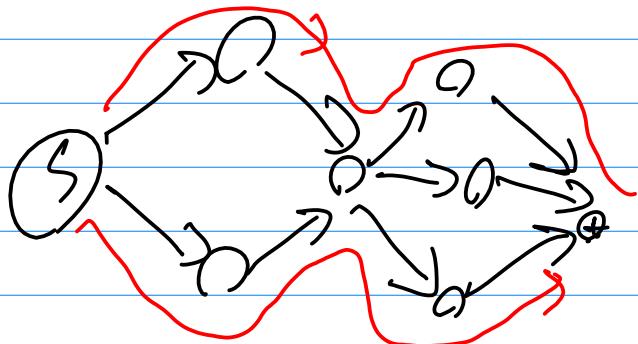
\Rightarrow (total value of all projects)
- (profit made by choice)

independent
of choice

\Rightarrow min cut = max profit

Edge Disjoint Paths:

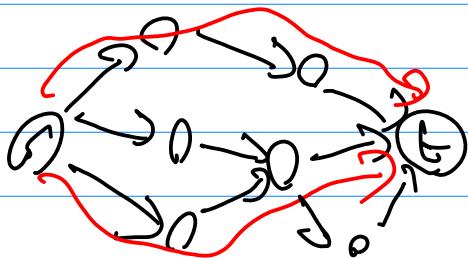
Find max # edge-disjoint s-t paths
in graph



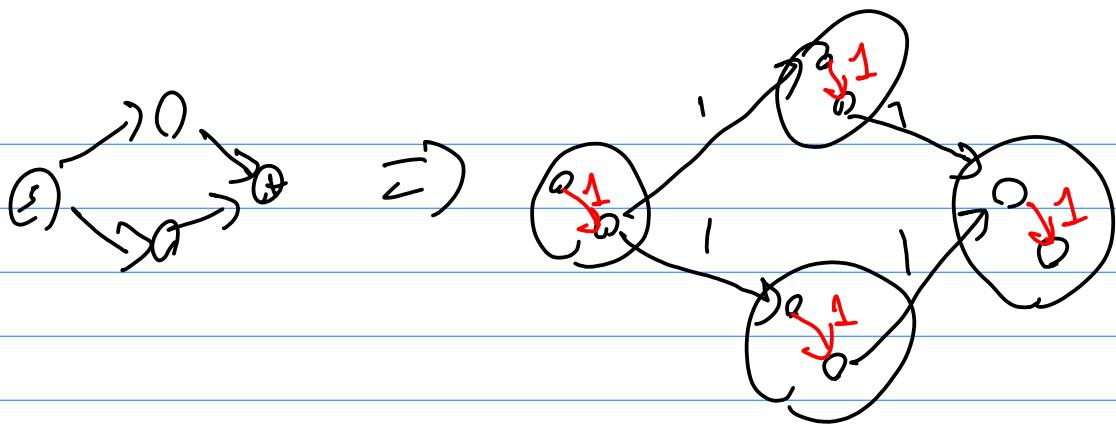
A: Max flow w/ capacity 1

Vertex-disjoint Paths:

Max # vertex-disjoint s-t paths
[reusing s/t is OK]



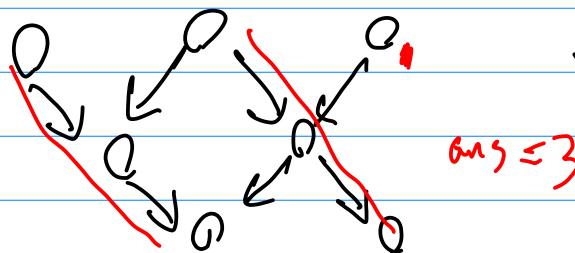
A: Split vertices into "in" - and "out" side



regular edge $(u, v) \Rightarrow (u_{in}, v_{out})$
add (v_{in}, u_{out}) edges & u
all capacity 1
Ans = flow (s_{out}, t_{in})

Min disjoint Path Cover

Given a DAG find min #
vertex-disjoint paths
that cover all vertices



Soln: bipartite matching between in-nodes & out-nodes

- \Rightarrow each vertex gets ≤ 1 selected edge in & out
- \Rightarrow can start at unmatched in-nodes and follow edges to get paths covering graph

$$(\# \text{ paths}) = (\# \text{ unmatched in vertices}) = V - (\# \text{ matched pairs})$$

\Rightarrow min # paths (\Leftarrow max matching pairs)

