

Problem 1 : [7 points] In a protocol for maintaining global topology information, the processes $p[i : 0..n-1]$ exchange messages of the form:

$st(k, vp, r)$

where k is the index of process $p[k]$ that originally generated the message, vp is the local topology information of process $p[k]$ when $p[k]$ generated the message, and r is the number of remaining hops that the message can still make. For simplicity, assume that the exchanged messages in the protocol can be lost but neither corrupted nor reordered. A process in this protocol is specified as follows:

process $p[i:0..n-1]$

const $rmax$

inp $N: \text{set } \{g \mid p[g] \text{ is a neighbor of } p[i]\},$
 $up: \text{array } [N] \text{ of boolean}$

var $net : \text{array } [0..n-1, 0..n-1] \text{ of boolean},$
 $vp : \text{array } [0..n-1] \text{ of boolean},$
 $r : 0..rmax-1,$
 $m : 0..n,$
 $f, h : N,$
 $k : 0..n-1$

par $g : N$

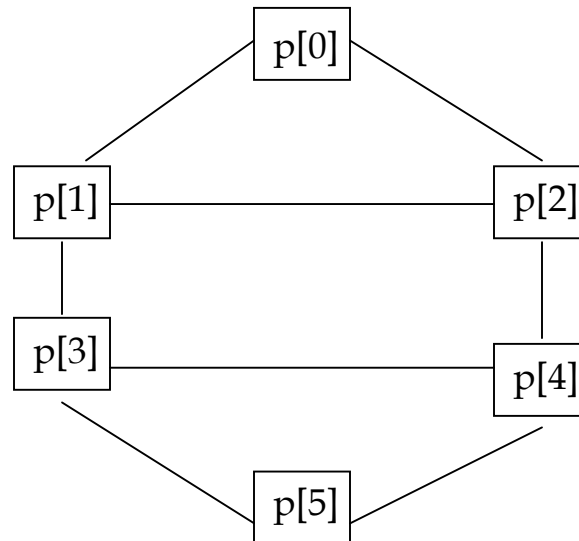
begin

$\text{true} \quad \quad \quad \rightarrow S.0$
[] $\text{rcv } st(k, vp, r) \text{ from } p[g] \quad \rightarrow S.1$

end

Specify statement S.1.

Problem 2 : [6 points] Assume that the protocol described in Problem 1 is to be deployed in a network whose topology is as follows:



What is a good value for r_{max} in this case? Explain your answer.

Problem 3 : [7 points] Two processes p and q circulate a $chk(b)$ message in order to maintain their local topology information according to the protocol in Section 11.2. Each of the two processes has a boolean variable named “up” whose value is true when the two channels between p and q are both up. The two processes have different timeout actions. When p times out, it assigns the value false to its up variable and sends a $chk(0)$ message to q . When q times out, it only assigns the value false to its up variable. Draw the state transition diagram for this protocol under the following assumptions:

- Initially, the two up variables, namely $up.p$ and $up.q$, have the value true, and the two channels, namely $ch.p.q$ and $ch.q.p$, are empty.
 - The channel $ch.p.q$ is up and will stay up indefinitely, and the channel $ch.q.p$ is down and will stay down indefinitely.
 - Messages, that are sent to an up channel, are neither corrupted nor lost until they are received.
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