# Proving and Explaining the Unfeasibility of Message Sequence Charts for Hybrid Systems

#### Alessandro Cimatti Sergio Mover Stefano Tonetta

Fondazione Bruno Kessler

October 31, 2011

#### **Hybrid Systems**

- Mix discrete (e.g. hardware) and continuous (e.g. sensor) behaviors.
- Complex critical systems: train control system (ETCS), airplane traffic control system (TCAS), ...
- Network of components.



#### Scenario-verification

#### Existing approaches:

- Reduction to reachability:
  - Can prove both feasibility and unfeasibility.
  - Inefficient.
- Scenario-based encoding [CAV11]:
  - Cannot prove unfeasibility.
  - Efficient.

- Our contribution is a SMT-based technique that:
  - Efficiently proves unfeasibility.
  - Extracts explanations for the unfeasibility.

#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

- 3 Explanations of Unfeasibility
- Experimental Evaluation
- 5 Conclusions and future work

#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

- 3 Explanations of Unfeasibility
- Experimental Evaluation
- 5 Conclusions and future work

#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

- 3 Explanations of Unfeasibility
- Experimental Evaluation
- 5 Conclusions and future work

## Hybrid Automata

Hybrid automata ([Henzinger 96]):

- Framework for representing hybrid systems.
- Discrete instantaneous mode switches.
- Continuous evolution according to flow conditions.



## Hybrid Automata Network

Network of hybrid automata  $\mathcal{H} = H_1 || \dots || H_n$ :

- Move asynchronously on local events  $(\tau)$ .
- Synchronize on shared events.



#### Different semantics:

- Global-time ([Henzinger 96]).
- Local-time ([Bengstsson 98]).

Sergio Mover (FBK)

#### Local-time semantics

- The time evolves independently in each automaton:
  - Local time scale.
  - The continuous evolution is a local transition.
- The local time of the automata must be the same:
  - On synchronizations.
  - At the end of a run.



 $\tau$  = local event (no stutter or time).

# SMT analysis of Hybrid Systems

- Each automaton is encoded in a symbolic transition system  $H_i = \langle Init_i, Trans_i \rangle$ .
- Bounded model checking:



#### • k-induction.

- Base case: BMC up to k.
- Inductive case: BMC and simple path condition up to k + 1.
- Use SMT solvers as decision procedure.

# Background SMT analysis of Hybrid Syster

Scenario-Verification

- 3 Explanations of Unfeasibility
- Experimental Evaluation
- 5 Conclusions and future work

 $\langle m, \phi \rangle$ : Message sequence chart *m* with constraints  $\phi$ .



- The CMSC is translated in a monitor automaton S<sub>m</sub>.
- The automaton is composed with the network.
- Enables off-the-shelf verification techniques:
  - BMC: feasibility.
  - k-induction: unfeasibility.



- The CMSC is translated in a monitor automaton S<sub>m</sub>.
- The automaton is composed with the network.
- Enables off-the-shelf verification techniques:
  - BMC: feasibility.
  - k-induction: unfeasibility.



- The CMSC is translated in a monitor automaton S<sub>m</sub>.
- The automaton is composed with the network.
- Enables off-the-shelf verification techniques:
  - BMC: feasibility.
  - k-induction: unfeasibility.



- The CMSC is translated in a monitor automaton S<sub>m</sub>.
- The automaton is composed with the network.
- Enables off-the-shelf verification techniques:
  - BMC: feasibility.
  - k-induction: unfeasibility.



Sergio Mover (FBK)

Unfeasibility and Explanations of MSC

- The CMSC is translated in a monitor automaton S<sub>m</sub>.
- The automaton is composed with the network.
- Enables off-the-shelf verification techniques:
  - BMC: feasibility.
  - k-induction: unfeasibility.



- The CMSC is translated in a monitor automaton S<sub>m</sub>.
- The automaton is composed with the network.
- Enables off-the-shelf verification techniques:
  - BMC: feasibility.
  - k-induction: unfeasibility.



#### Scenario-based encoding

- For all the automata:
  - Fix the position of the shared events. transitions are simplified wrt shared event







#### Scenario-based encoding

- For all the automata:
  - Fix the position of the shared events. transitions are simplified wrt shared event
  - Add the synchronization constraints.









#### Scenario-based encoding

- For all the automata:
  - Fix the position of the shared events. transitions are simplified wrt shared event
  - Add the synchronization constraints.
  - Encode the "local segments". transitions are simplified wrt τ





#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

- 3 Explanations of Unfeasibility
- Experimental Evaluation
- 5 Conclusions and future work

	Reduction to reachability	SMT-based approach
Feasibility	BMC Inefficient	Scenario-driven encoding Efficient
Unfeasibility	K-induction Inefficient	Partitioned k-induction Efficient

- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.



- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.



- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





SAT - new states are reachable

- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





SAT - new states are reachable

- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





UNSAT - no new states are reachable

- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.







- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.







- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.







- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.







SAT - new states are reachable

simple path
- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.







- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.





- Inductive step: proved incrementally following the partial order of the MSC.
- Base case: bounded feasibility check.







• Base case: bounded feasibility check.







• Base case: bounded feasibility check.





#### Outline

#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

Proving the unfeasibility of scenarios

- 3 Explanations of Unfeasibility
- 4 Experimental Evaluation
- 5 Conclusions and future work

- Typical use case:
  - We expect that a scenario is feasible.
  - The analysis proves that the scenario is unfeasible in the network.
  - How do we explain the unfeasibility?
- We extract three types of explanations for the unfeasibility.

#### Unfeasibility due to a component

Explained with a formula that:

- Is required by the component when simulating its MSC events.
- Is not consistent with the other components when they simulate the events of the MSC.



Unfeasibility and Explanations of MSC

#### Unfeasibility due to a component

Explained with a formula that:

- Is required by the component when simulating its MSC events.
- Is not consistent with the other components when they simulate the events of the MSC.
- It is the interpolant of *A* and *B*:
  - A is the encoding of the component and its MSC events.
  - B is the encoding of the other components and their MSC events.



#### Unfeasibility due the network

Explained with a formula that:

- Is required by the network when simulating the MSC.
- Is not consistent with the additional constraints of the MSC.



#### Unfeasibility due the network

Explained with a formula that:

- Is required by the network when simulating the MSC.
- Is not consistent with the additional constraints of the MSC.

It is the interpolant of *A* and *B*:

- A is the encoding of the network and the MSC.
- B are the CMSC constraints.



Subset of the original CMSC that is still unfeasible with the network.



Subset of the original CMSC that is still unfeasible with the network. Extracted from the unsatisfiable core of the encoding.



#### Outline

#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

Proving the unfeasibility of scenarios

- 3 Explanations of Unfeasibility
- Experimental Evaluation
- 5 Conclusions and future work

#### Implementation:

- Approach implemented on top of the NUSMV model checker.
- We use the MATHSAT SMT solver.

#### Settings:

- Linear hybrid automata benchmarks.
- Several handcrafted (unsatisfiable) MSCs.
- We scaled the dimension of the benchmarks (number of automata, length of the MSCs).

#### Comparison:

- MSC partitioned k-induction.
- Monolithic k-induction on the system composed with the monitor automata.

# Partitioned k-induction vs. Monolithic k-induction (run times)



#### Outline

#### Background

- SMT analysis of Hybrid Systems
- Scenario-Verification

2 Proving the unfeasibility of scenarios

- 3 Explanations of Unfeasibility
- 4 Experimental Evaluation
- 5 Conclusions and future work

- Efficient approach for proving the unfeasibility of CMSC.
  - The encoding exploits the structure of the CMSC.
  - Partitioned k-induction.
- Unfeasibility explanations:
  - Useful to localize and correct the errors.
  - Extracted exploiting the SMT solver functionalities.

Future works:

- More expressive MSCs (e.g. partial MSCs specifications).
- Validate the extracted explanations by real users.
- Automatic refinement loop in the abstraction.
- Non-linear hybrid systems.

## Thank you for your attention.