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- In C and C++ perennial source of security vulnerabilities
- In Java, C# program crashes due to exceptions



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  - + Can guarantee absence of errors
  - Can be hard to understand and fix detected errors
  - Does not help programmer write safe-by-construction code
- **Dynamic approaches** to memory safety
  - + Widely used in managed languages
  - But only transforms vulnerability into program crash
  - Run-time overhead

# Key Idea: Use Program Synthesis

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- 1 Programmer specifies **which** parts of the program should be guarded

Example:

```
if(???) {R} else { /* handle error */ }
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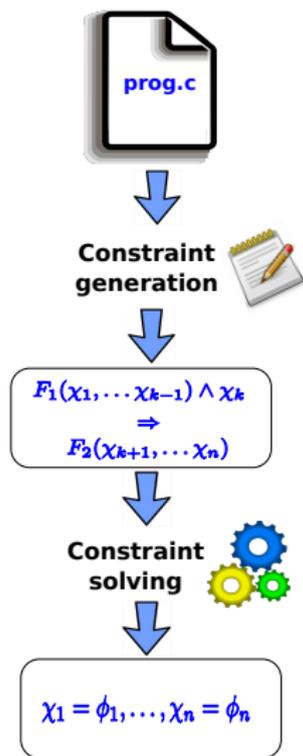
# Key Idea: Use Program Synthesis

## Key Idea: Program synthesis to guarantee memory safety

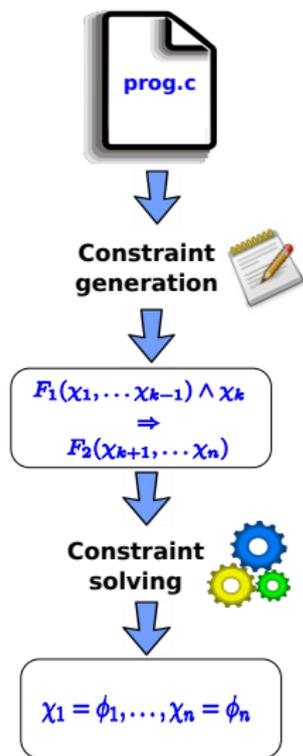
- 1 Programmer specifies **which** parts of the program should be guarded
- 2 Our technique **synthesizes** correct and optimal guards that guarantee memory safety
  - Optimal means as weak and as simple as possible

Example:

```
if(???) {R} else { /* handle error */ }
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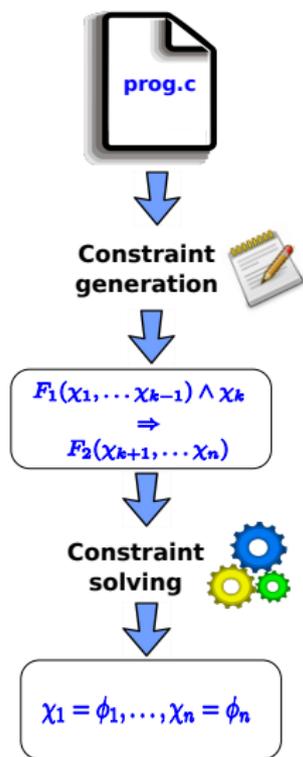


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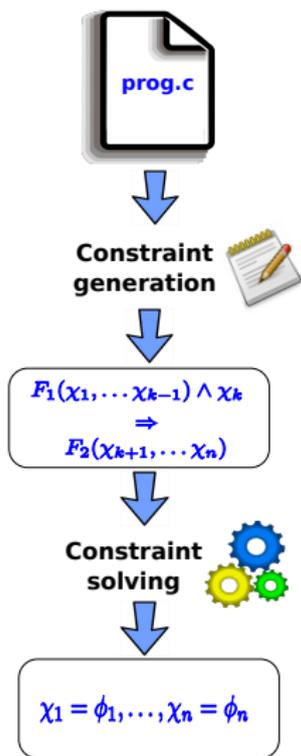
- Represent unknown guards using placeholders



## 1 Constraint Generation:

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- Perform dual forward and backward analysis to generate constraint for each unknown

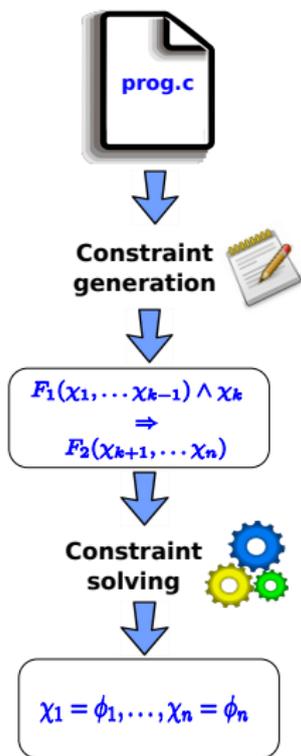
# Solution Overview



## 1 Constraint Generation:

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## 2 Constraint Solving:

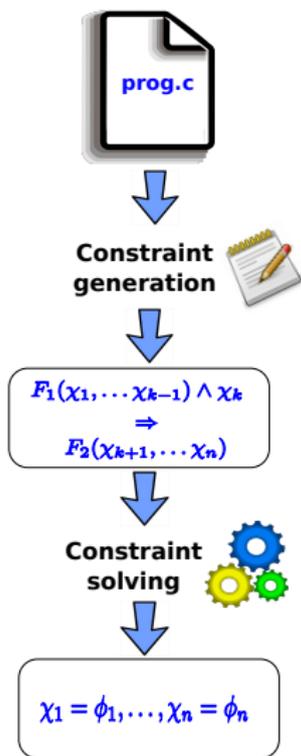


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## 2 Constraint Solving:

- An extended abduction algorithm for solving constraint system with multiple unknowns
- Guarantees Pareto-optimality

# Constraint Generation Overview

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if(??)
{
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$\psi$

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# Constraint Generation Overview

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- At synthesis point, compute postcondition  $\phi$  of code above ??
- Compute precondition  $\psi$  that ensures **memory safety** of code guarded by ??
- Condition to guarantee memory safety:

$$\phi \wedge ?? \models \psi$$

# Key Insight: Abduction

**Solution:** Abductive inference





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- Given facts  $F$  and desired outcome  $O$ , find simple explanatory hypothesis  $E$  such that

$$F \wedge E \models O \text{ and } \text{SAT}(F \wedge E)$$



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- $O \equiv$  memory safety precondition  $\psi$
- $E \equiv$  Solution for ??



$$F_1(\vec{\chi}_1) \wedge \chi_1 \Rightarrow F'_1(\vec{\chi}_1)$$

$$F_2(\vec{\chi}_2) \wedge \chi_2 \Rightarrow F'_2(\vec{\chi}_2)$$

⋮

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- Cannot directly use abduction because constraints have multiple unknowns
- New iterative, stratification-based algorithm for solving constraint system
- Uses abduction as a helper procedure
- Resulting solution is Pareto-optimal
  - Cannot improve solution for one unknown without making others worse

## Example

- Code snippet from Unix Coreutils with protected memory access

```
int main(int argc,
        char** argv)
{
    if(argc<=1) return -1;
    argv++; argc--;

    optind=0;
    while(...) {
        optind++;
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    }
    if(??) {
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## Example

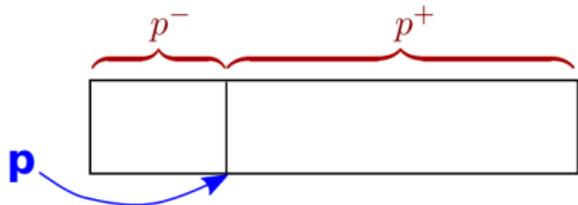
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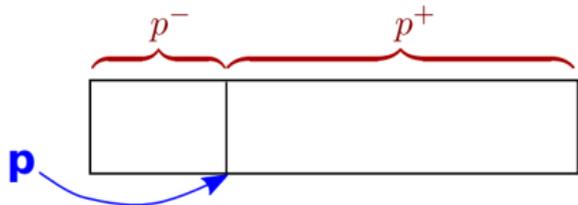


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# Example

- Code snippet from Unix Coreutils with protected memory access
- **Convention:** For pointer  $p$ :
  - $p^+$  represents distance to end of memory block
  - $p^-$  represents distance from beginning of memory block



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int main(int argc,
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## Example Cont.

- **First Step:** Compute what is known at ??  $\Rightarrow$  **postcondition**  $\phi$

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$$argv^+ = argc \wedge argv^- = 0$$

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- From language semantics:

$$argv^+ = argc \wedge argv^- = 0$$

- From computing the strongest postcondition:

$$argv^+ = argc \wedge \\ argv^- \geq 1 \wedge optind \geq 0$$

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- **Second Step:** Compute what needs to hold at ?? to ensure memory safety  
⇒ precondition  $\psi$

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## Example Cont.

- **Second Step:** Compute what needs to hold at ?? to ensure memory safety  
⇒ precondition  $\psi$
- Buffer access:

$$\begin{aligned}optind + 1 &< argv^+ \wedge \\optind + 1 &\geq -argv^-\end{aligned}$$

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## Example Cont.

- Solve abduction problem

$\phi \wedge ?? \models \psi$  where

$$\phi : \quad \begin{array}{l} \text{argv}^+ = \text{argc} \wedge \\ \text{argv}^- \geq 1 \wedge \text{optind} \geq 0 \end{array}$$

$$\psi : \quad \begin{array}{l} \text{optind} + 1 < \text{argv}^+ \wedge \\ \text{optind} + 1 \geq -\text{argv}^- \end{array}$$

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- **Solution:**  $argc - optind > 1$

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A word cloud of various Unix/Linux commands. The most prominent words are 'mkdir' in dark red, 'cp' in green, and 'ls' in dark red. Other visible words include 'rm', 'ln', 'mv', 'dd', 'df', 'truncate', 'dir', 'chmod', 'chown', and 'install'.

- Evaluated technique on the Unix Coreutils and parts of OpenSSH
- Removed conditionals used to prevent memory safety errors
- Used our new technique to synthesize the missing guards



## Experiments Cont.

Program	Lines	# holes	Time (s)	Memory	Synthesis successful?	Bug?
Coreutils hostname	160	1	0.15	10 MB	Yes	No
Coreutils tee	223	1	0.84	10 MB	Yes	Yes
Coreutils runcon	265	2	0.81	12 MB	Yes	No
Coreutils chroot	279	2	0.53	23 MB	Yes	No
Coreutils remove	710	2	1.38	66MB	Yes	No
Coreutils nl	758	3	2.07	80 MB	Yes	No
SSH - sshconnect	810	3	1.43	81 MB	Yes	No
Coreutils mv	929	4	2.03	42 MB	Yes	No
SSH - do_authentication	1,904	4	3.92	86 MB	Yes	Yes
SSH - ssh_session	2,260	5	4.35	81 MB	Yes	No

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# Summary

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- Two key ingredients:
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  - **Constraint solving:** New abduction-based algorithm for finding optimal solutions for placeholders representing unknown guards
- Experimental validation of our approach

# Questions?

