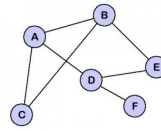


# CS311H: Discrete Mathematics

## Introduction to Graph Theory

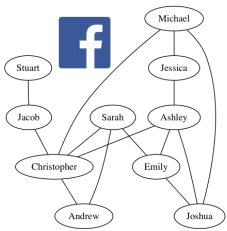
Instructor: Işıl Dillig

### Motivation



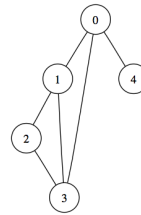
- ▶ Graph is a fundamental mathematical structure in computer science
- ▶ Graph  $G = (V, E)$  consists of a set of **vertices (nodes)**  $V$  and **edges**  $E$  between these nodes
- ▶ Lots of applications in many areas: web search, networking, databases, AI ...

### Example: Social Network as a Graph



- ▶ Nodes represent users (Michael, Jessica, Stuart ...)
- ▶ Edges represent friendship (e.g., Michael is friends with Jessica)
- ▶ Edge between nodes  $u$  and  $v$  is written as  $(u, v)$
- ▶ e.g.,  $(Sarah, Andrew)$  is an edge in this graph.

### Terminology



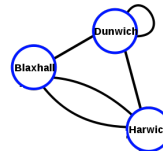
- ▶ Two nodes  $u$  and  $v$  are **adjacent** if there exists an edge between them (e.g., nodes 1 and 3)
- ▶ An edge  $(u, v)$  is **incident with** nodes  $u$  and  $v$
- ▶ **Degree** of a vertex  $v$ , written  $\deg(v)$ , is the number of edges incident with it
- ▶ **Neighborhood** of a vertex is the set of vertices adjacent to it

### Question

Consider a graph  $G$  with vertices  $v_1, v_2, v_3, v_4$  and edges  $(v_1, v_2), (v_2, v_3), (v_1, v_3), (v_2, v_4)$ .

1. Draw this graph.
2. What is the degree of each vertex?

### Simple Graphs



- ▶ Graph contains a **loop** if any node is adjacent to itself
- ▶ A **simple graph** does not contain loops and there exists at most one edge between any pair of vertices
- ▶ Graphs that have multiple edges connecting two vertices are called **multi-graphs**
- ▶ Most graphs we will look at are simple graphs

## Question

Consider a simple graph  $G$  where two vertices  $A$  and  $B$  have the same neighborhood.

Which of the following statements **must** be true about  $G$ ?

- A. The degree of each vertex must be even.
- B. Both  $A$  and  $B$  have a degree of 0.
- C. There cannot be an edge between  $A$  and  $B$ .

## Handshaking Theorem

Let  $G = (V, E)$  be a graph with  $m$  edges. Then:

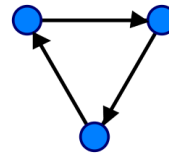
$$\sum_{v \in V} \deg(v) = 2m$$

- ▶ **Intuition:** Each edge contributes two to the sum of the degrees
- ▶ **Proof:**
- ▶
- ▶

## Applications of Handshaking Theorem

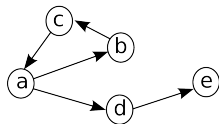
- ▶ Is it possible to construct a graph with 5 vertices where each vertex has degree 3?
- ▶ Prove that every graph has an even number of vertices of odd degree.
- ▶ If  $n$  people go to a party and everyone shakes everyone else's hand, how many handshakes occur?

## Directed Graphs



- ▶ All graphs we considered so far are **undirected**
- ▶ In undirected graphs, edge  $(u, v)$  same as  $(v, u)$
- ▶ A **directed edge (arc)** is an ordered pair  $(u, v)$  (i.e.,  $(u, v)$  not same as  $(v, u)$ )
- ▶ A **directed graph** is a graph with directed edges

## In-Degree and Out-Degree of Directed Graphs

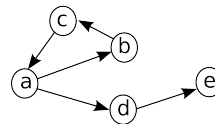


- ▶ The **in-degree** of a vertex  $v$ , written  $\deg^-(v)$ , is the number of edges going into  $v$
- ▶  $\deg^-(a) =$
- ▶ The **out-degree** of a vertex  $v$ , written  $\deg^+(v)$ , is the number of edges leaving  $v$
- ▶  $\deg^+(a) =$

## Handshaking Theorem for Directed Graphs

Let  $G = (V, E)$  be a directed graph. Then:

$$\sum_{v \in V} \deg^-(v) = \sum_{v \in V} \deg^+(v) = |E|$$

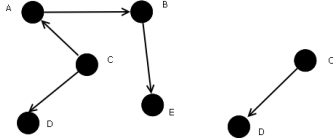


- ▶  $\sum_{v \in V} \deg^-(v) =$
- ▶  $\sum_{v \in V} \deg^+(v) =$

## Subgraphs

▶ A graph  $G = (V, E)$  is a **subgraph** of another graph  $G' = (V', E')$  if  $V \subseteq V'$  and  $E \subseteq E'$

▶ Example:



▶ Graph  $G$  is a **proper subgraph** of  $G'$  if  $G \neq G'$ .

## Question

Consider a graph  $G$  with vertices  $\{v_1, v_2, v_3, v_4\}$  and edges  $(v_1, v_3), (v_1, v_4), (v_2, v_3)$ .

Which of the following are subgraphs of  $G$ ?

1. Graph  $G_1$  with vertex  $v_1$  and edge  $(v_1, v_3)$
2. Graph  $G_2$  with vertices  $\{v_1, v_3\}$  and no edges
3. Graph  $G_3$  with vertices  $\{v_1, v_2\}$  and edge  $(v_1, v_2)$

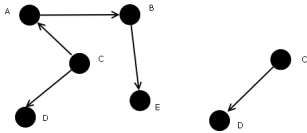
## Induced Subgraph

▶ Consider a graph  $G = (V, E)$  and a set of vertices  $V'$  such that  $V' \subseteq V$

▶ Graph  $G'$  is the **induced subgraph** of  $G$  with respect to  $V'$  if:

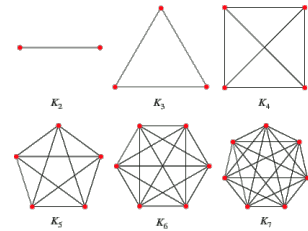
1.  $G'$  contains exactly those vertices in  $V'$
2. For all  $u, v \in V'$ , edge  $(u, v) \in G'$  iff  $(u, v) \in G$

▶ Subgraph induced by vertices  $\{C, D\}$ :



## Complete Graphs

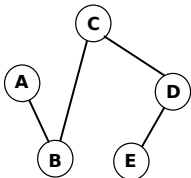
▶ A **complete graph** is a simple undirected graph in which every pair of vertices is connected by one edge.



▶ How many edges does a complete graph with  $n$  vertices have?

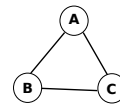
## Bipartite graphs

▶ A simple undirected graph  $G = (V, E)$  is called **bipartite** if  $V$  can be partitioned into two disjoint sets  $V_1$  and  $V_2$  such that every edge in  $E$  connects a  $V_1$  vertex to a  $V_2$  vertex

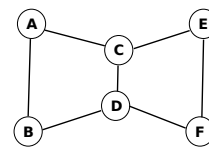


## Examples Bipartite and Non-Bi-partite Graphs

▶ Is this graph bipartite?



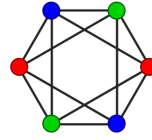
▶ What about this graph?



## Questions about Bipartite Graphs

- ▶ Does there exist a complete graph that is also bipartite?
- ▶ Consider a graph  $G$  with 5 nodes and 7 edges. Can  $G$  be bipartite?

## Graph Coloring



- ▶ A **coloring** of a graph is the assignment of a color to each vertex so that no two adjacent vertices are assigned the same color.
- ▶ A graph is  **$k$ -colorable** if it is possible to color it using  $k$  colors.
  - ▶ e.g., graph on left is 3-colorable
  - ▶ Is it also 2-colorable?
- ▶ The **chromatic number** of a graph is the least number of colors needed to color it.
  - ▶ What is the chromatic number of this graph?

## Question

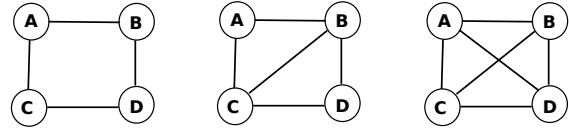
Consider a graph  $G$  with vertices  $\{v_1, v_2, v_3, v_4\}$  and edges  $(v_1, v_2), (v_1, v_3), (v_2, v_3), (v_2, v_4)$ .

Which of the following are valid colorings for  $G$ ?

1.  $v_1 = \text{red}, v_2 = \text{green}, v_3 = \text{blue}$
2.  $v_1 = \text{red}, v_2 = \text{green}, v_3 = \text{blue}, v_4 = \text{red}$
3.  $v_1 = \text{red}, v_2 = \text{green}, v_3 = \text{red}, v_4 = \text{blue}$

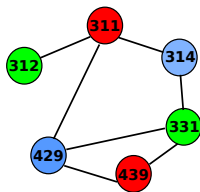
## Examples

What are the chromatic numbers for these graphs?



## Applications of Graph Coloring

- ▶ Graph coloring has lots of applications, particularly in scheduling.
- ▶ **Example:** What's the minimum number of time slots needed so that no student is enrolled in conflicting classes?



## A Scheduling Problem

- ▶ The math department has 6 committees  $C_1, \dots, C_n$  that meet once a month.
- ▶ The committee members are:
 

$C_1 = \{\text{Allen, Brooks, Marg}\}$	$C_2 = \{\text{Brooks, Jones, Morton}\}$
$C_3 = \{\text{Allen, Marg, Morton}\}$	$C_4 = \{\text{Jones, Marg, Morton}\}$
$C_5 = \{\text{Allen, Brooks}\}$	$C_6 = \{\text{Brooks, Marg, Morton}\}$
- ▶ How many different meeting times must be used to guarantee that no one has conflicting meetings?

## Bipartite Graphs and Colorability

Prove that a graph  $G = (V, E)$  is **bipartite** if and only if it is **2-colorable**.

## Complete graphs and Colorability

Prove that any complete graph  $K_n$  has chromatic number  $n$ .

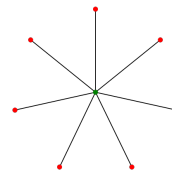
## Degree and Colorability

**Theorem:** A simple graph  $G$  is always  $\max\_degree(G) + 1$ -colorable.. Then,  $G$  is  $n + 1$ -colorable.

## Degree and Colorability, cont.

## Degree and Colorability, cont.

## Star Graphs and Colorability



- ▶ A **star graph**  $S_n$  is a graph with one vertex  $u$  at the center and the only edges are from  $u$  to each of  $v_1, \dots, v_{n-1}$ .
- ▶ Draw  $S_2, S_3, S_4, S_5$ .
- ▶ What is the chromatic number of  $S_n$ ?

## Question About Star Graphs

Suppose we have two star graphs  $S_k$  and  $S_m$ . Now, pick a random vertex from each graph and connect them with an edge.

Which of the following statements must be true about the resulting graph  $G$ ?

1. The chromatic number of  $G$  is 3
2.  $G$  is 2-colorable.
3.  $\max\_degree(G) = \max(k, m)$ .