

CS311H: Discrete Mathematics

Mathematical Induction

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Introduction to Mathematical Induction

- ▶ Many mathematical theorems assert that a property holds for **all** natural numbers, odd positive integers, etc.
- ▶ **Mathematical induction**: very important proof technique for proving such universally quantified statements
- ▶ Induction will come up over and over again in other classes:
 - ▶ algorithms, programming languages, automata theory, ...

Analogy



- ▶ Suppose we have an infinite ladder, and we know two things:
 1. We can reach the first rung of the ladder
 2. If we reach a particular rung, then we can also reach the next rung
- ▶ From these two facts, can we conclude we can reach **every** step of the infinite ladder?
- ▶ Answer is **yes**, and mathematical induction allows us to make arguments like this

Mathematical Induction

- ▶ Used to prove statements of the form $\forall x \in \mathbb{Z}^+. P(x)$
- ▶ An inductive proof has two steps:
 1. **Base case:** Prove that $P(1)$ is true
 2. **Inductive step:** Prove $\forall n \in \mathbb{Z}^+. P(n) \rightarrow P(n+1)$
- ▶ Induction says if you can prove (1) and (2), you can conclude:

$$\forall x \in \mathbb{Z}^+. P(x)$$

Inductive Hypothesis

- ▶ In the **inductive step**, need to show:

$$\forall n \in \mathbb{Z}^+. P(n) \rightarrow P(n + 1)$$

- ▶ To prove this, we assume $P(n)$ holds, and based on this assumption, prove $P(n + 1)$
- ▶ The assumption that $P(n)$ holds is called the **inductive hypothesis**

Example 1

- Prove the following statement by induction:

$$\forall n \in \mathbb{Z}^+. \sum_{i=1}^n i = \frac{(n)(n+1)}{2}$$

- **Base case:** $n = 1$. In this case, $\sum_{i=1}^1 i = 1$ and $\frac{(1)(1+1)}{2} = 1$; thus, the base case holds.
- **Inductive step:** By the inductive hypothesis, we assume $P(k)$:

$$\sum_{i=1}^k i = \frac{(k)(k+1)}{2}$$

- Now, we want to show $P(k+1)$:

$$\sum_{i=1}^{k+1} i = \frac{(k+1)(k+2)}{2}$$

Example 1, cont.

- First, observe:

$$\sum_{i=1}^{k+1} i = \sum_{i=1}^k i + (k+1)$$

- By the inductive hypothesis, $\sum_{i=1}^k i = \frac{(k)(k+1)}{2}$; thus:

$$\sum_{i=1}^{k+1} i = \frac{(k)(k+1)}{2} + (k+1)$$

- Rewrite left hand side as:

$$\sum_{i=1}^{k+1} i = \frac{k^2 + 3k + 2}{2} = \frac{(k+1)(k+2)}{2}$$

- Since we proved both base case and inductive step, property holds.

Example 2

- ▶ Prove the following statement for **all non-negative integers** n :

$$\sum_{i=0}^n 2^i = 2^{n+1} - 1$$

- ▶ Left as “do-it-at-home exercise” with solution!
- ▶ Since need to show for all $n \geq 0$, base case is $P(0)$, not $P(1)$!
- ▶ **Base case ($n = 0$):** $2^0 = 1 = 2^1 - 1$
- ▶ **Inductive step:**

$$\sum_{i=0}^{k+1} 2^i = \sum_{i=0}^k 2^i + 2^{k+1}$$

Example 2, cont.

$$\sum_{i=0}^{k+1} 2^i = \sum_{i=0}^k 2^i + 2^{k+1}$$

- By the **inductive hypothesis**, we have:

$$\sum_{i=0}^k 2^i = 2^{k+1} - 1$$

- Therefore:

$$\sum_{i=0}^{k+1} 2^i = 2^{k+1} - 1 + 2^{k+1}$$

- Rewrite as:

$$\sum_{i=0}^{k+1} 2^i = 2 \cdot 2^{k+1} - 1 = 2^{k+2} - 1$$

Example 3

► Prove that $2^n < n!$ for all integers $n \geq 4$



Example 4

► Prove that $3 \mid (n^3 - n)$ for all positive integers n .



The Horse Paradox

- ▶ Easy to make subtle errors when trying to prove things by induction – pay attention to details!
- ▶ Consider the statement: All horses have the same color
- ▶ What is wrong with the following bogus proof of this statement?
- ▶ $P(n)$: A collection of n horses have the same color
- ▶ Base case: $P(1)$ ✓

Bogus Proof, cont.

- ▶ Induction: Assume $P(k)$; prove $P(k + 1)$
- ▶ Consider a collection of $k + 1$ horses: $h_1, h_2, \dots, h_{k+1}$
- ▶ By IH, $h_1, h_2, \dots, h_k$ have the same color; let this color be c
- ▶ By IH, $h_2, \dots, h_{k+1}$ have same color; call this color c'
- ▶ Since h_2 has color c and c' , we have $c = c'$
- ▶ Thus, $h_1, h_2, \dots, h_{k+1}$ also have same color
- ▶ What's the fallacy?

Strengthening the Inductive Hypothesis

- ▶ Suppose we want to prove $\forall x \in \mathbb{Z}^+. P(x)$, but proof doesn't go through
- ▶ **Common trick:** Prove a stronger property $Q(x)$
- ▶ If $\forall x \in \mathbb{Z}^+. Q(x) \rightarrow P(x)$ and $\forall x \in \mathbb{Z}^+. Q(x)$ is provable, this implies $\forall x \in \mathbb{Z}^+. P(x)$
- ▶ In many situations, strengthening inductive hypothesis allows proof to go through!

Example

- ▶ Prove the following theorem: “For all $n \geq 1$, the sum of the first n odd numbers is a perfect square.”
- ▶ We want to prove $\forall x \in \mathbb{Z}^+. P(x)$ where:

$$P(n) = \sum_{i=1}^n 2i - 1 = k^2 \text{ for some integer } k$$

- ▶ Try to prove this using induction...

Example, cont.

- ▶ Let's use a stronger predicate:

$$Q(n) = \sum_{i=1}^n 2i - 1 = n^2$$

- ▶ Clearly $Q(n) \rightarrow P(n)$
- ▶ Now, prove $\forall n \in \mathbb{Z}^+. Q(n)$ using induction!

Strong Induction

- ▶ Slight variation on the inductive proof technique is **strong induction**
- ▶ Regular and strong induction only differ in the inductive step
- ▶ **Regular induction:** assume $P(k)$ holds and prove $P(k + 1)$
- ▶ **Strong induction:** assume $P(1), P(2), \dots, P(k)$; prove $P(k + 1)$
- ▶ Strong induction can be viewed as standard induction with **strengthened inductive hypothesis!**

Motivation for Strong Induction

- ▶ Prove that if n is an integer greater than 1, then it is either a prime or can be written as the product of primes.
- ▶ Let's first try to prove the property using regular induction.
- ▶ **Base case ($n=2$):** Since 2 is a prime number, $P(2)$ holds.
- ▶ **Inductive step:** Assume k is either a prime or the product of primes.
- ▶ But this doesn't really help us prove the property about $k + 1$!
- ▶ Claim is proven much easier using strong induction!

Proof Using Strong Induction

Prove that if n is an integer greater than 1, then it is either a prime or can be written as the product of primes.

- ▶ **Base case:** same as before.
- ▶ **Inductive step:** Assume each of $2, 3, \dots, k$ is either prime or product of primes.
- ▶ Now, we want to prove the same thing about $k + 1$
- ▶ Two cases: k is either (i) prime or (ii) composite
- ▶ If it is prime, property holds.

Proof, cont.

- ▶ If composite, $k + 1$ can be written as pq where $2 \leq p, q \leq k$
- ▶ By the IH, p, q are either primes or product of primes.
- ▶ Thus, $k + 1$ can also be written as product of primes
- ▶ **Observe:** Much easier to prove this property using strong induction!

A Word about Base Cases

- ▶ In all examples so far, we had only one base case
 - ▶ i.e., only proved the base case for one integer
- ▶ In some inductive proofs, there may be **multiple base cases**
 - ▶ i.e., prove base case for the first k numbers
- ▶ In the latter case, inductive step only needs to consider numbers greater than k

Example

- ▶ Prove that every integer $n \geq 12$ can be written as $n = 4a + 5b$ for some non-negative integers a, b .
- ▶ Proof by **strong induction** on n and consider 4 base cases
- ▶ **Base case 1 ($n=12$):** $12 = 3 \cdot 4 + 0 \cdot 5$
- ▶ **Base case 2 ($n=13$):** $13 = 2 \cdot 4 + 1 \cdot 5$
- ▶ **Base case 3 ($n=14$):** $14 = 1 \cdot 4 + 2 \cdot 5$
- ▶ **Base case 4 ($n=15$):** $15 = 0 \cdot 4 + 3 \cdot 5$

Example, cont.

Prove that every integer $n \geq 12$ can be written as $n = 4a + 5b$ for some non-negative integers a, b .

- ▶ **Inductive hypothesis:** Suppose every $12 \leq i \leq k$ can be written as $i = 4a + 5b$.
- ▶ **Inductive step:** We want to show $k + 1$ can also be written this way for $k + 1 \geq 16$
- ▶ **Observe:** $k + 1 = (k - 3) + 4$
- ▶ By the inductive hypothesis, $k - 3 = 4a + 5b$ for some a, b because $k - 3 \geq 12$
- ▶ But then, $k + 1$ can be written as $4(a + 1) + 5b$

Another Example

- ▶ For $n \geq 1$, prove there exist natural numbers a, b such that:

$$5^n = a^2 + b^2$$

- ▶ **Insight:** $5^{k+1} = 5^2 \cdot 5^{k-1}$

Matchstick Example

- ▶ **The Matchstick game:** There are two piles with same number of matches initially
- ▶ Two players take turns removing any positive number of matches from one of the two piles
- ▶ Player who removes the last match wins the game
- ▶ **Prove:** Second player always has a winning strategy.

Matchstick Proof

- ▶ $P(n)$: Player 2 has winning strategy if initially n matches in each pile
- ▶ Base case:
- ▶ Induction: Assume $\forall j. 1 \leq j \leq k \rightarrow P(j)$; show $P(k+1)$
- ▶ Inductive hypothesis:
- ▶ Prove Player 2 wins if each pile contains $k+1$ matches

Matchstick Proof, cont.

- ▶ Case 1: Player 1 takes $k + 1$ matches from one of the piles.
- ▶ What is winning strategy for player 2
- ▶ Case 2: Player 1 takes r matches from one pile, where $1 \leq r \leq k$
- ▶ Now, player 2 takes r matches from other pile
- ▶ Now, the inductive hypothesis applies $\Rightarrow$ player 2 has winning strategy for rest of the game