

CS311H: Discrete Mathematics

Introduction to Number Theory

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Introduction to Number Theory

- ▶ Number theory is the branch of mathematics that deals with integers and their properties
- ▶ Number theory has a number of applications in computer science, esp. in modern **cryptography**
- ▶ **Next few lectures:** Basic concepts in number theory and its application in crypto

Divisibility

- ▶ Given two integers a and b where $a \neq 0$, we say a divides b if there is an integer c such that $b = ac$
- ▶ If a divides b , we write $a|b$; otherwise, $a \nmid b$
- ▶ Example: $2|6$, $2 \nmid 9$
- ▶ If $a|b$, a is called a factor of b
- ▶ b is called a multiple of a

Example

- ▶ **Question:** If n and d are positive integers, how many positive integers not exceeding n are divisible by d ?
- ▶ **Recall:** All positive integers divisible by d are of the form dk
- ▶ We want to find how many numbers dk there are such that $0 < dk \leq n$.
- ▶ In other words, we want to know how many **integers** k there are such that $0 < k \leq \frac{n}{d}$
- ▶ How many integers are there between 1 and $\frac{n}{d}$?

Properties of Divisibility

► **Theorem 1:** If $a|b$ and $b|c$, then $a|c$



Divisibility Properties, cont.

- ▶ **Theorem 2:** If $a|b$ and $a|c$, then $a|(mb + nc)$ for any int m, n
- ▶ **Proof:**
- ▶ **Corollary 1:** If $a|b$ and $a|c$, then $a|(b + c)$ for any int c
- ▶ **Corollary 2:** If $a|b$, then $a|mb$ for any int m

The Division Theorem

- ▶ **Division theorem:** Let a be an integer, and d a positive integer. Then, there are **unique** integers q, r with $0 \leq r < d$ such that $a = dq + r$
- ▶ Here, d is called **divisor**, and a is called **dividend**
- ▶ q is the **quotient**, and r is the **remainder**.
- ▶ We use the $r = a \bmod d$ notation to express the remainder
- ▶ The notation $q = a \operatorname{div} d$ expresses the quotient
- ▶ What is $101 \bmod 11$?
- ▶ What is $101 \operatorname{div} 11$?

Congruence Modulo

- ▶ In number theory, we often care if two integers a, b have same remainder when divided by m .
- ▶ If so, a and b are **congruent modulo** m , $a \equiv b \pmod{m}$.
- ▶ More technically, if a and b are integers and m a positive integer, $a \equiv b \pmod{m}$ iff $m \mid (a - b)$
- ▶ **Example:** 7 and 13 are congruent modulo 3.
- ▶ **Example:** Find a number congruent to 7 modulo 4.

Congruence Modulo Theorem

- ▶ **Theorem:** $a \equiv b \pmod{m}$ iff $a \bmod m = b \bmod m$
- ▶ **Part 1, $\Rightarrow$:** Suppose $a \equiv b \pmod{m}$.
- ▶ Then, by definition of $\equiv$, $m \mid (a - b)$
- ▶ By definition of $\mid$, there exists k such that $a - b = mk$, i.e.,
 $a = b + mk$
- ▶ By division thm, $b = mp + r$ for some $0 \leq r < m$
- ▶ Then, $a = mp + r + mk = m(p + k) + r$
- ▶ Thus, $a \bmod m = r = b \bmod m$

Congruence Modulo Theorem Proof, cont.

- ▶ **Theorem:** $a \equiv b \pmod{m}$ iff $a \bmod m = b \bmod m$
- ▶ Part 2, $\Leftarrow$: Suppose $a \bmod m = b \bmod m$
- ▶ Then, there exists some p_1, p_2, r such that $a = p_1 \cdot m + r$ and $b = p_2 \cdot m + r$ where $0 \leq r < m$
- ▶ Then, $a - b = p_1 \cdot m + r - p_2 \cdot m - r = m \cdot (p_1 - p_2)$
- ▶ Thus, $m \mid (a - b)$
- ▶ By definition of $\equiv$, $a \equiv b \pmod{m}$

Example

- ▶ Prove that if $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then:

$$a + c \equiv b + d \pmod{m}$$



Applications of Congruence in Cryptography

- ▶ Congruences have many applications in cryptography, e.g., shift ciphers
- ▶ **Shift cipher** with **key k** encrypts message by shifting each letter by k letters in alphabet (if past Z , then wrap around)
- ▶ What is encryption of "KILL HIM" with shift cipher of key 3?
- ▶ Shift ciphers also called **Cesar ciphers** because Julius Caesar encrypted secret messages to his generals this way

Mathematical Encoding of Shift Ciphers

- ▶ First, let's number letters A-Z with 0 – 25
- ▶ Represent message with sequence of numbers
- ▶ **Example:** The sequence "25 0 2" represents "ZAC"
- ▶ To encrypt, apply **encryption function** f defined as:

$$f(x) = (x + k) \bmod 26$$

- ▶ Because f is bijective, its inverse yields decryption function:

$$g(x) = (x - k) \bmod 26$$

Ciphers and Congruence Modulo

- ▶ Shift cipher is a very primitive and insecure cipher because very easy to infer what k is
- ▶ But contains some useful ideas:
 - ▶ Encoding words as sequence of numbers
 - ▶ Use of modulo operator
- ▶ Modern encryption schemes much more sophisticated, but also share these principles (coming lectures)

Prime Numbers

- ▶ A positive integer p that is greater than 1 and divisible only by 1 and itself is called a **prime number**.
- ▶ First few primes: 2, 3, 5, 7, 11, ...
- ▶ A positive integer that is greater than 1 and that is not prime is called a **composite number**
- ▶ Example: 4, 6, 8, 9, ...

Fundamental Theorem of Arithmetic

- ▶ **Fundamental Thm:** Every positive integer greater than 1 is either prime or can be written **uniquely** as a product of primes.
- ▶ This unique product of prime numbers for x is called the **prime factorization** of x
- ▶ Examples:
 - ▶ $12 =$
 - ▶ $21 =$
 - ▶ $99 =$

Determining Prime-ness

- ▶ In many applications, such as crypto, important to determine if a number is prime – following thm is useful for this:
- ▶ **Theorem:** If n is composite, then it has a prime divisor less than or equal to $\sqrt{n}$
- ▶
- ▶
- ▶
- ▶

Consequence of This Theorem

Theorem: If n is composite, then it has a prime divisor $\leq \sqrt{n}$

- ▶ Thus, to determine if n is prime, only need to check if it is divisible by primes $\leq \sqrt{n}$
- ▶ **Example:** Show that 101 is prime
- ▶ Since $\sqrt{101} < 11$, only need to check if it is divisible by 2, 3, 5, 7.
- ▶ Since it is not divisible by any of these, we know it is prime.

Infinitely Many Primes

- ▶ **Theorem:** There are infinitely many prime numbers.
- ▶ **Proof:** (by contradiction) Suppose there are finitely many primes: $p_1, p_2, \dots, p_n$
- ▶ Now consider the number $Q = p_1 p_2 \dots p_n + 1$. Q is either prime or composite
- ▶ **Case 1:** Q is prime. We get a contradiction, because we assumed only prime numbers are $p_1, \dots, p_n$
- ▶ **Case 2:** Q is composite. In this case, Q can be written as product of primes.
- ▶ But Q is not divisible by any of $p_1, p_2, \dots, p_n$
- ▶ Hence, by Fundamental Thm, not composite $\Rightarrow \perp$



Greatest Common Divisors

- ▶ Suppose a and b are integers, not both 0.
- ▶ Then, the largest integer d such that $d|a$ and $d|b$ is called **greatest common divisor** of a and b , written $\gcd(a,b)$.
- ▶ Example: $\gcd(24, 36) =$
- ▶ Example: $\gcd(2^3 5, 2^2 3) =$
- ▶ Example: $\gcd(14, 25) =$
- ▶ Two numbers whose gcd is 1 are called **relatively prime**.
- ▶ Example: 14 and 25 are relatively prime

Least Common Multiple

- ▶ The **least common multiple** of two positive integers a and b , written $\text{lcm}(a,b)$, is the smallest integer c such that $a|c$ and $b|c$.
- ▶ Example: $\text{lcm}(9, 12)=$
- ▶ Example: $\text{lcm}(2^3 3^5 7^2, 2^4 3^3)=$

Theorem about LCM and GCD

- ▶ **Theorem:** Let a and b be positive integers. Then,
 $ab = \gcd(a, b) \cdot \text{lcm}(a, b)$
- ▶ **Proof:** Let $a = p_1^{i_1} p_2^{i_2} \dots p_n^{i_n}$ and $b = p_1^{j_1} p_2^{j_2} \dots p_n^{j_n}$
- ▶ Then, $ab = p_1^{i_1+j_1} p_2^{i_2+j_2} \dots p_n^{i_n+j_n}$
- ▶ $\gcd(a, b) = p_1^{\min(i_1, j_1)} p_2^{\min(i_2, j_2)} \dots p_n^{\min(i_n, j_n)}$
- ▶ $\text{lcm}(a, b) = p_1^{\max(i_1, j_1)} p_2^{\max(i_2, j_2)} \dots p_n^{\max(i_n, j_n)}$
- ▶ Thus, we need to show $i_k + j_k = \min(i_k, j_k) + \max(i_k, j_k)$

Proof, cont.

- Show $i_k + j_k = \min(i_k, j_k) + \max(i_k, j_k)$

Computing GCDs

- ▶ Simple algorithm to compute gcd of a, b :
 - ▶ Factorize a as $p_1^{i_1} p_2^{i_2} \dots p_n^{i_n}$
 - ▶ Factorize b as $p_1^{j_1} p_2^{j_2} \dots p_n^{j_n}$
 - ▶ $\gcd(a, b) = p_1^{\min(i_1, j_1)} p_2^{\min(i_2, j_2)} \dots p_n^{\min(i_n, j_n)}$
- ▶ But this algorithm is not good because prime factorization is **computationally expensive!** (not polynomial time)
- ▶ Much more efficient algorithm to compute gcd, called the **Euclidian algorithm**

Insight Behind Euclid's Algorithm

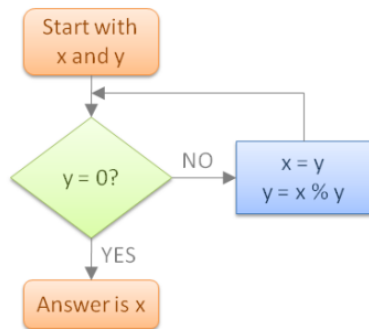
- ▶ **Theorem:** Let $a = bq + r$. Then, $\gcd(a, b) = \gcd(b, r)$
 - ▶ e.g., Consider $a = 12$, $b = 8$ and $a = 12$, $b = 5$
 - ▶ **Proof:** We'll show that a, b and b, r have the same common divisors – implies they have the same gcd.
- ⇒ Suppose d is a common divisor of a, b , i.e., $d|a$ and $d|b$
- ▶ By theorem we proved earlier, this implies $d|a - bq$
 - ▶ Since $a - bq = r$, $d|r$. Hence d is common divisor of b, r .
- ⇐ Now, suppose $d|b$ and $d|r$. Then, $d|bq + r$
- ▶ Hence, $d|a$ and d is common divisor of a, b □

Using this Theorem

Theorem: Let $a = bq + r$. Then, $\gcd(a, b) = \gcd(b, r)$

- ▶ Suggests following recursive strategy to compute $\gcd(a, b)$:
 - ▶ **Base case:** If b is 0, then $\gcd$ is a
 - ▶ **Recursive case:** Compute $\gcd(b, a \bmod b)$
- ▶ **Claim:** We'll eventually hit base case – why?

Euclidian Algorithm



- ▶ Find gcd of 72 and 20
- ▶ $12 = 72 \% 20$
- ▶ $8 = 20 \% 12$
- ▶ $4 = 12 \% 8$
- ▶ $0 = 8 \% 4$
- ▶ gcd is 4!

GCD as Linear Combination

- ▶ $\gcd(a, b)$ can be expressed as a **linear combination** of a and b
- ▶ **Theorem:** If a and b are positive integers, then there exist integers s and t such that:

$$\gcd(a, b) = s \cdot a + t \cdot b$$

- ▶ Furthermore, Euclidian algorithm gives us a way to compute these integers s and t (known as **extended Euclidian** algorithm)

Example

- ▶ Express $\gcd(72, 20)$ as a linear combination of 72 and 20
- ▶ First apply Euclid's algorithm (write $a = bq + r$ at each step):

1. $72 = 3 \cdot 20 + 12$

2. $20 = 1 \cdot 12 + 8$

3. $12 = 1 \cdot 8 + 4$

4. $8 = 2 \cdot 4 + 0 \Rightarrow \gcd \text{ is } 4$

- ▶ Now, using (3), write 4 as $12 - 1 \cdot 8$
- ▶ Using (2), write 4 as $12 - 1 \cdot (20 - 1 \cdot 12) = 2 \cdot 12 - 1 \cdot 20$
- ▶ Using (1), we have $12 = 72 - 3 \cdot 20$, thus:

$$4 = 2 \cdot (72 - 3 \cdot 20) - 1 \cdot 20 = 2 \cdot 72 + (-7) \cdot 20$$

Exercise

Use the extended Euclid algorithm to compute $\gcd(38, 16)$.

A Useful Result

- ▶ **Lemma:** If a, b are relatively prime and $a|bc$, then $a|c$.
- ▶ **Proof:** Since a, b are relatively prime $\gcd(a, b) = 1$
- ▶ By previous theorem, there exists s, t such that $1 = s \cdot a + t \cdot b$
- ▶ Multiply both sides by c : $c = csa + ctb$
- ▶ By earlier theorem, since $a|bc$, $a|ctb$
- ▶ Also, by earlier theorem, $a|csa$
- ▶ Therefore, $a|csa + ctb$, which implies $a|c$ since $c = csa + ctb$ □

Example

Lemma: If a, b are relatively prime and $a|bc$, then $a|c$.

- ▶ Suppose $15 \mid 16 \cdot x$
- ▶ Here 15 and 16 are relatively prime
- ▶ Thus, previous theorem implies: $15 \mid x$

Question

► Suppose $ca \equiv cb \pmod{m}$. Does this imply $a \equiv b \pmod{m}$?



Another Useful Result

- ▶ **Theorem:** If $ca \equiv cb \pmod{m}$ and $\gcd(c, m) = 1$, then $a \equiv b \pmod{m}$



Examples

- ▶ If $15x \equiv 15y \pmod{4}$, is $x \equiv y \pmod{4}$?
- ▶ If $8x \equiv 8y \pmod{4}$, is $x \equiv y \pmod{4}$?
- ▶