

CS311H: Discrete Mathematics

More Number Theory

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Linear Congruences

- ▶ A congruence of the form $ax \equiv b \pmod{m}$ where a, b, m are integers and x a variable is called a **linear congruence**.
- ▶ Given such a linear congruence, often need to answer:
 1. Are there any solutions?
 2. What are the solutions?
- ▶ **Example:** Does $8x \equiv 2 \pmod{4}$ have any solutions?
- ▶ **Example:** Does $8x \equiv 2 \pmod{7}$ have any solutions?
- ▶ **Question:** Is there a systematic way to solve linear congruences?

Determining Existence of Solutions

- ▶ **Theorem:** The linear congruence $ax \equiv b \pmod{m}$ has solutions iff $\gcd(a, m) \mid b$.
- ▶ Proof involves two steps:
 1. If $ax \equiv b \pmod{m}$ has solutions, then $\gcd(a, m) \mid b$.
 2. If $\gcd(a, m) \mid b$, then $ax \equiv b \pmod{m}$ has solutions.
- ▶ First prove (1), then (2).

Proof, Part I

If $ax \equiv b \pmod{m}$ has solutions, then $\gcd(a, m) \mid b$.

- ▶
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Proof, Part II

If $\gcd(a, m) \mid b$, then $ax \equiv b \pmod{m}$ has solutions.

- ▶ Let $d = \gcd(a, m)$ and suppose $d \mid b$
- ▶ Then, there is a k such that $b = dk$
- ▶ By earlier theorem, there exist s, t such that $d = s \cdot a + t \cdot m$
- ▶ Multiply both sides by k : $dk = a \cdot (sk) + m \cdot (tk)$
- ▶ Since $b = dk$, we have $b = a \cdot (sk) + m \cdot (tk)$
- ▶ Thus, $b \equiv a \cdot (sk) \pmod{m}$
- ▶ Hence, sk is a solution. \square

Examples

- ▶ Does $5x \equiv 7 \pmod{15}$ have any solutions?
- ▶ Does $3x \equiv 4 \pmod{7}$ have any solutions?

Finding Solutions

- ▶ Can determine existence of solutions, but how to find them?
- ▶ **Theorem:** Let $d = \gcd(a, m) = sa + tm$. If $d|b$, then the solutions to $ax \equiv b \pmod{m}$ are given by:

$$x = \frac{sb}{d} + \frac{m}{d}u \text{ where } u \in \mathbb{Z}$$

Example

Let $d = \gcd(a, m) = sa + tm$. If $d|b$, then the solutions to $ax \equiv b \pmod{m}$ are given by:

$$x = \frac{sb}{d} + \frac{m}{d}u \text{ where } u \in \mathbb{Z}$$

- ▶ What are the solutions to the linear congruence $3x \equiv 4 \pmod{7}$?

▶

Another Example

Let $d = \gcd(a, m) = sa + tm$. If $d|b$, then the solutions to $ax \equiv b \pmod{m}$ are given by:

$$x = \frac{sb}{d} + \frac{m}{d}u \text{ where } u \in \mathbb{Z}$$

- ▶ What are the solutions to the linear congruence $3x \equiv 1 \pmod{7}$?

▶

▶

Inverse Modulo m

- ▶ The **inverse of a modulo m** , written \bar{a} has the property:

$$a\bar{a} \equiv 1 \pmod{m}$$

- ▶ **Theorem:** Inverse of a modulo m exists if and only if a and m are relatively prime.
- ▶ **Proof:** Inverse must satisfy $ax \equiv 1 \pmod{m}$
- ▶
- ▶
- ▶ Does 3 have an inverse modulo 7?

Example

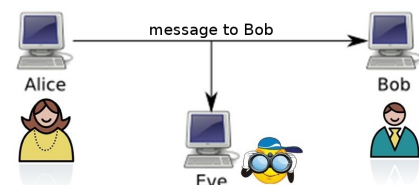
- ▶ Find an inverse of 3 modulo 7.
- ▶ An inverse is any solution to $3x \equiv 1 \pmod{7}$
- ▶ Earlier, we already computed solutions for this equation as:

$$x = -2 + 7u$$

- ▶ Thus, -2 is an inverse of 3 modulo 7
- ▶ $5, 12, -9, \dots$ are also inverses

Cryptography

- ▶ Cryptography is the study of techniques for secure transmission of information in the presence of adversaries



- ▶ How can Alice send secret messages to Bob without Eve being able to read them?

Private vs. Public Crypto Systems

- ▶ Two different kinds of cryptography systems:
 1. Private key cryptography (also known as **symmetric**)
 2. Public key cryptography (**asymmetric**)
- ▶ In private key cryptography, sender and receiver agree on **secret key** that both use to encrypt/decrypt the message
- ▶ In public key cryptography, a **public key** is used to encrypt the message, and **private key** is used to decrypt the message

Private Key Cryptography

- ▶ Private key crypto is classical method, used since antiquity
- ▶ Caesar's cipher is an example of private key cryptography
- ▶ Caesar's cipher is **shift cipher** where $f(p) = (p + k) \pmod{26}$
- ▶ Both receiver and sender need to know k to encrypt/decrypt
- ▶ Modern symmetric algorithms: RC4, DES, AES, ...
- ▶ **Main problem:** How do you exchange secret key in a secure way?

Public Key Cryptography

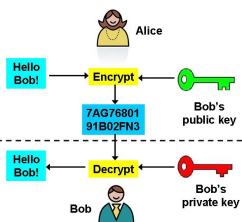
- ▶ Public key cryptography is the modern method: different keys are used to encrypt vs. decrypt message
- ▶ Most commonly used public key system is **RSA**
- ▶ Great application of number theory and things we've learned

RSA History



- ▶ Named after its inventors Rivest, Shamir, and Adleman, all researchers at MIT (1978)
- ▶ Actually, similar system invented earlier by British researcher Clifford Cocks, but classified – unknown until 90's

RSA Overview



- ▶ Bob has two keys: public and private
- ▶ Everyone knows Bob's public key, but only he knows his private key
- ▶ Alice encrypts message using Bob's public key
- ▶ Bob decrypts message using private key
- ▶ Since public key cannot decrypt, no one can read message except Bob

High Level Math Behind RSA

- ▶ In the RSA system, **private key** consists of two **very large prime numbers** p, q
- ▶ **Public key** consists of a number n , which is the product of p, q and another number e , which is relatively prime with $(p - 1)(q - 1)$
- ▶ Encrypt messages using n, e , but to decrypt, must know p, q
- ▶ In theory, can extract p, q from n using **prime factorization**, but this is intractable for very large numbers
- ▶ **Security of RSA relies on inherent computational difficulty of prime factorization**

Encryption in RSA

- ▶ To send message to Bob, Alice first represents message as a sequence of numbers
- ▶ Call this number representing message M
- ▶ Alice then uses Bob's public key n, e to perform encryption as:

$$C = M^e \pmod{n}$$

- ▶ C is called the **ciphertext**

RSA Decryption

- ▶ **Decryption key d** is the inverse of e modulo $(p-1)(q-1)$:

$$d \cdot e \equiv 1 \pmod{(p-1)(q-1)}$$

- ▶ **Decryption function:** $C^d \pmod{n}$
- ▶ As we saw earlier, d can be computed reasonably efficiently if we know $(p-1)(q-1)$
- ▶ However, since adversaries do not know p, q , they cannot compute d with reasonable computational effort!

Security of RSA

- ▶ The encryption function used in RSA is a **trapdoor function**
- ▶ Trapdoor function is easy to compute in one direction, but very difficult in reverse direction without additional knowledge
- ▶ Decryption without private key is very hard because requires prime factorization (which is intractable for large enough numbers)
- ▶ **Interesting fact:** There are efficient (poly-time) prime factorization algorithms for quantum computers (e.g., Shor's algorithm)
- ▶ If we could build quantum computers with sufficient "qubits", RSA would no longer be secure!