

Precise Reasoning for Programs Using Containers

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Containers

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General-purpose data structures for inserting, retrieving, removing, and iterating over elements



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- **Examples:** Array, vector, list, map, set, stack, queue, . . .

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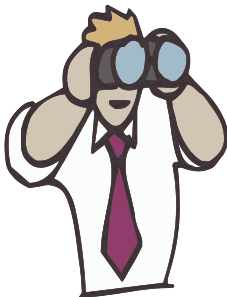


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 - Widely used; provided by common programming languages or standard libraries
- ⇒ Associate arrays in scripting languages, data structures provided by C++ STL, etc.



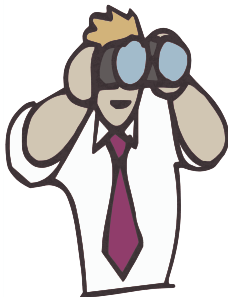
Precise static reasoning about
containers crucial for successful
verification

Observation #1



- Many different kinds of containers, varying in the **convenience** or **efficiency** of certain operations

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- Many different kinds of containers, varying in the **convenience** or **efficiency** of certain operations
- But **functionally**, there are only two kinds.

Classification of Containers

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1 Position-dependent Containers



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- Well-defined meaning of position

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⋮

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- Keys of arbitrary type
- Iteration order may be undefined

Observation #2:

Container Client



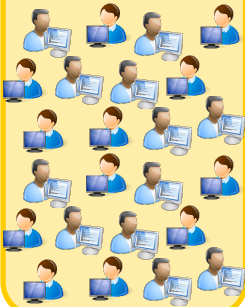
Implementation



- Orders of magnitude more **clients** of containers than there are container **implementations**

Observation #2:

Container Client



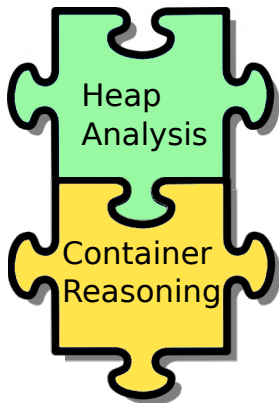
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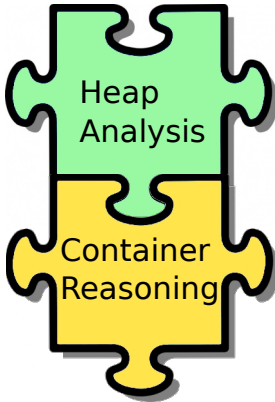
⇒ Need fully automatic, scalable techniques for reasoning about client-side use of container data structures

This Talk



Precise, fully-automatic technique that integrates **container reasoning** into **heap analysis**

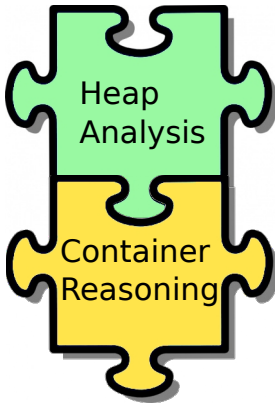
This Talk



Precise, fully-automatic technique that integrates **container reasoning** into **heap analysis**

- ① tracks **key-value correlations**

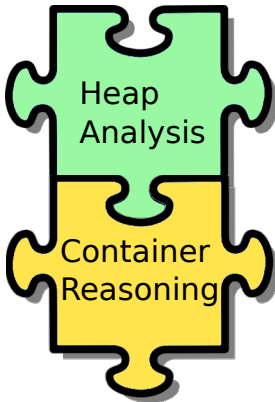
This Talk



Precise, fully-automatic technique that integrates **container reasoning** into **heap analysis**

- ① tracks **key-value correlations**
- ② can model **nested containers** in a precise way

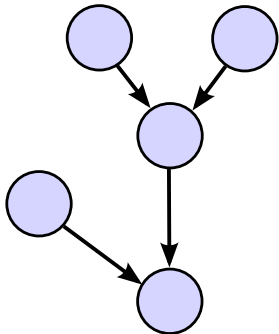
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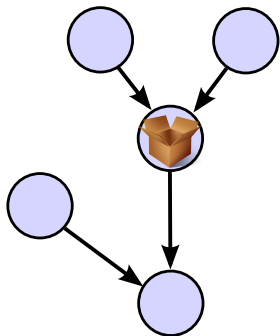
- ① tracks **key-value correlations**
- ② can model **nested containers** in a precise way
- ③ unifies **heap** and container analysis

Integrating Container Reasoning into Heap Analysis



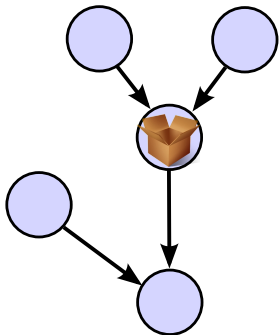
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Integrating Container Reasoning into Heap Analysis



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Integrating Container Reasoning into Heap Analysis



- To integrate containers into heap analysis, we model containers as **abstract memory locations** in the heap abstraction
- For precise, per-element reasoning, we model containers using **indexed locations** we introduced in ESOP'10 for reasoning about arrays

Indexed Locations

$\langle container \rangle_i$

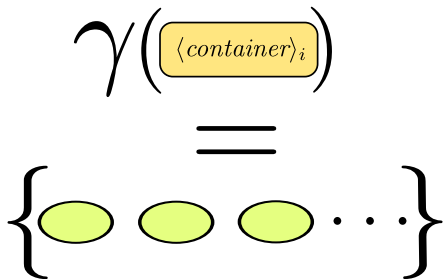
- Container represented using a single abstract location qualified by **index variable**

Indexed Locations

$\langle container \rangle_i$

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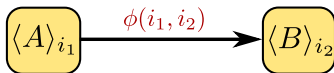
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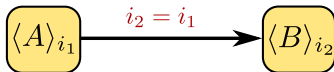
- Container represented using a single abstract location qualified by **index variable**
- Index variable ranges over **possible elements** of container
- **Key advantage:** Can refer to individual elements in container using only one abstract location

Symbolic Points-to Relations



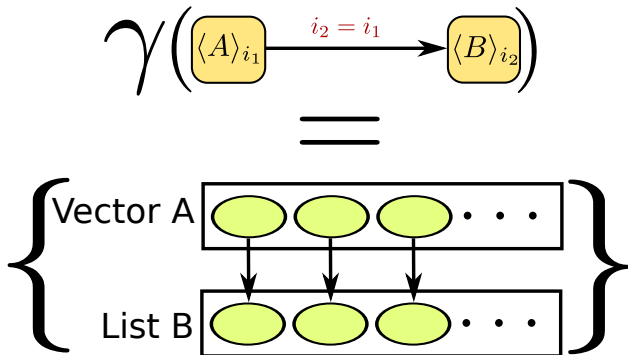
Points-to edges are qualified by **constraints** on index variables.

Symbolic Points-to Relations



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Symbolic Points-to Relations



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Modeling Value-Dependent Containers

Problem

- Natural representation for position-dependent containers

Modeling Value-Dependent Containers

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- But how do we represent points-to relations for value-dependent containers?

Modeling Value-Dependent Containers



Problem

- Natural representation for position-dependent containers
- But how do we represent points-to relations for value-dependent containers?

Solution

Introduce a level of indirection mapping **keys** to abstract **indices**

Key-to-Index Mapping for Value-Dependent Containers

- For value-dependent containers, any such **key-to-index mapping** M must satisfy the axiom:

$$\forall k_1, k_2. M(k_1) = M(k_2) \Rightarrow k_1 = k_2$$

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- Otherwise, distinct keys may map to same index, overwriting each other's value
- Thus, for soundness, M 's inverse is a function

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Two Alternatives

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 \Rightarrow make M a function

Using Invertible, Uninterpreted Functions

$$pos(\text{key}) = \chi$$



$$pos^{-1}(\chi) = \text{key}$$

Thus, map key to index in
abstract location using
invertible, uninterpreted
function

Simple Example

- Consider map `scores` mapping student names (strings) to a vector of their grades.

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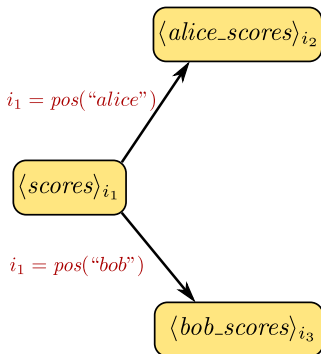
$\langle scores \rangle_{i_1}$

Simple Example

$\langle scores \rangle_{i_1}$

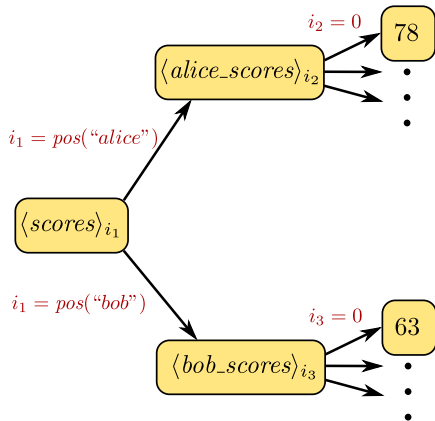
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- Consider map `scores` mapping student names (strings) to a vector of their grades.
- Map initially contains scores associated with two students: Alice and Bob
- Alice's first score is 78; Bob's first score is 63

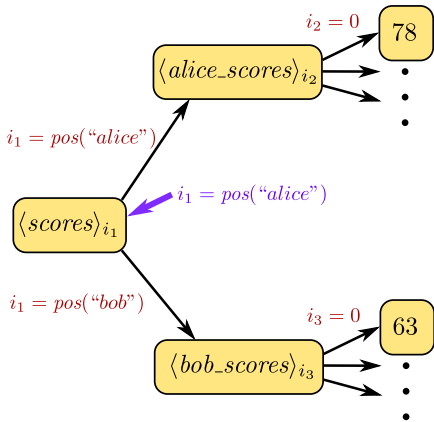


- We have seen how to represent containers



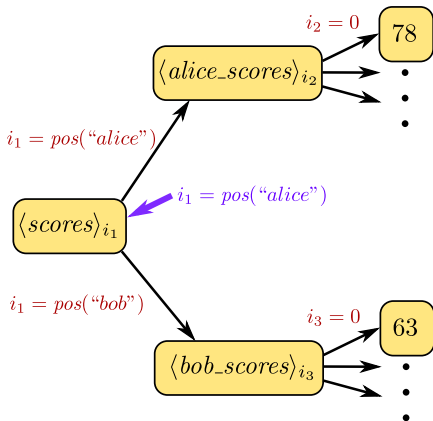
- We have seen how to represent containers
- But how do we **statically analyze** statements that manipulate them?

Simple Example: Reading from Containers



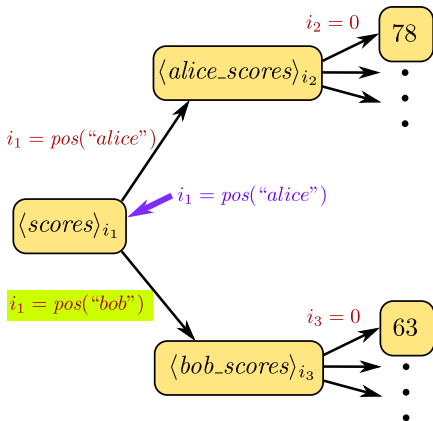
- What is the value of `scores["alice"][0]`?

Simple Example: Reading from Containers



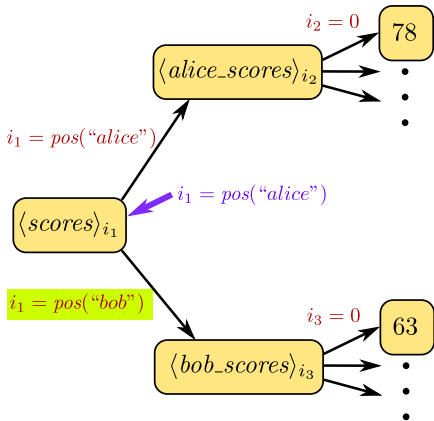
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- Determine where `scores` points to under $i_1 = pos("alice")$

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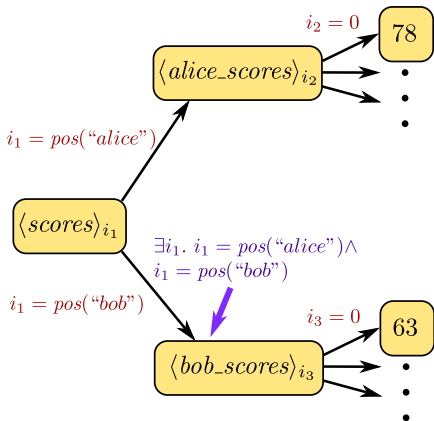
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- $i_1 = pos("bob")$

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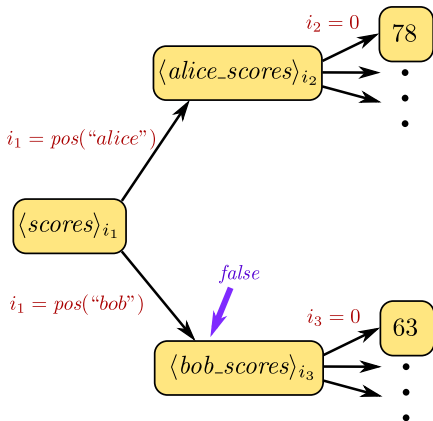
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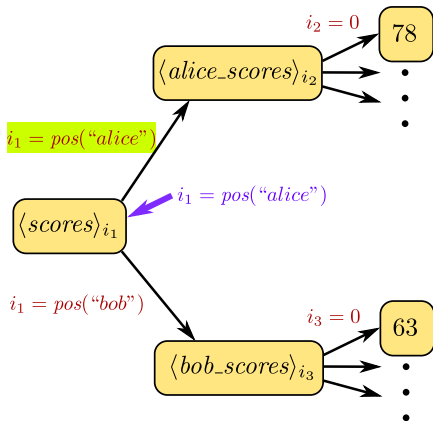
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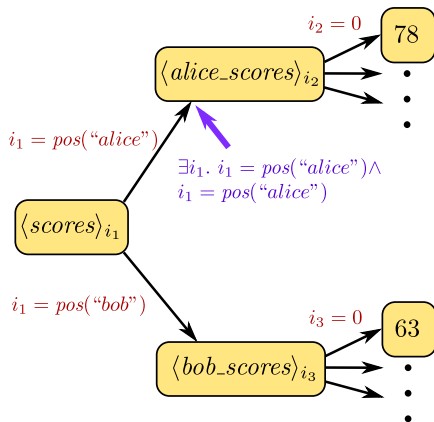
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- $\exists i_1. i_1 = pos("bob") \wedge i_1 = pos("alice")$

\Rightarrow **UNSAT** because `pos` is **invertible**

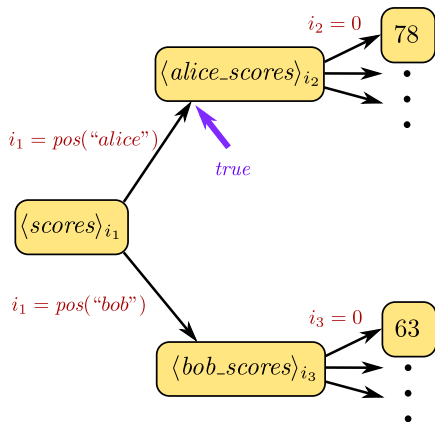
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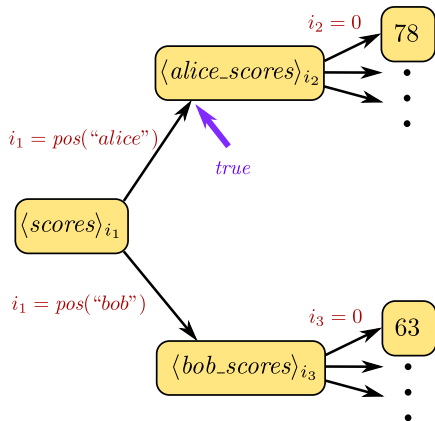
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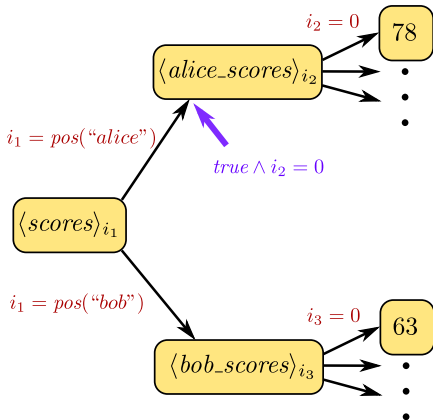


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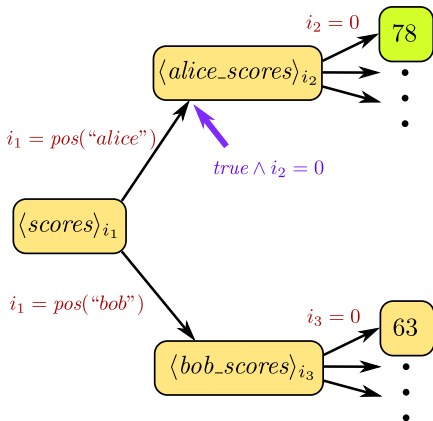
- Thus, entry for "alice" points to vector represented by $\langle alice_scores \rangle_{i_2}$

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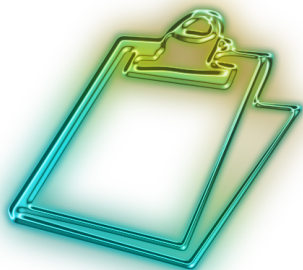
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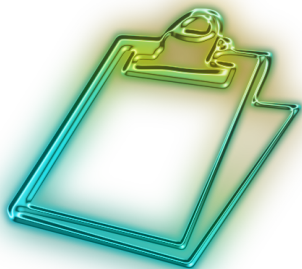
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Summary: Reading from Containers



- Statically analyzing reads from containers requires checking for **satisfiability** and **existential quantifier elimination**

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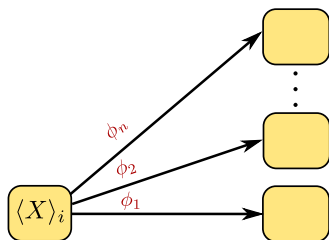


- Statically analyzing reads from containers requires checking for **satisfiability** and **existential quantifier elimination**
- Use of **invertible functions** for key-value mapping is crucial for precisely tracking key-value correlations



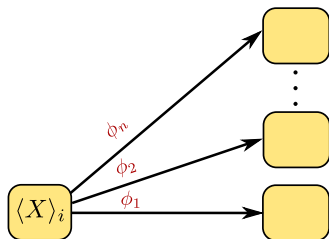
How do we analyze
stores to containers?

Writing to Containers



Consider storing object Y for key k in container X :

Writing to Containers

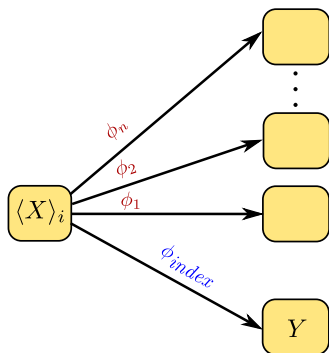


Consider storing object Y for key k in container X :

1 Compute

$$\phi_{index} : \begin{cases} i = k & \text{X position-dependent} \\ i = pos(k) & \text{X value-dependent} \end{cases}$$

Writing to Containers



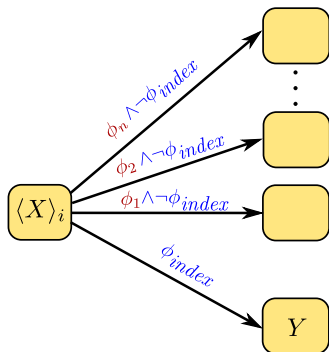
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Writing to Containers



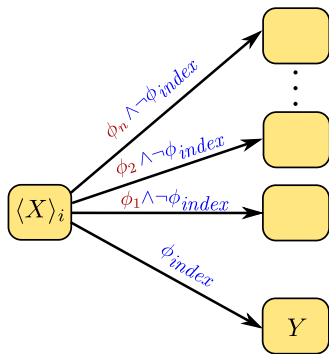
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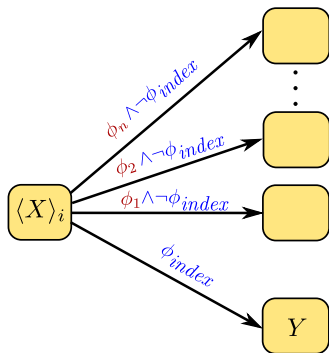
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Writing to Containers



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Need **bracketing constraints** $\langle \phi_{may}, \phi_{must} \rangle$ for sound negation

$$\Rightarrow \neg \langle \phi_{may}, \phi_{must} \rangle = \langle \neg \phi_{must}, \neg \phi_{may} \rangle$$

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- ⇒ Precise reasoning about nested containers requires **precise** reasoning about memory allocations

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- ⇒ Precise reasoning about nested containers requires **precise** reasoning about memory allocations
- Need to distinguish between allocations in different loop iterations or recursive calls

Consider the following example

```
for(int i=0; i<N; i++)  
    v.push_back(new map());
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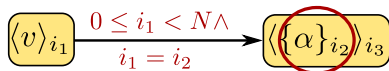
Difficulty

Statically unknown number of allocations

Allocations

Consider the following example

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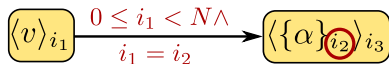
Solution

Model allocation with indexed location

Allocations

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Solution

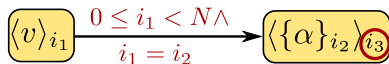
Model allocation with **indexed location**

- i_2 differentiates allocations from different loop iterations

Allocations

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for(int i=0; i<N; i++)  
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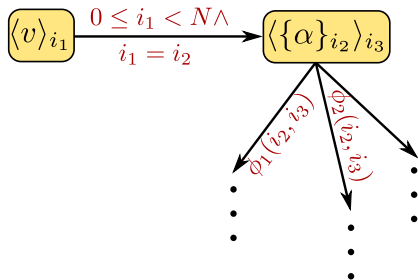
Model allocation with **indexed location**

- i_2 differentiates allocations from different loop iterations
- i_3 differentiates indices in map

Allocations

Consider the following example

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Solution

Model allocation with **indexed location**

- i_2 differentiates allocations from different loop iterations
- i_3 differentiates indices in map
- Outgoing edges from $\langle \{\alpha\}_{i_2} \rangle_{i_3}$ qualify both i_2 and i_3

Implementation



- Implemented heap/container analysis in our **Compass** program analysis framework for C and C++ programs

Implementation



- Implemented heap/container analysis in our **Compass** program analysis framework for C and C++ programs
- Analysis requires solving constraints in combined theory of **linear inequalities** over integers and **uninterpreted functions** and quantifier elimination
⇒ used our **Mistral** SMT solver



- Analyzed real open-source C++ applications using containers



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 - LiteSQL, 16,030 LOC



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 - DigiKam, 128,318 LOC

Application

```
switch (iFilterType)
{
    case CM_FILTERHIGHPASS:
    case CM_FILTERBWHIGHPASS:
        hrgn[2] = CreateEllipticRgn(x11, y11, x12, y12);
        break;
    case CM_FILTERLOWPASS:
    case CM_FILTERBLOWPASS:
        hrgn[0] = CreateEllipticRgn(x11, y11, x12, y12);
        hrgn[1] = CreateRectRgn(0, 0, 10, 10);
        hrgn[2] = CreateRectRgn(0, 0, 10, 10);
        CombineRgn(hrng[2], hrng[0], hrng[1], RGN_XOR);
        DeleteObject(hrng[0]);
        DeleteObject(hrng[1]);
        break;
    case CM_FILTERBANDSTOP:
        hrng[0] = CreateEllipticRgn(x21, y21, x22, y22);
        hrng[1] = CreateRectRgn(0, 0, 10, 10);
        hrng[2] = CreateRectRgn(0, 0, 10, 10);
        CombineRgn(hrng[2], hrng[0], hrng[1], RGN_XOR);
        DeleteObject(hrng[0]);
        DeleteObject(hrng[1]);
        break;
    case CM_FILTERBANDPASS:
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        hrng[1] = CreateRectRgn(0, 0, 10, 10);
        hrng[2] = CreateRectRgn(0, 0, 10, 10);
        CombineRgn(hrng[2], hrng[0], hrng[1], RGN_XOR);
        DeleteObject(hrng[0]);
        DeleteObject(hrng[1]);
        hrng[3] = CreateRectRgn(0, 0, 10, 10);
        CombineRgn(hrng[2], hrng[0], hrng[3], RGN_OR);
        DeleteObject(hrng[0]);
        DeleteObject(hrng[1]);
        break;
}
hrng[0] = CreateRectRgn(0, 0, 10, 10);
hrng[3] = CreateRectRgn(0, 0, 10, 10);
CombineRgn(hrng[3], hrng[0], hrng[2], RGN_AND);
// Tegnet endelige karakterer i radet
FillRgn(FaintInfo.hdc, hrng[3], hbrRed);
// Fjern de allokerede regionene, de er bare midlertidige
for(i=0; i<4; i++)
    if (hrng[i] != NULL)
        DeleteObject(hrng[i]);
```

Ran our Compass verification tool

- Detect all possible segmentation faults or run-time exceptions caused by:
 - null dereference errors
 - accessing deleted memory
- Also checked memory leaks

First Experiment:

- Represent containers as bags of values



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- Existing tools that **analyze programs of this size** use this abstraction



First Experiment:

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- To achieve this effect, we modeled containers using summary nodes



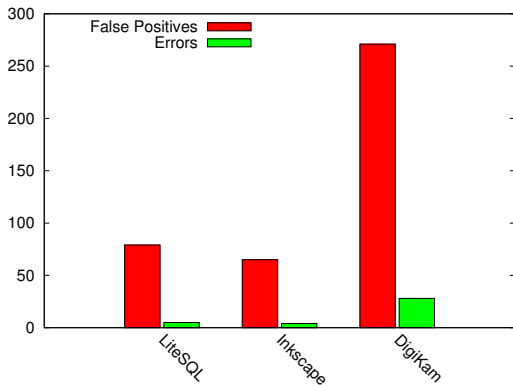
First Experiment



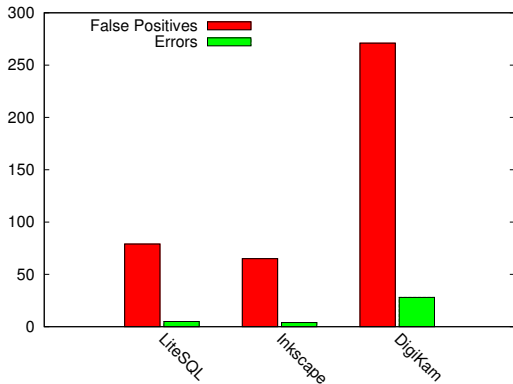
First Experiment:

- Represent containers as bags of values
 - Existing tools that **analyze programs of this size** use this abstraction
 - To achieve this effect, we modeled containers using summary nodes
- ⇒ Cannot track index-to-value correlations, modification to one container element contaminates all others

Containers as Bags



Containers as Bags



Conclusion



Treating containers as bags leads to **unacceptable** number of false alarms.

Second Experiment



Second Experiment:

- Used the techniques described in this talk:
indexed locations, symbolic points-to
relations

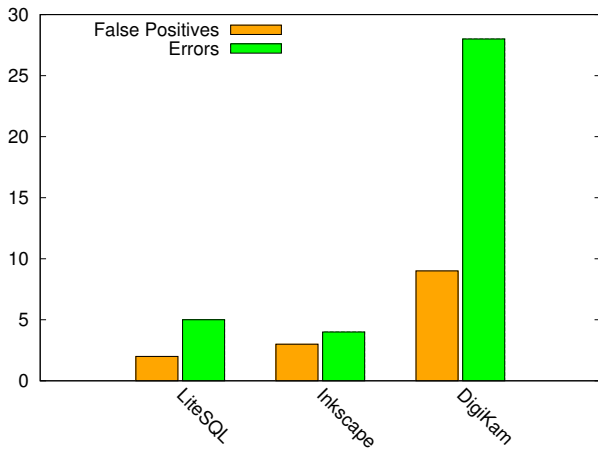
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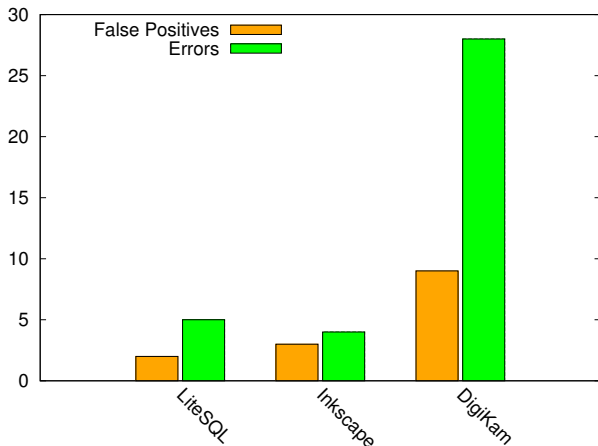
Second Experiment:

- Used the techniques described in this talk:
indexed locations, symbolic points-to
relations
- ⇒ Able to track key-value correlations;
precise reasoning about heap objects
stored in containers

Containers Modeled as Indexed Locations

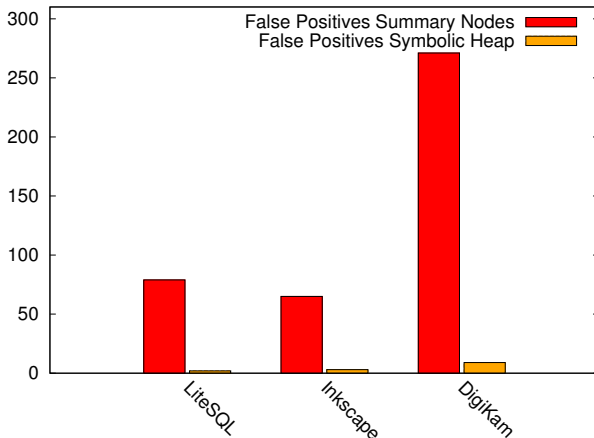


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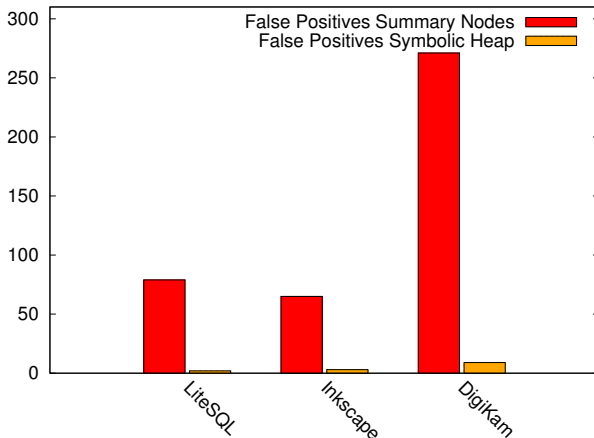


✓ Analysis reports very few false positives

Containers Modeled as Indexed Locations

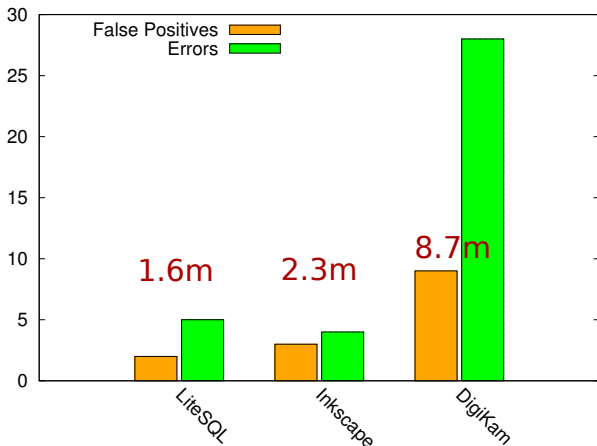


Containers Modeled as Indexed Locations



✓ More than an order of magnitude reduction compared to less precise analysis

Containers Modeled as Indexed Locations



✓ Cost of the analysis is tractable



- A sound, precise, and automatic technique for client-side reasoning about contents of an important family of data structures



- A sound, precise, and automatic technique for client-side reasoning about contents of an important family of data structures
- Precise reasoning for key-value correlations, nested data structures, and dynamic allocations



- A sound, precise, and automatic technique for client-side reasoning about contents of an important family of data structures
- Precise reasoning for key-value correlations, nested data structures, and dynamic allocations
- First practical verification of container- and heap-manipulating programs

Related Work

-  Dillig, I., Dillig, T., Aiken, A.:
Fluid Updates: Beyond Strong vs. Weak Updates.
In: ESOP. (2010)
-  Lam, P., Kuncak, V., Rinard, M.:
Hob: A Tool for Verifying Data Structure Consistency.
In: CC. 237–241
-  Reps, T.W., Sagiv, S., Wilhelm, R.:
Static Program Analysis via 3-Valued Logic.
In: CAV. (2004) 15–30
-  Deutsch, A.:
Interprocedural May-Alias Analysis for Pointers:
Beyond k-limiting.
In: PLDI. (1994) 230–241
-  Marron, M., Stefanovic, D., Hermenegildo, M., Kapur, D.:
Heap Analysis in the Presence of Collection Libraries.
In: PASTE. (2007)

