Fast LCF-Style Proof Reconstruction for Z3

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interactive theorem prover















SMT: Satisfiability Modulo Theories







































Z3 Proofs

Natural deduction proofs:



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34 axiom schemata and inference rules:





19 core rules

les

Core rules Asserted, MP, UNIT-RESOLUTION, ...

• mostly propositional reasoning



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Structure of the proof, simple semantics





REWRITE: simplification rules

Example $\overline{(P \longrightarrow Q) \longleftrightarrow (Q \lor \neg P)}$ $\overline{0 + x = x}$
Z3 Inference Rules



$T\operatorname{H-LEMMA:}$ inconsistency of theory atoms

Example

$$0 < x - 3y \land x - 3y < 0$$

False





Proof Reconstruction

Overall approach:

- one proof method for every Z3 inference rule
- depth-first traversal of Z3 proof



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Apply existing automated proof tools:

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Avoiding automated proof tools:

Speed-up of up to 3 orders of magnitude!

Reconstruction techniques:

single primitive inference rule or theorem instantiation



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- especific proof tools



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- schematic theorems
- conjunction/disjunction-reasoning
- oppositional reasoning
- Iinear arithmetic
- o quantifier elimination for integer/real arithmetic

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Example: UNIT-RESOLUTION

$$\frac{P \lor \neg Q \lor R}{P \lor R}$$

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Theorem memoization:

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Generalization:

- replace complex expressions by fresh variables
- avoids expensive preprocessing in the decision procedures
- enables shallow reasoning, more potential for theorem re-use

Logic		Z3		Isabelle		Rates	
		Med.	Med.		Med.		Time-
	#	Time	Size	#	Time	Succ.	out
AUFLIA+p	187	0.03 s	5 KB	187	0.06 s	100%	0%
AUFLIA-p	192	0.04 s	4 KB	190	0.06 s	98%	0%
AUFLIRA	189	0.02 s	16 KB	144	0.04 s	76%	0%
QF_AUFLIA	92	0.02 s	132 KB	49	2.82 s	53%	42%
QF_IDL	40	0.27 s	2 MB	19	3.25 s	47%	52%
QF_LIA	100	3.74 s	10 MB	26	45.27 s	26%	65%
QF_LRA	88	0.15 s	1 M B	55	23.87 s	62%	36%
QF₋RDL	52	0.70 s	1 MB	26	19.44 s	50%	50%
QF_UF	87	0.87 s	4 MB	73	3.40 s	83%	16%
QF_UFIDL	55	0.30 s	2 MB	8	3.59 s	14%	85%
QF_UFLIA	91	0.02 s	107 KB	85	2.12 s	93%	6%
QF_UFLRA	100	0.06 s	699 KB	100	3.35 s	100%	0%
Total	1273	3.66 s	13 MB	962	11.31 s	75%	19%

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Profiling:

- highly optimized implementation
- bottleneck: theory reasoning requires expensive proof search
- $\bullet~50\%$ of the runtime is spent on 15% of all Z3 proof steps





Conclusion

LCF-style proof reconstruction for Z3:

- highly optimized
- feasible and efficient
- slower than proof search in Z3
- Z3's proofs could be easier to check (costly theory reasoning)

Benefits for Isabelle/HOL and HOL4:

powerful automation

Benefits for Z3:

- proof checker
- identification of several bugs
Proof bug

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Documentation problem

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Soundness bug

$$(\forall x^T. x = c^T) \land a^T \neq b^T$$