

Case-Analysis for Rippling and Inductive Proof

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Interactive Theorem Proving Conference
12 July 2010

Introduction: Automating Inductive Proofs

- Induction: reasoning about recursion.
 - Programs, data-structures...
- Automation difficult. Guidance needed:
 - e.g. case-splitting for inductively defined datatypes.

Aim: Automation of theorems involving conditional definitions.

- Half of the functions in Isabelle's list library involve conditional statements.
- e.g. *member*, *delete*, *subtraction*.

Proof Planning and Rippling

- **Proof Planning:** Families of proofs with similar structure.
 - e.g. inductive proofs.
- **Rippling:** Heuristic for guiding rewriting in step cases.
 - Annotate differences between induction hypothesis and conclusion.
 - Rewrites must reduce differences.
 - Top of term tree: Get instance of IH.
 - Position of \forall quantified variable in IH.
 - Fully automatic and guarantees termination.
 - Does not require rewrite rules to be oriented.

Case Analysis for Rippling

- **IsaPlanner:** Higher-order proof-planner for Isabelle.
- Extend rippling in IsaPlanner to cover case analysis for:
 - If-statements.
 - Case-statements over datatypes.
- **Retain termination:**
 - Ripple-step becomes: rewriting + case-splitting.

A Simple Example

Definition of *max*:

$$\text{max } 0 \ y = \ y \quad (1)$$

$$\begin{aligned} \text{max } (\text{Suc } x) \ y = & \text{case } y \text{ of } 0 \Rightarrow \text{Suc } x \\ & | \text{Suc } z \Rightarrow \text{Suc}(\text{max } x \ z) \end{aligned} \quad (2)$$

Commutativity of *max*:

Inductive hypothesis (IH): $\forall b. \text{max } a \ b = \text{max } b \ a$

Step-case goal: $\text{max } \boxed{\text{Suc } a} \ [b'] = \text{max } [b'] \ \boxed{\text{Suc } a}$

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Apply (2):

$$\boxed{\text{case } b' \text{ of } 0 \Rightarrow \text{Suc } a \mid \text{Suc } z \Rightarrow \text{Suc}(\text{max } a \ \lfloor z \rfloor)} = \text{max } \lfloor b' \rfloor \ \boxed{\text{Suc } \underline{a}}$$

Applying the Split

Isabelle automatically derives splitting rules for case-statements for each datatype:

$$\begin{aligned} \llbracket ?n = 0 \implies ?P(?f_1); \forall x. (?n = \text{Suc } x) \implies ?P(?f_2 \ x) \rrbracket \implies \\ ?P(\text{case } ?n \text{ of } 0 \Rightarrow ?f_1 \mid (\text{Suc } x) \Rightarrow (?f_2 \ x)) \end{aligned}$$

$?P$ matches context of case-statement.

- Case-split as resolution step.
- Interactive setting: User gives instantiation of $?P$.
- Automatic proof: $?P$ in head position - match any term, many trivial unifiers, large search space.

Restricted Unification

$$\begin{array}{c}
 ?P(\underbrace{\text{case } ?n \text{ of } 0 \Rightarrow ?f_1 \mid (\text{Suc } x) \Rightarrow (?f_2 \ x)}_{\text{meta-variable argument}}) \\
 \underbrace{\text{case } b' \text{ of } 0 \Rightarrow \text{Suc } a \mid \text{Suc } z \Rightarrow \text{Suc}(\text{max } a \ z)}_{\text{subterm of interest}} = \text{max } b' \ \text{Suc } a
 \end{array}$$

- Only apply to terms containing argument of meta-variable $?P$.
- Find instantiation for $?P$, then safe to apply regular resolution.
- Algorithm using *Zipper*.
- Traverse term, find matching subterm. Zipper keeps track of its context.
- Use context of subterm to provide instantiation of $?P$:
 $\lambda v. v = \text{max } b' \ \text{Suc } a$

Example cont.

A ripple step with case-analysis:

$$\left\{ \begin{array}{l} \max \boxed{\text{Suc } \underline{a}} \lfloor b' \rfloor = \max \lfloor b' \rfloor \boxed{\text{Suc } \underline{a}} \\ \quad \downarrow \text{ Rewrite using: max def.} \\ \text{case } b' \text{ of } 0 \Rightarrow \text{Suc } a \mid \text{Suc } z \Rightarrow \text{Suc}(\max a z) = \max b' \text{ Suc } a \\ \quad \downarrow \text{ Apply case-split.} \end{array} \right.$$

$$b' = 0 \implies \text{Suc } a = \max 0 (\text{Suc } a) \quad (3)$$

$$b' = \text{Suc } z \implies \boxed{\text{Suc}(\underline{\max a \lfloor z \rfloor})} = \max \lfloor \text{Suc } z \rfloor \boxed{\text{Suc } \underline{a}} \quad (4)$$

Goal 3 is solved by simplification (no rippling embedding).

Rippling continues on goal 4.

Evaluation

- Case-statements common, not previously covered by rippling.
- Implementation in IsaPlanner:
Rippling + Case Analysis + Lemma Calculation
- Corpus of 87 inductive theorems. If- and case statements.
 - Lists, natural numbers, binary trees.
 - Isabelle library, inductive TP literature, dependently typed programming.
- Prover given only function definitions. No extra lemmas supplied.

Evaluation

- 47/87 new theorems proved automatically.
- Remaining theorems:
 - More sophisticated reasoning about side-conditions.
 - Conjecturing of conditional lemmas.
- Example failed proof: $\text{sorted}(\text{insertionSort}(l))$
- Need lemma with assumption:
 $\text{sorted } m \implies \text{sorted}(\text{insert } x \ m).$

Comparison to Simplification-Based Technique

- Isabelle's simplifier allows splitting on if-statements.
- Case-statements not attempted: may cause non-termination.

Coverage:

- Proved by Induction+Simp technique: 37
- Proved by Rippling: 47

Termination:

- Rippling terminates on all examples.
- Induction+Simp often fails to terminate:
 - Proofs it cannot solve.
 - When asked for alternative proofs.
 - Stuck trying to conjecture increasingly complex lemmas.

Other Approaches

- **Recursion Analysis:** Choose induction scheme avoiding need for case-splits.
 - **ACL2, VeriFun:** No datatypes, destructor style, recursion on several arguments instead of case statements.
- **Isabelle:** HO, datatypes, constructor style.
Want to work directly with these.
- **Further Work:** Automatic construction/selection of induction schemes in IsaPlanner.
- May still need case-analysis even with more elaborate induction schemes.
 - Case-statement introduced by auxiliary lemma, rewrite from other conditional function definition.

Summary

- Case-analysis needed for inductive proofs about conditional functions.
- Fully automatic technique implemented in IsaPlanner.
- Incorporated with rippling.
- Ensures termination also in presence of case-statements over datatypes.
- Proves 47/87 theorems in evaluation corpus.
- Further work:
 - More sophisticated reasoning with assumptions
 - Conditional lemma discovery.