

Ultrawrap: Using SQL Views for RDB2RDF

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Goals

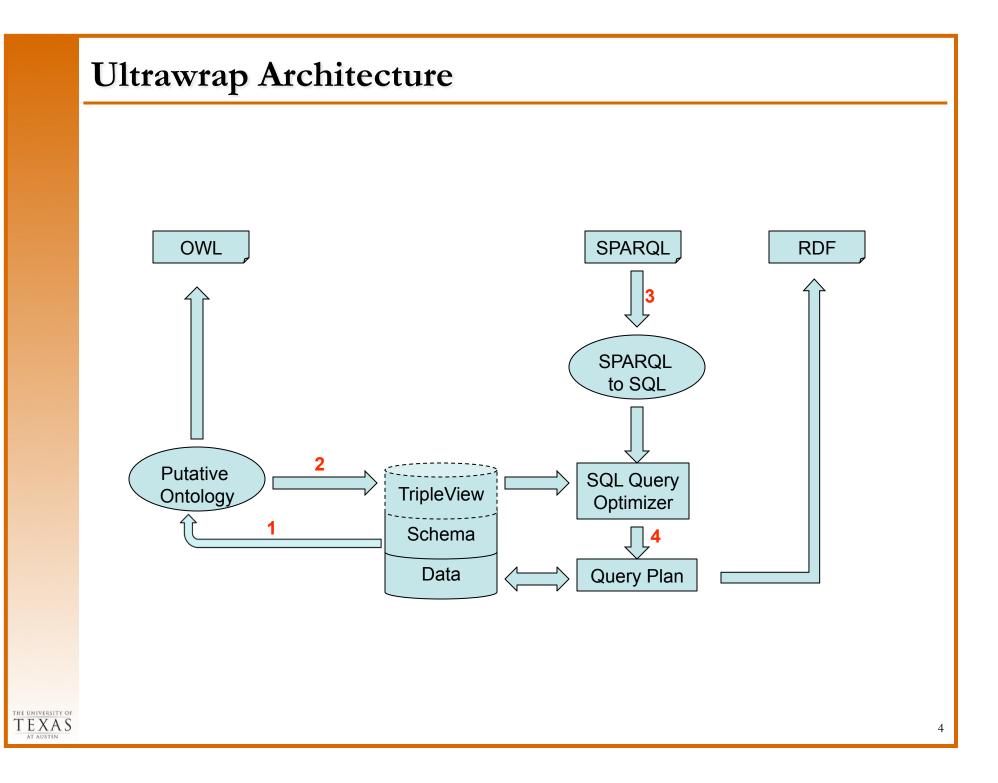
- ✓ Automatically create SPARQL endpoint for legacy relational databases.
- ✓ Real-time consistency between the relational and RDF data
- ✓ Making maximal use of existing SQL infrastructure
- Research question: Do existing commercial SQL query engines already subsume all the algorithms needed to support effective SPARQL execution on relational data?

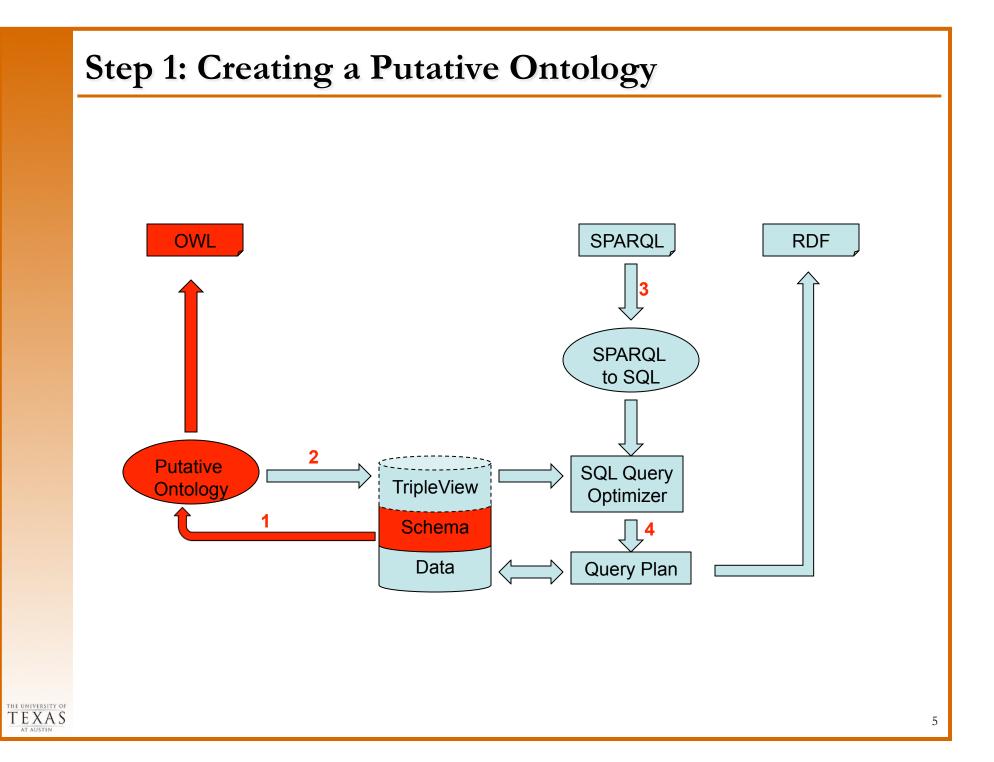


Ultrawrap Architecture

- Compile Time
 - 1. Create Putative Ontology (PO)
 - 2. Create Virtual Triple Store
- Run Time
 - 3. Naïve SPARQL to SQL translation
 - 4. SQL Optimizer is the rewriter
- Future
 - 5. Putative Ontology to Domain Ontology mapping







Ontology Quality? I.e. putative

- Putative: "commonly regarded as such"
- Putative Ontology (PO): automatic syntactic transformation from a data source schema to an ontology
 - data or information source ontology
- Evidence: SQL schema from E-R models make "interesting" ontologies



FOL rules transform SQL DDL to OWL

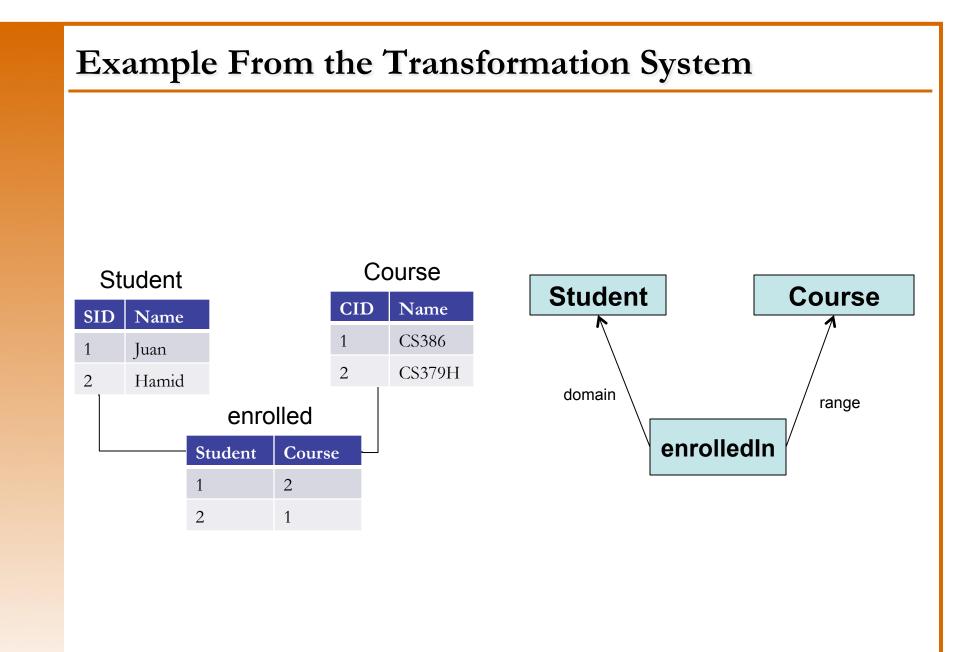
- Full mapping in Datalog
 - Stratified and safe
- Proof of total coverage of all key combinations

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Documentation

- Tirmizi, S. H., Sequeda, J.F., and Miranker, D.P.
 Translating SQL Applications to the Semantic
 Web. In Proceedings of the 19th international Conference on Database and Expert Systems Applications (DEXA2008).
- Sequeda, J.F. (1), Tirmizi, S.H.(1), Corcho, O. (2), Miranker, D.P. (1). "Direct Mapping SQL Databases to the Semantic Web: A Survey". The University of Texas at Austin, Department of Computer Sciences(1), Universidad Politecnica de Madrid(2). Report# TR-09-04 (regular tech report). January 19th, 2009. 35 pages. ftp://ftp.cs.utexas.edu/pub/techreports/tr09-04.pdf

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Transformation System

Rel(r) \land FK(xtr,r,_,t) \land FK(xsr,r,_,s) \land xtr \neq xsr \land Attr(y,r) \land *BinRel(r,s,t)* \neg NonFK(y,r) \land FK(z,r,_,u) \land fkey(z,r,u) \in {fkey(xsr,r,s), fkey(xtr,r,t)}

create table **STUDENT**{ SID integer primary key, **NAME** varchar not null }

create table **COURSE**{ CID integer primary key, **NAME** varchar not null }



<owl:Class rdf:ID="Student"/>

<owl:Class rdf:ID="Course"/>

 $ObjP(r,s,t) \leftarrow BinRel(r,s,t) \land Rel(s) \land Rel(t) \land \neg BinRel(s,_,) \land \neg BinRel(t,_,)$

create table **ENROLLED**{ Student integer foreign key <owl:ObjectProperty rdf:ID="enrolledIn"> references STUDENT(SID), <rdfs:domain rdf:resource="#Student"/> **Course** integer foreign key <rdfs:range rdf:resource="#Course"/> references COURSE(CID), </owl:ObjectProperty> constraint REG PK primary key (Student, Course) } HE UNIVERSITY OF EXAS

Proof of total coverage of all key combinations

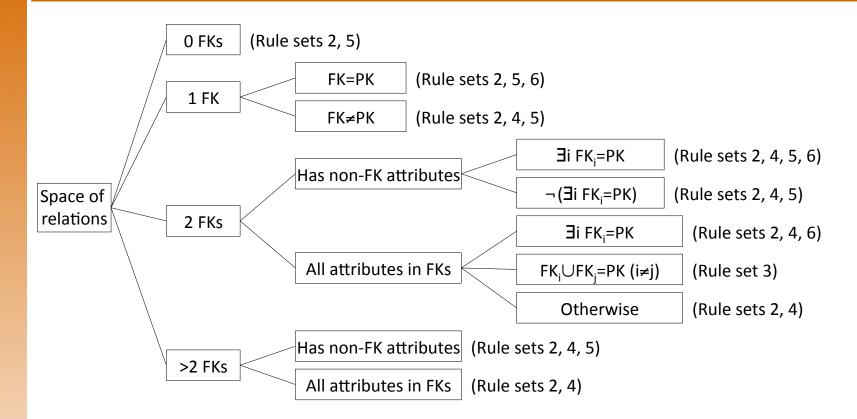
- PK: a relation only has a Primary Key
- C-PK: a relation only has a composite Primary Key
- S-FK: a relation only has one Foreign Key
- N-FK: a relation has at least two or more Foreign Keys

Grammar

- $E \rightarrow PK + T \mid C PK + T$
- E → S-FK
- $E \rightarrow N$ -FK
- $T \rightarrow S$ -FK | N-FK



LR(0) Item set construction represents all possibilities



The tree describes the complete space of relations when all possible combinations of primary and foreign keys are considered.



All Key Combinations Enumerated

- PK + S-FK: a relation has a Primary Key and only one Foreign Key
 - PK = S-FK: the Foreign Key is the Primary Key
 - PK \cap S-FK = 0: the Foreign Key and the Primary Key do not share any attributes
- PK + N-FK: a relation has a Primary Key and two (2) Foreign Keys
 - PK \cap N-FK = 0: the Foreign Key and the Primary Key do not share any attributes
 - PK \subset N-FK: one of the Foreign Keys is also the Primary Key
- PK + N-FK: a relation has a Primary Key and more than two (> 2) Foreign Keys
 - PK \cap N-FK = 0: the Foreign Key and the Primary Key do not share any attributes
 - PK \subseteq N-FK: one of the Foreign Keys is also the Primary Key
- C-PK + S-FK: a relation has a Composite Primary Key and only one Foreign Key.
 - C-PK \cap S-FK = 0: the Foreign Key and the Primary Key do not share any attributes
 - S-FK \subseteq C-PK: the Foreign Key is part of the Primary Key
- C-PK + N-FK: a relation has a Composite Primary Key and two (2) Foreign Keys
 - C-PK \cap N-FK = 0: all the Foreign Keys and the Primary Key do not share any attributes
 - N-FK \subseteq C-PK: all the Foreign Keys are part of the Primary Key
 - C-PK ∩ N-FK ≠0, C-PK N-FK ≠ 0, N-FK C-PK ≠ 0: The Foreign Keys and Primary Key share common attributes
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Ultrawrap creates PO automatically

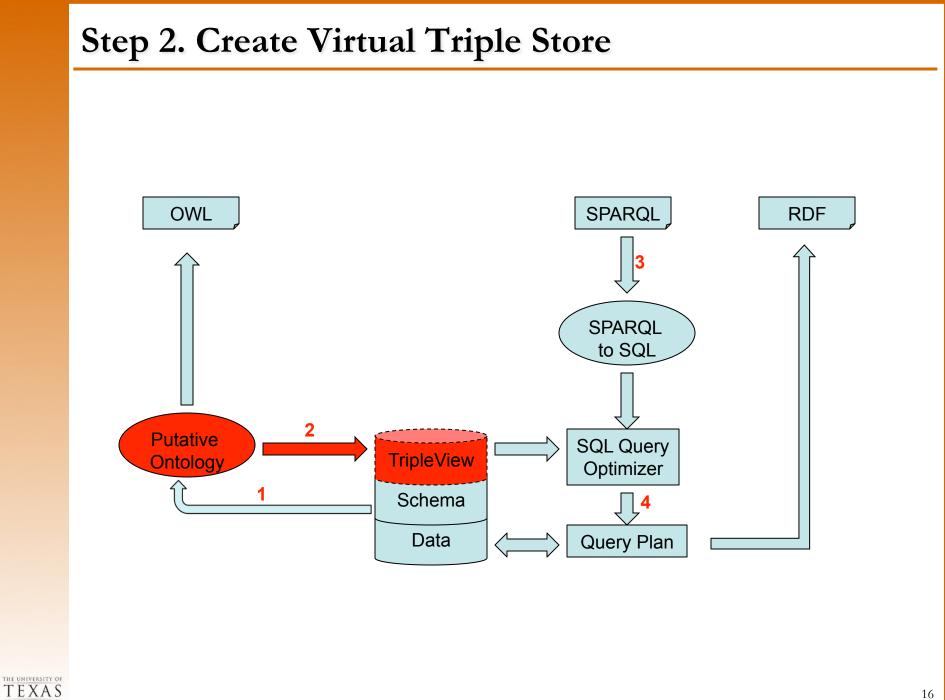
- Reads the Data Dictionary
 - Specific to each vendor
- Currently supporting Microsoft SQL Server and MySQL



Quality of a Putative Ontology

- Quality of the PO depends on the SQL DDL
 - Is it normalized?
 - Are all the constraints explicit?
- If the Quality is sufficient, all we need to do is rename
- Need to map the Putative Ontology to a Domain Ontology
- Evidence: SQL schema from E-R models make "interesting" ontologies
- SQL schemas made without any previous modeling make "poor" ontologies





Step 2. Create Virtual Triple Store

- Represent all relational data as triples using a view definition
 - Promise of avoiding self joins (optimizer will do this)
- Triple table approach: one table with three columns (s,p,o)
 - No symbol/lookup table. Strings are in the view
- Actually, the view is (s,spk, p, o, opk) where spk and opk are the index values
 - Optimizer needs to know the index values



Step 2. Create Virtual Triple Store

• Create SELECT statements that output triples

SELECT "Product"+id as s, id as spk, "product_label" as p, label as o, null as opk FROM Product

S	SPK	Р		OPK
Product1	1	product_label	Label of Product 1	null
Product2	2	product_label	Label of Product 2	null

SELECT "Product"+ProductID as s, ProductID as spk, "product_productfeature"
 as p, "ProductFeature"+ProductFeatureID as o, ProductFeatureID as opk
 FROM ProductFeatureProduct

S	SPK	Р	0	ОРК
Product1	1	product_productfeature	ProductFeature45	45
Product1	1	product_productfeature	ProductFeature98	98

• Use the PO as basis to create all the SELECT statements

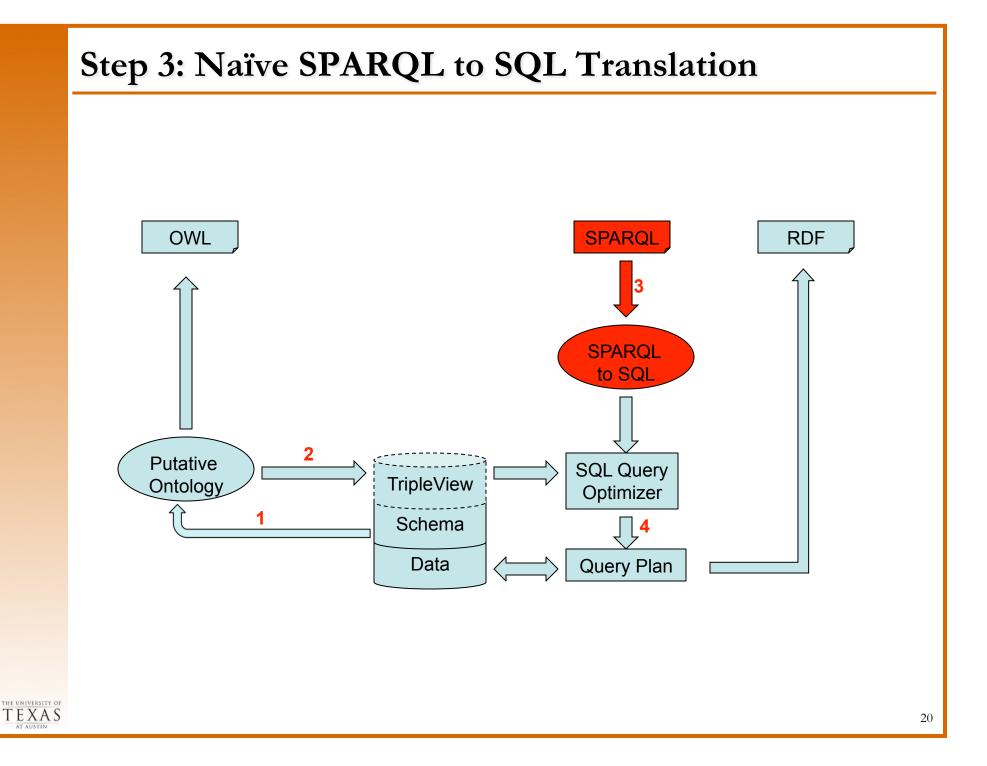
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Step 2. Create Virtual Triple Store

Triple View is a union of all the SELECT statementss
 CREATE VIEW TripleView(s,spk, p, o, opk) as
 SELECT "Product"+id as s, id as spk, "rdf:type" as p, "Product" as o, null
 as opk FROM Product
 UNION
 SELECT "Product"+id as s, id as spk, "label" as p, label as o, null as opk
 FROM Product
 UNION
 SELECT "Product"+ProductID as s, ProductID as spk, "product_productfeature"
 as p, "ProductFeature"+ProductFeatureID as o, ProductFeatureID as opk
 FROM Product

UNION ...

 BSBM generates ~80 select statements in order to represent all relational data as triples



Step 3: Naïve SPARQL to SQL Translation

SPARQL Query

```
SELECT ?product ?label
WHERE{
?product producttype_product
    ProductType47.
?product product_label ?label.
?product product_productfeature
    ProductFeature76.
?product product_productfeature
    ProductFeature4242.
?product product_propertyNum1 ?v.
FILTER (?v >500)
}
```

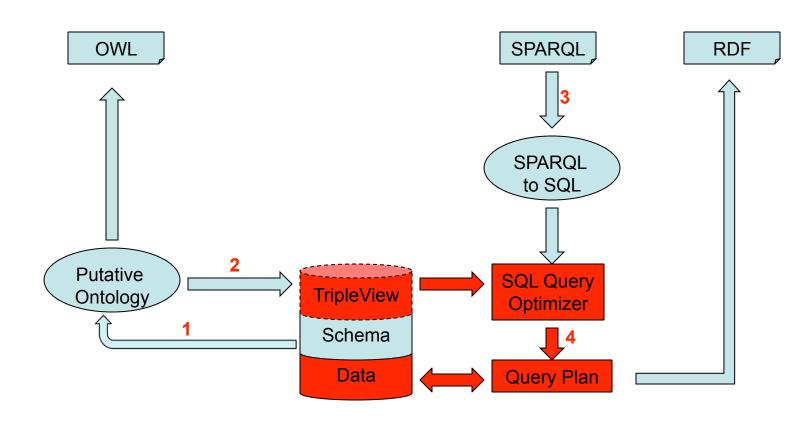
SQL Query on the Triple View

```
SELECT t1.o as product, t2.o as
   label
FROM TripleView t1, t2, t3, t4, t5
WHERE
    t1.p = 'producttype product'
and t1.opk = 47
and t2.p = 'product label'
and t3.spk = t1.spk
and t3.p = 'product productfeature'
and t3.opk = 76
and t4.spk = t1.spk
and t4.p = 'product productfeature'
and t4.opk = 4242
and t5.spk = t1.spk
and t5.p = 'product propertyNum1'
and t5.0 > 500
```

Syntactic transformation from a SPARQL query to an equivalent SQL query on the Triple View



Step 4: SQL Query Optimizer is the Rewrite system



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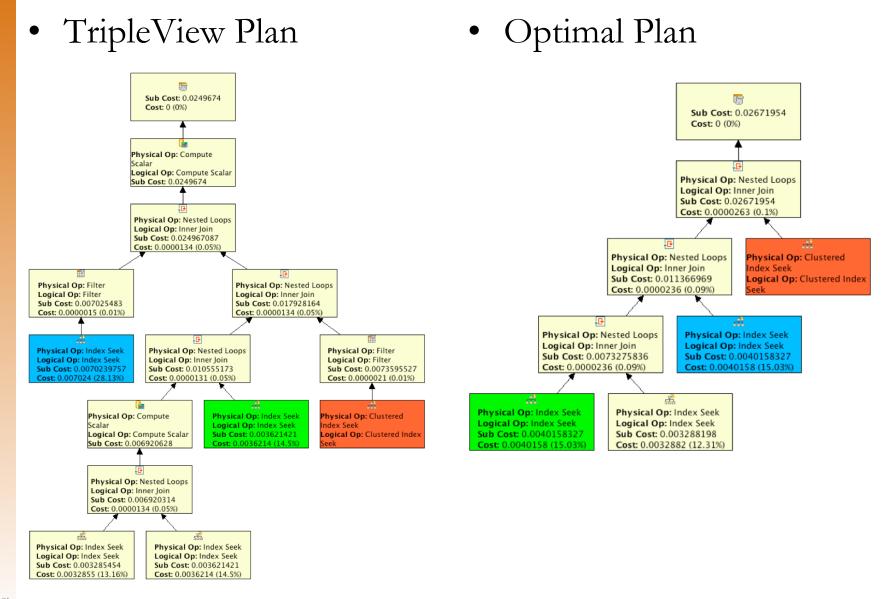
```
TripleView(1, label, ABC) :- Product(1,ABC, _, _)
TripleView(1, propNum1, 1) :- Product(1,_, 1, _)
TripleView(1, propNum1, 2) :- Product(1,_, _, 2)
```

```
SQL Query on the TripleView
Query(X, Y):-TripleView(X, label, Y),
TripleView(X, propNum1, 1),
TripleView(X, propNum2, 2)
```

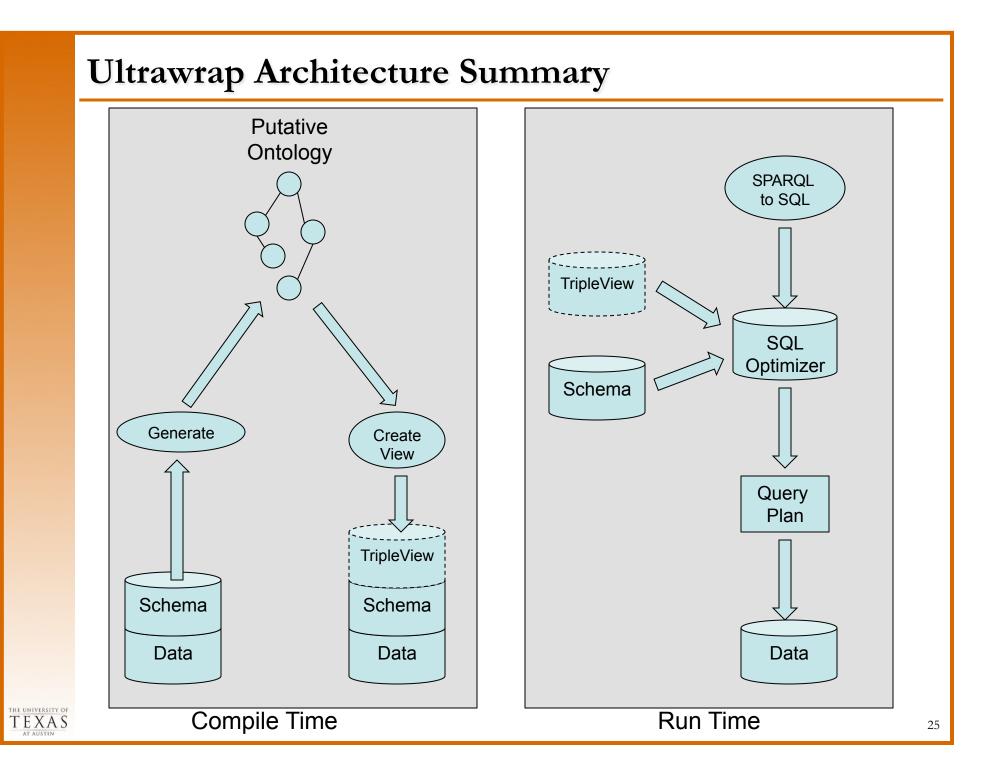
```
SQL Query on the Relational Data
SELECT id, label FROM product
WHERE propNum1 = 1 and propNum2 = 2
Query(X, Y) :- Product(X, Y, 1, 2)
```

```
Evaluate SQL Query on the TripleView
Query(X,Y):-Product(X,Y,1,_),Product(X,Y,_, 2)
Query(X,Y):- Product(X, Y, 1, 2)
```

Step 4: SQL Query Optimizer is the Rewrite system



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Current Implementation

- Running on Microsoft SQL Server
- Jess Rule Engine
- Initial test on BSBM on 1 million triples, execution time is close to running time of native SQL queries on relational data

Query	1	2	3	4	5	6	7	8	9	10	11	12
Jena SDB	4.41	4.33	6.27	7.12	12.36	1.45	11.94	6.69	8.38	5.39	2.76	4.34
Sesame	2.49	0.86	3.52	3.78	7.31	1.76	15.51	3.02	1.17	3.63	1.49	0.65
Virtuoso RDF Views	8.29	2.77	9.79	16.13	1.89	0.09	16.59	6.83	2.14	6.96	9.50	3.44
D2R Server	5.03	5.28	7.93	7.63	222.73	0.94	10.96	12.46	13.37	7.16	30.61	2.55
Ultrawrap	1	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00
Native SQL	0.94	0.67	0.90	0.80	1.09	0.62	0.67	0.72	1.03	1.02	0.94	0.30

Current Work

- Shifted problem to ontology-to-ontology mapping
 - Version 0: query only the Putative Ontology
 - Version 1: Manually mapping layer between Domain Ontology to Putative Ontology
 - Version 2: Automatic identify mappings
- Testing on Oracle, PostgreSQL, Virtuoso
- Road Map
 - Dec 2009: Version 0
 - Feb 2010: Version 1 running on other RDBMS



Thank You





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