

# **Mechanically checked proof on Dijkstra's shortest path algorithm**

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# Introduction

- Dijkstra's shortest path algorithm: a classical algorithm to find the shortest path between two vertices in a finite graph with non-negative weighted edges
- Directed Finite Graph with non-negative weighted edges
- Correctness of the algorithm: "if both vertices  $a$  and  $b$  are in the graph  $g$ , then the algorithm does return a shortest path from  $a$  to  $b$  in the graph  $g$ "

# Algorithm

1.  $\lambda(u) \leftarrow 0$ ; for each vertex  $t$  other than  $u$  in  $V$ ,  $\lambda(t) \leftarrow \infty$ ;  
and  $T \leftarrow V$ ;
2. Let  $s$  be a vertex in  $T$  such that  $\lambda(s)$  is minimum;
3. If  $s = v$ , stop (or If  $T = \{\}$ , stop);
4. For every edge from  $s$  to  $t$ , if  $t \in T$  and  
 $\lambda(t) > \lambda(s) + wt(st)$ , then  $\lambda(t) \leftarrow \lambda(s) + wt(st)$ ;
5.  $T \leftarrow T - \{s\}$  and go to step 2.

# Formalization

- Graph representation: an association list ((u1 (v1 . w1) (v2 . w2) ...) ...)
- path table pt: ((u . path-from-a-to-u) ...)
- Function returns the result

```
(defun dijkstra-shortest-path (a b g)
  (let ((p (dsp (all-nodes g) (list (cons a (list a))) g)))
    (path b p)))
```

- Function maintains the iteration

```
(defun dsp (ts pt g)
  (cond ((endp ts) pt)
        (t (let ((u (choose-next ts pt g)))
              (dsp (del u ts)
                  (reassign u (neighbors u g) pt g)
                  g))))))
```

# Formalization

1. Let  $ts$  be initially all vertices in  $g$ ;
2. Let  $pt$  be initially  $(list (cons a (list a)))$ ;
3.  $(path\ n\ pt)$  returns the already discovered path associated with  $n$  in  $pt$ , i.e. initially  $(path\ a\ pt) = (list\ a)$  and  $(path\ n\ pt) = nil$  for all other vertices; and  $(d\ n\ pt\ g)$  returns the weight of  $(path\ n\ pt)$  in  $g$ . It is convenient to use  $NIL$  as "infinity";
4. Repeat until  $ts$  is empty:
  - (a) Choose  $u$  in  $ts$  such that  $(d\ u\ pt\ g)$  is minimal;
  - (b) for each edge from  $u$  to some neighbor  $v$  with weight  $wt$ , if  $(d\ v\ pt\ g) > (d\ u\ pt\ g) + wt$ , then reassign  $(path\ v\ pt)$  to be  $(append\ (path\ u\ pt)\ (list\ v))$ ;
  - (c) Delete  $u$  from  $ts$ .

# Traditional Proof

- When a vertex  $u$  is chosen by step 4(a), the path associated with  $u$  in the path table is the shortest path from the start vertex to  $u$  in the graph
- When a vertex  $u$  is chosen by step 4(a), for any vertex  $v$  chosen after  $u$ , the path associated with  $v$  in the path table is the shortest path from the start vertex to  $v$  through the vertices (i.e. the internal vertices), which are chosen before  $u$

# Mechanical Proof

## ● Main Theorem:

```
(defthm main-theorem
  (implies (and (nodep a g)
                (nodep b g)
                (graphp g))
           (shortest-path a
                          b
                          (dijkstra-shortest-path a b g)
                          g)))
```

## ● Invariant:

```
(defun inv (ts pt g a)
  (let ((fs (comp-set ts (all-nodes g))))
    (and (prop-ts-node a ts fs pt g)
         (prop-fs-node a fs fs pt g)
         (paths-from-s-table a pt g))))
```

# Function details

- **(all-nodes g)** returns all the nodes in the graph g

```
(defun all-nodes (g)
  (cond ((endp g) nil)
        (t (cons-set (caar g)
                     (my-union (strip-cars (cdar g))
                               (all-nodes (cdr g)))))))
```

- **(nodep n g)** returns t iff a is a vertex in the graph g

```
(defun nodep (n g) (mem n (all-nodes g)))
```

- **(graphp g)** returns t iff g is a legal graph:

```
(defun graphp (g)
  (cond ((endp g) (equal g nil))
        ((and (consp (car g))
              (edge-weightsp (cdar g))
              (graphp (cdr g)))
         (t nil)))
```



# Function details

- **(edge-weightsp lst)** returns t iff lst is a legal list of edges:

```
(defun edge-weightsp (lst)
  (cond ((endp lst) (equal lst nil))
        ((and (consp (car lst))
              (rationalp (cдар lst))
              (<= 0 (cдар lst))
              (not (assoc (caar lst) (cdr lst))))
         (edge-weightsp (cdr lst)))
        (t nil)))
```

- **(comp-set ts s)** returns the set deleting ts from s

```
(defun comp-set (ts s)
  (if (endp s) nil
      (if (mem (car s) ts)
          (comp-set ts (cdr s))
          (cons (car s) (comp-set ts (cdr s))))))
```

# Function details

- `(shortest-path a b p g)` returns `t` iff `p` is the shortest path from `a` to `b` in `g`

```
(defun-sk shortest-path (a b p g)
  (forall path (implies (path-from-to path a b g)
                        (shorter p path g))))
```

- `(paths-from-s-table s pt g)` returns `t` iff for any path in `pt`, it is associated with a key vertex `u`, then the path is a path from `s` to `u` in `g`

```
(defun paths-from-s-table (s pt g)
  (if (endp pt) t
      (and (if (not (cdar pt)) t
               (path-from-to (cdar pt) s (caar pt) g))
           (paths-from-s-table s (cdr pt) g))))
```

# Function details

## ● (prop-ts-node a ts fs pt g)

```
(defun prop-ts-node (a ts fs pt g)
  (if (endp ts) t
      (and (shorter-all-inter-path a (car ts)
                                     (path (car ts) pt) fs g)
           (all-but-last-node (path (car ts) pt) fs)
           (prop-ts-node a (cdr ts) fs pt g))))
```

## ● (all-but-last-node p fs)

```
(defun all-but-last-node (p fs)
  (if (endp p) t
      (if (endp (cdr p)) t
          (and (mem (car p) fs)
               (all-but-last-node (cdr p) fs)))))
```

# Function details

## • (shorter-all-inter-path a b p fs g)

```
(defun-sk shorter-all-inter-path (a b p fs g)
  (forall path (implies (and (path-from-to path a b g)
                              (all-but-last-node path fs))
                        (shorter p path g))))
```

## • (prop-fs-node a fs s pt g)

```
(defun prop-fs-node (a fs s pt g)
  (if (endp fs) t
      (and (shortest-path a (car fs) (path (car fs) pt) g)
            (all-but-last-node (path (car fs) pt) s)
            (prop-fs-node a (cdr fs) s pt g))))
```

# Proof sketch

- initially the invariant is correct

```
(defthm inv-0
  (implies (nodep a g)
            (inv (all-nodes g) (list (cons a (list a))) g a)))
```

- the invariant is maintained by the iteration

```
(defthm inv-choose-next
  (implies (and (inv ts pt g a)
                (my-subsetp ts (all-nodes g))
                (graphp g)
                (consp ts)
                (setp ts)
                (nodep a g)
                (equal (path a pt) (list a)))
            (let ((u (choose-next ts pt g)))
              (inv (del u ts)
                   (reassign u (neighbors u g) pt g) g a))))
```

# Proof sketch

- the final form of the invariant is correct

```
(defthm inv-last
  (implies (and (nodep a g)
                (graphp g))
    (inv nil
      (dsp (all-nodes g)
            (list (cons a (list a)))
              g)
            g a)))
```

- main lemma

```
(defthm main-lemma
  (implies (and (inv nil pt g a)
                (nodep b g))
    (shortest-path a b (path b pt) g)))
```

# Prove inv-0

## ● sub-goal 1

```
(implies (mem a (all-nodes g))
          (prop-fs-node a
                        (comp-set (all-nodes g) (all-nodes g))
                        (comp-set (all-nodes g) (all-nodes g))
                        (list (list a a)) g))
```

## ● lemma 1

```
(defthm comp-set-id
  (not (comp-set s s)))
```

# Prove inv-0

## ● sub-goal 2

```
(implies (mem a (all-nodes g))  
         (prop-ts-node a (all-nodes g) nil (list (list a a)) g))
```

## ● lemma 2

```
(defthm prop-path-nil  
  (prop-ts-node a s nil (list (cons a (list a))) g))
```



# Prove inv-choose-next

- lemma 1

```
(defthm paths-from-s-table-reassign
  (implies (and (paths-from-s-table a pt g)
                (graphp g)
                (my-subsetp v-lst (all-nodes g)))
            (paths-from-s-table a (reassign u v-lst pt g) g)))
```

- not hard to prove this lemma

# Prove inv-choose-next

## ● lemma 2

```
(defthm prop-fs-node-choose
  (implies (and (inv ts pt g a)
                (my-subsetp ts (all-nodes g))
                (graphp g)
                (consp ts)
                (setp ts))
            (let ((u (choose-next ts pt g)))
              (prop-fs-node a
                (comp-set (del u ts) (all-nodes g))
                (comp-set (del u ts) (all-nodes g))
                (reassign u (neighbors u g) pt g)
                g))))))
```

# Prove inv-choose-next

## ● lemma 3

```
(defthm prop-ts-node-choose-next
  (implies (and (inv ts pt g a)
                (my-subsetp ts (all-nodes g))
                (setp ts)
                (consp ts)
                (graphp g)
                (nodep a g)
                (equal (path a pt) (list a)))
    (let ((u (choose-next ts pt g)))
      (prop-ts-node a (del u ts)
                    (comp-set (del u ts)
                              (all-nodes g))
                    (reassign u (neighbors u g) pt g)
                    g))))
```

# Prove prop-fs-node-choose-next

- the form of (prop-fs-node a ss ss pt g), has to be generalized
- (comp-set (del u ts) s) VS (cons u (comp-set ts s))
- u is the chosen vertex, which should have the shortest path
- **General lemma**

```
(defthm prop-fs-node-choose-lemma2
  (implies (and (prop-fs-node a fs s pt g)
                (my-subsetp fs (all-nodes g))
                (all-but-last-node (path u pt) s)
                (paths-from-s-table a pt g)
                (nodep u g)
                (graphp g)
                (shortest-path a u (path u pt) g))
            (prop-fs-node a (cons u fs) s
                          (reassign u (neighbors u g) pt g) g)))
```

# Prove prop-fs-node-choose-next

- consider `(comp-set (del u ts) s)` as a subset of `(cons u (comp-set ts s))`

```
(defthm prop-fs-node-choose-lemma3
  (implies (and (my-subsetp s fs)
                (my-subsetp fs (all-nodes g))
                (paths-from-s-table a pt g)
                (prop-fs-node a fs ss pt g))
            (prop-fs-node a s ss pt g)))
```

- compare `(comp-set ts s)` with `(comp-set (del u ts) s)`

```
(defthm prop-fs-node-choose-lemma4
  (implies (and (my-subsetp s ss)
                (prop-fs-node a fs s pt g))
            (prop-fs-node a fs ss pt g)))
```

# Prove prop-fs-node-choose-next

- has to establish  $(\text{shortest-path } a \ u \ (\text{path } u \ pt) \ g)$

```
(defthm choose-next-shortest
  (implies (and (graphp g)
                (consp ts)
                (my-subsetp ts (all-nodes g))
                (inv ts pt g a))
            (shortest-path a (choose-next ts pt g)
                          (path (choose-next ts pt g) pt) g)))
```

- traditional proof: for the chosen vertex  $u$  and any path  $p$  from  $a$  to  $u$  in  $g$ , find the leftmost vertex  $v$ , which is in  $ts$ , in the path  $p$ , then the path associated with  $v$  in  $pt$  is shorter than the partial path from  $a$  to  $v$  in  $p$ , and the partial path is shorter than  $p$ , while  $u$  is chosen before  $v$ , which means the path associated with  $u$  in  $pt$  is shorter than the one associated with  $v$

# Prove choose-next-shortest

- auxiliary function (find-partial-path p s)

```
(defun find-partial-path (p s)
  (if (endp p) nil
      (if (mem (car p) s)
          (cons (car p) (find-partial-path (cdr p) s))
          (list (car p)))))
```

- the partial path is shorter than the original one

```
(defthm partial-path-shorter
  (implies (graphp g)
            (shorter (find-partial-path p s) p g)))
```

# Prove choose-next-shortest

- (find-partial-path p s) returns a path, whose internal vertices are all in s

```
(defthm pathp-partial-path
  (implies (pathp p g)
    (and (path-from-to (find-partial-path p s)
      (car p)
      (car (last (find-partial-path p
        s))))
      g)
    (all-but-last-node (find-partial-path p s) s))))
```



# Prove choose-next-shortest

- the last vertex of (find-partial-path p (comp-set ts (all-nodes g))) is in ts

```
(defthm find-partial-path-last-mem
  (implies (and (mem (car (last p)) ts)
                (pathp p g)
                (my-subsetp ts (all-nodes g))))
    (mem (car
          (last
           (find-partial-path p
                               (comp-set ts
                                         (all-nodes g))))))
          ts)))
```

# Prove choose-next-shortest

- for any vertex  $v$  in  $ts$ , the path associated with the chosen vertex is shorter than the one associated with  $v$

```
(defthm choose-next-shorter-other
  (implies (mem v ts)
            (shorter (path (choose-next ts pt g) pt)
                     (path v pt) g)))
```

- the transitivity of shorter relation

```
(defthm shorter-trans
  (implies (and (shorter p1 p2 g)
                (shorter p2 p3 g))
            (shorter p1 p3 g)))
```

# Prove prop-ts-node-choose-next

```
(defthm prop-ts-node-choose-next
  (implies (and (inv ts pt g a)
                (my-subsetp ts (all-nodes g))
                (setp ts)
                (consp ts)
                (graphp g)
                (nodep a g)
                (equal (path a pt) (list a))))
    (let ((u (choose-next ts pt g)))
      (prop-ts-node a (del u ts)
                    (comp-set (del u ts)
                              (all-nodes g))
                    (reassign u (neighbors u g) pt g)
                    g))))
```

# Prove prop-ts-node-choose-next

- similarly consider (comp-set (del u ts) s) as (cons u (comp-set ts s))

```
(defthm prop-ts-node-lemma3
  (implies (and (paths-from-s-table a pt g)
                (graphp g)
                (nodep a g)
                (equal (path a pt) (list a))
                (prop-fs-node a fs fs pt g)
                (prop-ts-node a ts fs pt g)
                (mem u ts)
                (shortest-path a u (path u pt) g))
            (prop-ts-node a (del u ts) (cons u fs)
                          (reassign u (neighbors u g) pt g) g)))

(defthm prop-ts-node-lemma1
  (implies (and (my-subsetp s fs)
                (my-subsetp fs s)
                (prop-ts-node a ts fs pt g))
            (prop-ts-node a ts s pt g)))
```

# Prove prop-ts-node-lemma3

- 2 sub-goals to prove:
  - for any vertex  $v$  in  $(\text{del } u \text{ ts})$ , the path associated with  $v$  in the reassigned path table is shorter than any path from  $a$  to  $v$  with internal vertices in  $(\text{cons } u \text{ fs})$ , stated by prop-ts-node-lemma2
  - internal vertices of all paths in the reassigned path table are in the set  $(\text{cons } u \text{ fs})$ , stated by prop-ts-node-lemma3-3

```
(defthm prop-ts-node-lemma3-3
  (implies (and (paths-from-s-table a pt g)
                (all-but-last-node (path v pt) fs)
                (all-but-last-node (path u pt) fs))
            (all-but-last-node (path v (reassign u v-lst
                                                pt g))
                              (cons u fs))))
```

# Prove prop-ts-node-lemma2

## • prop-ts-node-lemma2

```
(defthm prop-ts-node-lemma2
  (implies (and (shorter-all-inter-path a v (path v pt) fs g)
                (graphp g)
                (nodep a g)
                (equal (path a pt) (list a))
                (prop-fs-node a fs fs pt g)
                (shortest-path a u (path u pt) g)
                (paths-from-s-table a pt g))
            (shorter-all-inter-path a v
              (path v (reassign u
                              (neighbors u g)
                              pt g))
                (cons u fs) g)))
```

# Prove prop-ts-node-lemma2

## • prop-ts-node-lemma2-3

```
(defthm prop-ts-node-lemma2-3
  (implies (and (shorter-all-inter-path a v (path v pt) fs g)
                (graphp g)
                (prop-fs-node a fs fs pt g)
                (nodep a g)
                (path-from-to p a v g)
                (all-but-last-node p (cons u fs))
                (shortest-path a u (path u pt) g)
                (paths-from-s-table a pt g)
                (equal (path a pt) (list a)))
            (shorter (path v (reassign u (neighbors u g) pt g)
                        p g)))
```

## • two cases to prove

- a and v are identical, easy to prove
- a and v are not equal, by prop-ts-node-lemma2-2

# Prove prop-ts-node-lemma2-2

## ● prop-ts-node-lemma2-2

```
(defthm prop-ts-node-lemma2-2
  (implies (and (shorter-all-inter-path a v (path v pt) fs g)
                (graphp g)
                (prop-fs-node a fs fs pt g)
                (path-from-to p a v g)
                (not (equal a v))
                (shortest-path a u (path u pt) g)
                (all-but-last-node p (cons u fs))
                (paths-from-s-table a pt g))
            (shorter (path v (reassign u (neighbors u g) pt) g)
                     p g)))
```



# Prove prop-ts-node-lemma2-2

- two cases to prove
  - (path u pt) is NIL, (not (all-but-last-node p fs)) happens in the hypotheses. We know (shortest-path a u (path u pt) g) holds and (path u pt) is NIL, therefore there is no path from a to u, then u won't happen in any path, especially in the path p; and we know (all-but-last-node p (cons u fs)) holds, therefore (all-but-last-node p fs) holds.

```
(defthm not-path-implies-path-in-fs
  (implies (and (shortest-path a u (path u pt) g)
                (not (path u pt))
                (graphp g)
                (path-from-to p a v g)
                (all-but-last-node p (cons u fs))))
           (all-but-last-node p fs))
```

- (path u pt) is not NIL, by prop-ts-node-lemma2-1

# Prove prop-ts-node-lemma2-1

## ● prop-ts-node-lemma2-1

```
(defthm prop-ts-node-lemma2-1
  (implies (and (shorter-all-inter-path a v
                                                       (path v pt) fs g)
               (graphp g)
               (prop-fs-node a fs fs pt g)
               (path-from-to p a v g)
               (not (equal a v))
               (path u pt)
               (shortest-path a u (path u pt) g)
               (all-but-last-node p (cons u fs))
               (paths-from-s-table a pt g))
           (shorter (path v (reassign u (neighbors u g) pt g))
                    p g)))
```

# Prove prop-ts-node-lemma2-1

- two cases to prove
  - for the path  $p$  from  $a$  to  $v$ , the vertex neighbored to  $v$  in  $p$  is  $u$ 
    - $(\text{path } u \text{ } pt)$  is the shortest path from  $a$  to  $u$ , so  $(\text{append } (\text{path } u \text{ } pt) (\text{list } v))$  is shorter than  $p$
    - $(\text{path } v \text{ } pt)$  is shorter than  $(\text{append } (\text{path } u \text{ } pt) (\text{list } v))$
    - $(\text{path } v \text{ } (\text{reassign } u \text{ } (\text{neighbors } u \text{ } g) \text{ } pt \text{ } g))$  is shorter than  $(\text{path } v \text{ } pt)$
  - for the path  $p$  from  $a$  to  $v$ , the vertex neighbored to  $v$  in  $p$  isn't  $u$ , we have to define two auxiliary functions

```
(defun find-last-next-path (p)
  (if (or (endp p) (endp (cdr p))) nil
      (cons (car p) (find-last-next-path (cdr p)))))
```

```
(defun last-node (p)
  (car (last (find-last-next-path p))))
```

# Prove prop-ts-node-lemma2-1

- (append (path (last-node p) pt) (list v)) is shorter than (append (find-last-next-path p) (list v)), by last-node-lemma1
- (append (find-last-next-path p) (list v)) is actually the path p, by last-node-lemma2
- (path v pt) is shorter than (append (path (last-node p) pt) (list v)), by shorter-than-append-fs

# Prove prop-ts-node-lemma2-1

## ● last-node-lemma2

```
(defthm last-node-lemma2
  (implies (and (equal (car (last p)) v)
                (pathp p g))
           (equal (append (find-last-next-path p) (list v)) p)))
```

## ● shorter-than-append-fs

```
(defthm shorter-than-append-fs
  (implies (and (shorter-all-inter-path a v (path v pt) s g)
                (prop-fs-node a fs s pt g)
                (my-subsetp fs s)
                (path w pt)
                (paths-from-s-table a pt g)
                (mem w fs))
           (shorter (path v pt)
                    (append (path w pt) (list v)) g)))
```

# Prove last-node-lemma1

- (last-node p) is not equal to u, but still in (cons u fs)
- (path (last-node p) pt) is the shortest path from a to (last-node p), so shorter than (find-last-next-path p)
- shorter-implies-append-shorter

```
(defthm shorter-implies-append-shorter
  (implies (and (shorter p1 p2 g)
                (graphp g)
                (true-listp p1)
                (equal (car (last p1)) (car (last p2)))
                (pathp p2 g))
            (shorter (append p1 (list v))
                    (append p2 (list v)) g)))
```

- to apply shorter-implies-append-shorter, establish (pathp (find-last-next-path p)), by path-from-to-implies-all-path-lemma

# Prove last-node-lemma1

- path-from-to-implies-all-path-lemma

```
(defthm path-from-to-implies-all-path-lemma
  (implies (and (path-from-to p a v g)
                (not (equal a v)))
            (and (pathp (find-last-next-path p) g)
                  (mem v
                       (neighbors
                        (car (last (find-last-next-path p)))
                        g))))))
```

- path p is from a to v, where a isn't equal to v, the length of p is at least 2, by path-length
- the length of path p is at least 2, then the conclusion of the lemma holds, by pathp-find-last-next

# Prove last-node-lemma1

## ● path-length

```
(defthm path-length
  (implies (and (pathp p g)
                (not (equal (car p) (car (last p)))))
           (<= 2 (len p))))
```

## ● pathp-find-last-next

```
(defthm pathp-find-last-next
  (implies (and (pathp p g)
                (<= 2 (len p)))
           (and (pathp (find-last-next-path p) g)
                (mem (car (last p))
                     (neighbors
                      (car (last (find-last-next-path p))
                       g)))))
```



# Prove inv-last

## ● maintain some hypotheses

```
(defthm del-subsetp
  (implies (my-subsetp ts s)
            (my-subsetp (del u ts) s)))
(defthm del-true-listp
  (implies (true-listp ts)
            (true-listp (del u ts))))
(defthm del-noduplicates
  (implies (setp ts)
            (setp (del u ts))))
(defthm path-a-pt-reassign
  (implies (and (paths-from-s-table a pt g)
                (graphp g)
                (nodep a g)
                (equal (path a pt) (list a)))
            (equal (path a (reassign u v-lst pt g)) (list a))))
```

# Conclusion

- Dijkstra's shortest path algorithm
- 122 lemmas and 48 goals proved by hints, within which 27 hints are only the hint of in-theory kind, 6 hints are given on sub-goal level, 19 hints are explicit instantiation of lemmas, 2 hints are explicit induction scheme, and 2 hints are explicit expansion of functions
- follow the traditional proof scheme
- trying to find some common schemes and propose a ACL2 book for further proof in graph algorithms