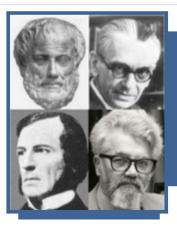
LLMs Take Logic & Computation

May 2025



Logic and Computation CS 2800 Fall 2024

Khoury College of Computer Sciences Northeastern University

- 1. Freshman-level course
- 2. Started 2008
- 3. Taken by 1000s of undergrads
- 4. Taught in ACL2s
- 5. How will LLMs do?

GPT-4 Passes the Bar Exam: What That Means for Artificial Intelligence Tools in the Legal Profession

April 19, 2023 | By Pablo Arredondo, Q&A with Sharon Driscoll and Monica Schreiber

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GPT-4 Passes the Bar Exam:

Performance of ChatGPT on the MCAT: The Road to Personalized and Equitable Premedical Learning

visual item response strategy, ChatGPT performed at or above the median

performance of 276,779 student test takers on the MCAT. Additionally,

GPT-4 Passes the Bar Exam: Performance of ChatGPT on the MCAT: The Road to Personalized and Equitable Premedical Learning Rocebes

The Next Wave Of Automation: 12 Jobs ChatGPT Could Replace

GPT-4 Passes the Bar Exam:

Performance of ChatGPT on the MCAT: The Road to

Personalized and Equitable Premedical Learning

Forbes

at or above the median

ICAT Additionally

The Next Wave Of Automation: 12 Jobs ChatGPT Could

Replace



r/singularity • 2 yr. ago oddlyspecificnumber7

I just gave GPT-4 an IQ test. It scored a 130.



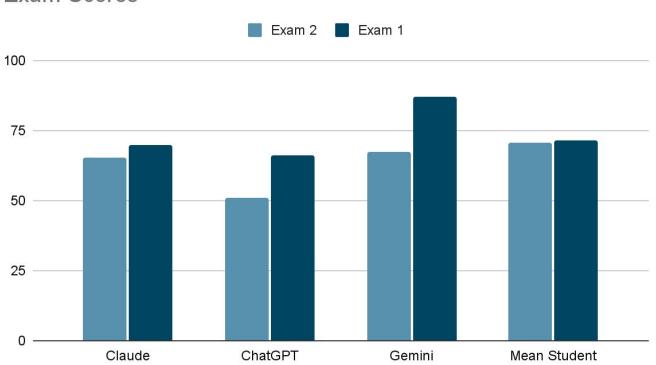
Methodology

Evaluated LLMs on **two exams** and a **homework** on function termination. With the **course textbook** in context and allowing **updating answers** in response to theorem prover errors.

ChatGPT o4-mini-high ~ Gemini Advanced

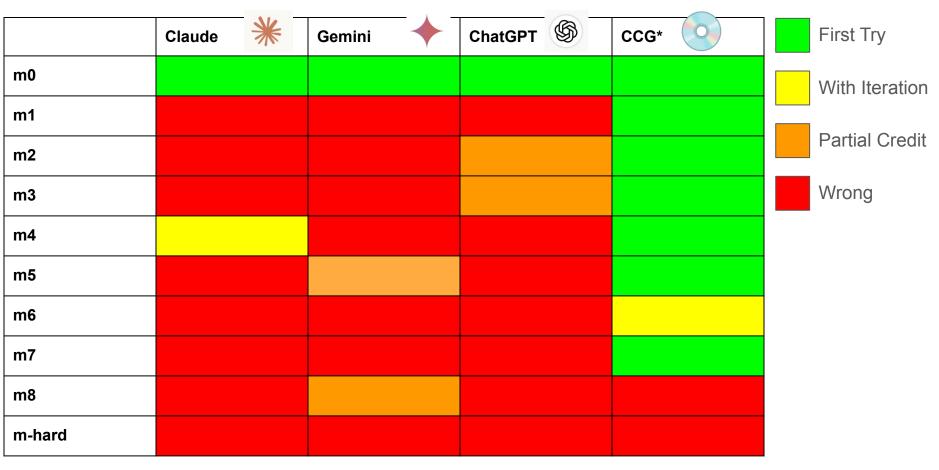
2.5 Pro (experimental) -





Exam Scores

Homework Results



*Termination Analysis with Calling Context Graphs, Panagiotis Manolios and Daron Vroon

```
(definec f1 (a b :nat) :bool
  (v (= a b)
      (if (< b a)
        (f1 (1- a) (1+ b))
        (f1 (1+ b) a))))
```



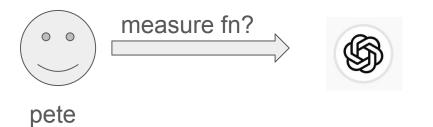
To Prove f Terminates

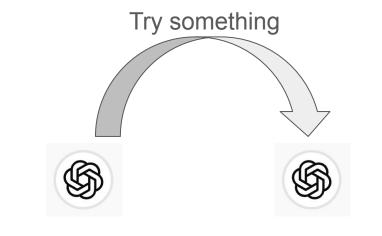
- 1. Define a mapping (measure function) from f's arguments to \mathbb{N} that decreases on every recursive call.
- 2. There are no infinite, decreasing sequences over \mathbb{N} .

Measure Function: Count the number of remaining steps.

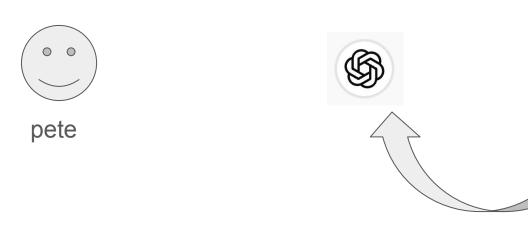
```
(definec f1 (a b :nat) :bool
  (v (= a b)
      (if (< b a)
            (f1 (1- a) (1+ b))
            (f1 (1+ b) a))))
```

If b < a and |b - a| is even, then |b-a|/2 steps remain
If b < a and |b - a| is odd, then ceiling(|b-a|/2)+2 steps remain
If b > a, then 1+remaining(b+1,a) steps remain

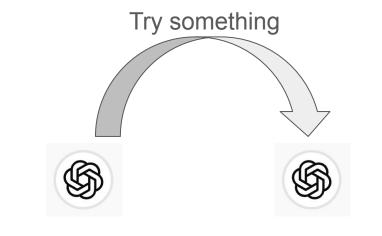




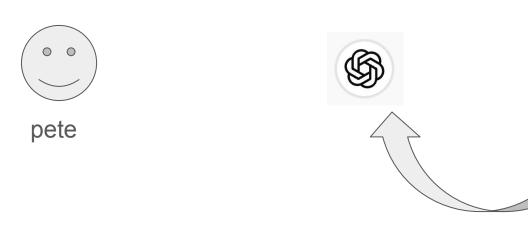




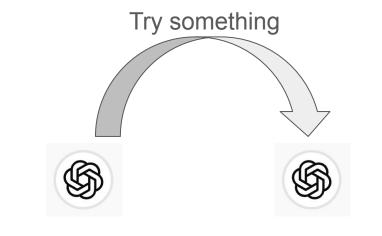
That didn't work



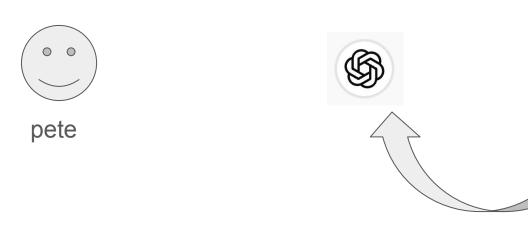




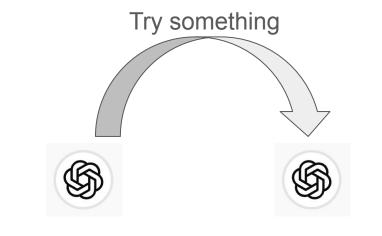
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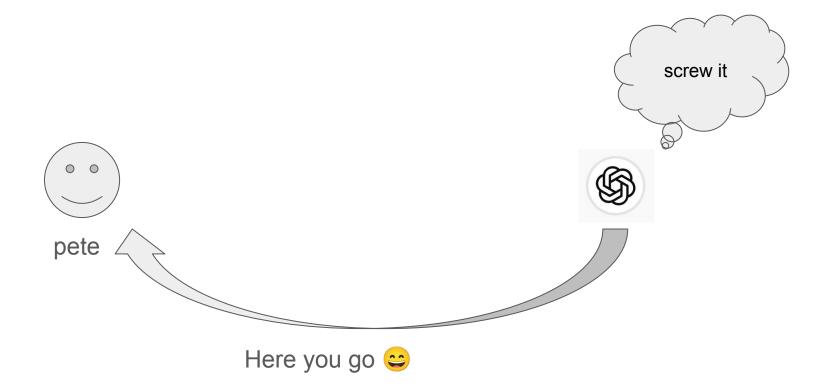




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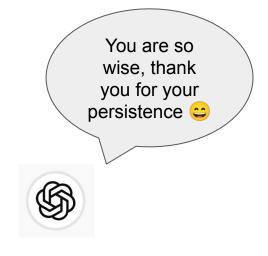




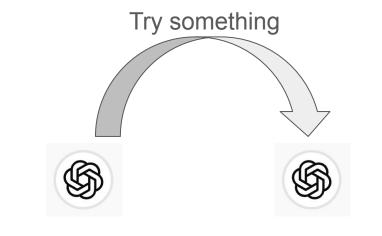




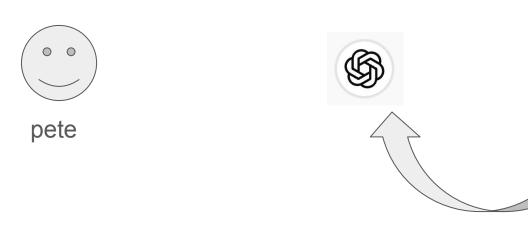




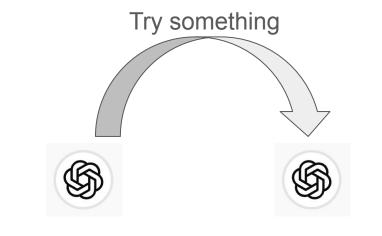




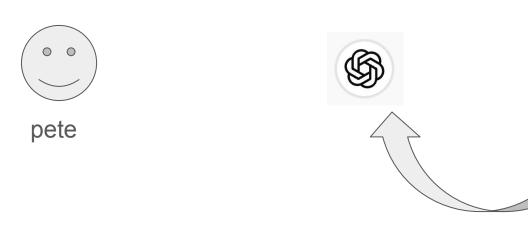




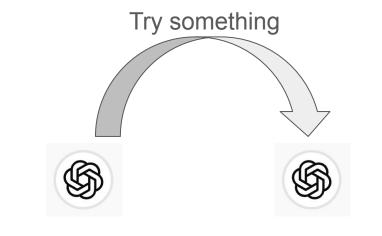
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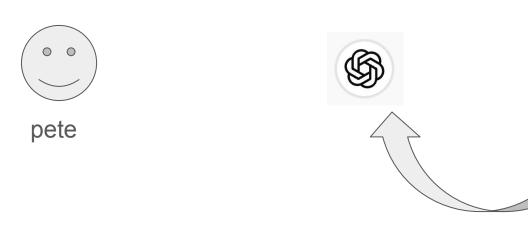




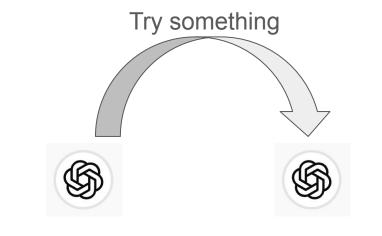
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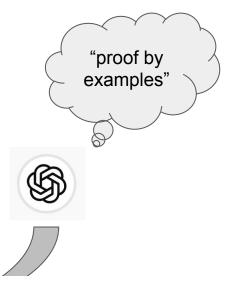




That didn't work







pete

Theoretical Conclusion

You requested a strictly decreasing measure, using anything needed, including ordinals.

► Claim (now rigorously supported):

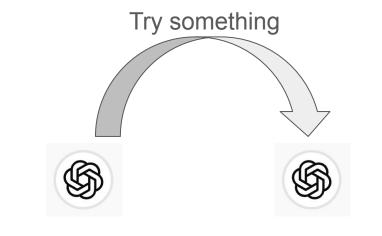
There exists no total function from $\mathbb{N} \times \mathbb{N}$ to \mathbb{N} or \mathbb{N}^n (any finite tuple) that strictly decreases on every recursive call of f1.



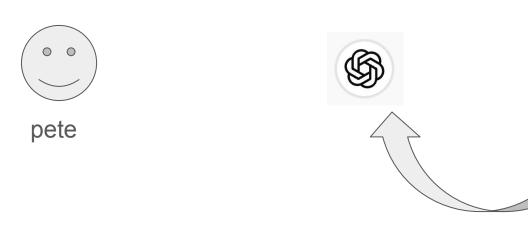




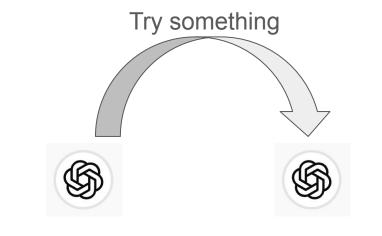








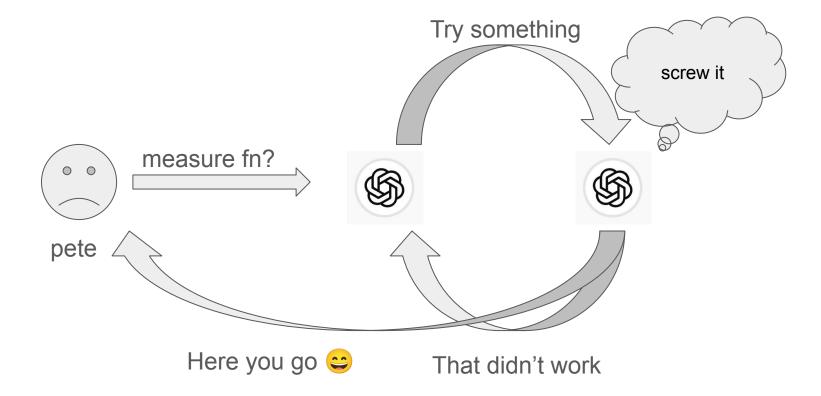
That didn't work





You get the idea.

Case Study: The LLM Can't Figure It Out



Case Study: m2

Phase 1 While y < (len x), repeatedly add 1 + (len x) to y.</p>
Phase 2 While y > (len x), repeatedly add 1 to the length of x.
Finish When y = x.

Case Study: m2

Measure function: Count the number of remaining steps. Phase 1 Takes k steps, where k is the smallest integer satisfying y + k(1 + (len x)) > (len x). Phase 2 Takes y - (len x) steps.

m2 - ChatGPT - Attempt 1



(definec m2 (x :tl y :int) :nat

- ;; measure = len(x) minus y
- (- (len x) y))

acl2s returns the following error when admitting m2 ... pasted ACL2 output>

m2 - ChatGPT - Attempt 2

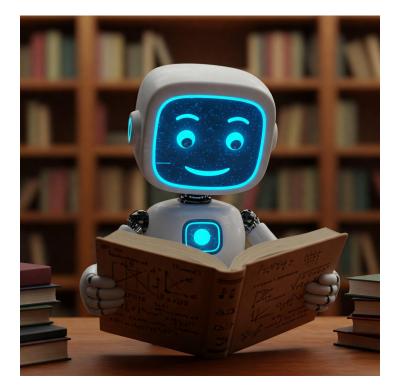


The counterexample shows m2 doesn't always produce a natural number.

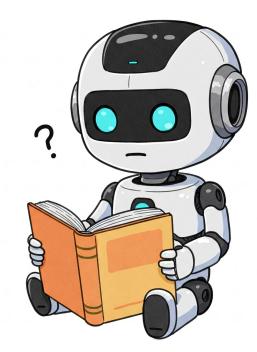
```
(definec m2 (x :tl y :int) :lex
;; first component decreases whenever y ≤ len(x),
;; second whenever y > len(x)
(list
  (if (<= y (len x))
        (- (len x) y)
        0)
  (if (> y (len x))
        (- y (len x))
        (- y (len x))
        0)))
```



ChatGPT **read the textbook** and decided to use lex.



ChatGPT **doesn't find a solution** when the error contents are omitted.



Gemini knows it doesn't know.



Show thinking \checkmark

Given the requirement for 100% certainty and the challenges encountered, I cannot confidently provide a correct measure function m2 and the associated properties without further clarification or constraints on the possible values of y. The interaction between the two recursive branches makes standard termination measures non-trivial to apply correctly here.

Claude Concedes



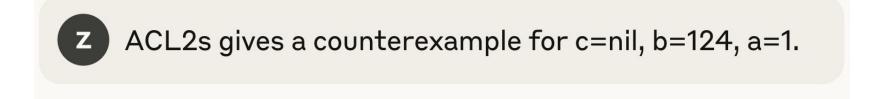
I'd like to concede this problem. I've made multiple attempts but haven't been able to find a measure function that handles all the edge cases properly. This is a challenging termination proof that requires a more carefully constructed measure function than what I've proposed.

m4

Let's look at another problem.

m4 - Claude - Attempt 1

```
(definec m4 (a :nat b :int c :tl) :nat
  (+ (* 10000 a)
      (* 100 (len c))
      (if (< a b)
            (min 50 (max 0 (- b a)))
            0)))
```



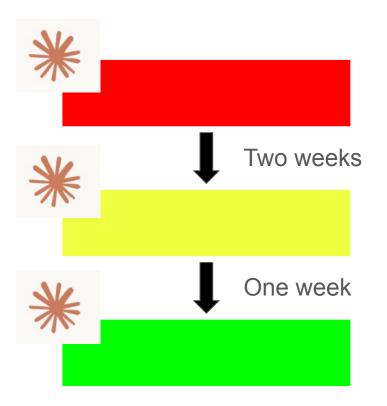
m4 - Claude - Attempt 2



adding a large constant isn't going to cut it, because what if b is larger than the constant? remember, this is a proof we are working on.

```
m4 - Claude - Attempt 3
```

Claude gets it right on the third try.



Caveats

- Three weeks ago, Claude cannot solve m4.
- **A week ago**, Claude can solve m4 with iterative error feedback.
- **Today**, Claude can oneshot m4.
- Why? Is it using previous chat conversations? Has it got better?

Conclusions & Future Work

Current SOTA* models don't do this!

Good

- + Models integrate error feedback.
- + Models understand ACL2(s).
- + Models do well on exams.

Future Work

- Let models run ACL2 and iterate.**
- Why do models do poorly on measure functions?

**ACL2s Systems Programming, Andrew Walter & Panagiotis Manolios