

A Formalization of Elementary Linear Algebra: Parts I and II

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OUTLINE

Part I:

- ▶ Matrix Algebra
- ▶ Determinants

Part II:

- ▶ Row reduction
- ▶ Invertibility
- ▶ Systems of Linear Equations

Part III: (2026?)

- ▶ Vector Spaces
- ▶ Linear Transformations
- ▶ Eigenvectors and Diagonalization

OUTLINE

Part I: Matrices over a Commutative Ring

- ▶ Matrix Algebra
- ▶ Determinants

Part II: Matrices over a Field

- ▶ Row reduction
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THE DIRECTORY `books/projects/linear/`

- ▶ `ring`: Axioms of commutative ring R , lists of elements of R
- ▶ `rmat`: Matrices over R , algebraic operations, transpose
- ▶ `rdet`: Determinant of a matrix over R
- ▶ `field`: Axioms of field F , lists of elements of F
- ▶ `fmat`: Matrices over F , algebraic operations, transpose
- ▶ `fdet`: Determinant of a matrix over F
- ▶ `reduction`: Row reduction, inverse, linear equations
- ▶ `rational`: Matrices over $\mathbb{Q}$

THE COMMUTATIVE RING R

```
(encapsulate (((rp *) => *)                               ;element recognizer
              ((r+ * *) => *) ((r* * *) => *)             ;operations
              ((r0) => *) ((r1) => *)                     ;identities
              ((r- *) => *))                               ;additive inverse
  (local (defun rp (x) (rationalp x)))
  (local (defun r+ (x y) (+ x y)))
  (local (defun r* (x y) (* x y)))
  (local (defun r0 () 0))
  (local (defun r1 () 1))
  (local (defun r- (x) (- x)))
  ;; Closure:
  (defthm r+closed (implies (and (rp x) (rp y)) (rp (r+ x y))))
  (defthm r*closed (implies (and (rp x) (rp y)) (rp (r* x y))))
  ;; Commutativity
  ...
)
```

Informally, we refer to the ring R characterized by these axioms.

MATRICES OVER R

An $m \times n$ matrix over R is a list of m rlists of length n :

```
(defun rmatp (a m n)
  (if (zp m)
      (null a)
      (and (consp a)
            (rlistnp (car a) n)
            (rmatp (cdr a) (1- m) n))))

(defun row (i a) (nth i a))

(defun col (j a)
  (if (consp a)
      (cons (nth j (car a)) (col j (cdr a)))
      ()))

(defun entry (i j a) (nth j (nth i a)))
```

MATRIX OPERATIONS

```
(defun rmat-add (a b)
  (if (consp a)
      (cons (rlist-add (car a) (car b))
            (rmat-add (cdr a) (cdr b)))
      ()))
```

```
(defun rmat-scalar-mul (c a)
  (if (consp a)
      (cons (rlist-scalar-mul c (car a))
            (rmat-scalar-mul c (cdr a)))
      ()))
```

```
(defun replace-row (a k r)
  (if (zp k)
      (cons r (cdr a))
      (cons (car a) (replace-row (cdr a) (1- k) r))))
```

```
(defun rmat-sum (a)
  (if (consp a)
      (r+ (rlist-sum (car a)) (rmat-sum (cdr a)))
      (r0)))
```

TRANPOSE OF A MATRIX

The *transpose* of a matrix is the list of its columns:

```
(defun transpose-mat-aux (a j n)
  (if (and (natp j) (natp n) (< j n))
      (cons (col j a) (transpose-mat-aux a (1+ j) n))
      ()))

(defun transpose-mat (a) (transpose-mat-aux a 0 (len (car a))))

(defthm transpose-rmat-entry
  (implies (and (posp m) (posp n) (rmatp a m n)
                (natp j) (< j n) (natp i) (< i m))
    (equal (entry j i (transpose-mat a))
           (entry i j a))))

(defthmd sum-rmat-transpose
  (implies (and (natp m) (natp n) (rmatp a m n))
    (equal (rmat-sum (transpose-mat a))
           (rmat-sum a))))
```

MATRIX MULTIPLICATION

The product of an $m \times n$ matrix and an $n \times p$ matrix is an $m \times p$ matrix:

```
(defund rmat* (a b)
  (if (consp a)
      (cons (rdot-list (car a) (transpose-mat b))
            (rmat* (cdr a) b))
      ()))

(defthm rmatp-rmat*
  (implies (and (rmatp a m n) (rmatp b n p)
                (posp m) (posp n) (posp p))
            (rmatp (rmat* a b) m p)))

(defthmd rmat*-entry
  (implies (and (posp m) (posp n) (posp p)
                (rmatp a m n) (rmatp b n p)
                (natp i) (< i m) (natp j) (< j p))
            (equal (entry i j (rmat* a b))
                    (rdot (row i a) (col j b)))))
```

IDENTITY MATRIX

```
(defun runit (i n)
  (if (zp n) ()
      (if (zp i) (cons (r1) (rlistn0 (1- n)))
          (cons (r0) (runit (1- i) (1- n))))))

(defun id-rmat-aux (i n)
  (if (and (natp i) (natp n) (< i n))
      (cons (runit i n) (id-rmat-aux (1+ i) n))
      ()))
(defun id-rmat (n) (id-rmat-aux 0 n))

(defthmd id-rmat-right
  (implies (and (posp m) (posp n) (rmatp a m n))
            (equal (rmat* a (id-rmat n)) a)))

(defthmd id-rmat-left
  (implies (and (posp m) (posp n) (rmatp a m n))
            (equal (rmat* (id-rmat m) a) a)))
```

ASSOCIATIVITY OF MATRIX MULTIPLICATION

If a , b , and c are matrices of dimensions $m \times n$, $n \times p$, and $p \times q$, then

$$((ab)c)_{ij} = \sum_{s=0}^{p-1} \sum_{r=0}^{n-1} a_{ir} b_{rs} c_{sj} = \sum_{r=0}^{n-1} \sum_{s=0}^{p-1} a_{ir} b_{rs} c_{sj} = (a(bc))_{ij}.$$

```
(defthmd rmat*-assoc-entry
  (implies (and (posp m) (posp n) (posp p) (posp q)
                (rmatp a m n) (rmatp b n p) (rmatp c p q)
                (natp i) (< i m) (natp j) (< j q))
    (equal (entry i j (rmat* a (rmat* b c)))
           (entry i j (rmat* (rmat* a b) c))))

(defthmd rmat*-assoc
  (implies (and (rmatp a m n) (rmatp b n p) (rmatp c p q)
                (posp m) (posp n) (posp p) (posp q))
    (equal (rmat* a (rmat* b c))
           (rmat* (rmat* a b) c))))
```

DEFINITION OF DETERMINANT

Each permutation p in the symmetric group $(\text{sym } n)$ contributes a term to the determinant of an $n \times n$ matrix a :

```
(defun rdet-prod (a p n)
  (if (zp n)
      (r1)
      (r* (rdet-prod a p (1- n))
           (entry (1- n) (nth (1- n) p) a))))

(defun rdet-term (a p n)
  (if (even-perm-p p n)
      (rdet-prod a p n)
      (r- (rdet-prod a p n))))

(defun rdet-sum (a l n)
  (if (consp l)
      (r+ (rdet-term a (car l) n) (rdet-sum a (cdr l) n))
      (r0)))

(defun rdet (a n) (rdet-sum a (slist n) n))
```

PROPERTIES OF DETERMINANT

The determinant is characterized by three properties:

```
(defthm rdet-id-rmat
  (implies (posp n) (equal (rdet (id-rmat n) n) (r1))))

(defthmd rdet-alternating
  (implies (and (rmatp a n n) (posp n)
                (natp i) (< i n) (natp j) (< j n) (not (= i j))
                (= (row i a) (row j a)))
    (equal (rdet a n) (r0))))

(defthm rdet-n-linear
  (implies (and (rmatp a n n) (posp n) (natp i) (< i n)
                (rlistnp x n) (rlistnp y n) (rp c))
    (equal (rdet (replace-row
                  a i (rlist-add (rlist-scalar-mul c x) y))
            n)
      (r+ (r* c (rdet (replace-row a i x) n))
        (rdet (replace-row a i y) n)))))
```

UNIQUENESS OF DETERMINANT

The determinant is uniquely determined by those properties:

```
(encapsulate (((rdet0 * *) => *))
  (local (defun rdet0 (a n) (rdet a n)))
  (defthm rp-rdet0
    (implies (and (rmatp a n n) (posp n)) (rp (rdet0 a n))))
  (defthmd rdet0-n-linear
    ... )
  (defthmd rdet0-adjacent-equal
    (implies (and (rmatp a n n) (posp n)
                  (natp i) (< i (1- n))
                  (= (row i a) (row (1+ i) a)))
      (equal (rdet0 a n) (r0)))))

(defthmd rdet-unique
  (implies (rmatp a n n)
    (equal (rdet0 a n)
      (r* (rdet a n) (rdet0 (id-rmat n) n)))))
```

MULTIPLICATIVITY OF DETERMINANT

Multiplicativity is proved by functional instantiation:

```
(defun rdet-rmat* (a b n)
  (rdet (rmat* a b) n))

(defthm rdet-rmat*-rewrite
  (implies (and (rmatp a n n) (rmatp b n n) (posp n))
    (equal (rdet-rmat* a b n)
      (r* (rdet a n) (rdet-rmat* (id-rmat n) b n))))
  :hints
  (("Goal" :use ((functional-instance rdet-unique
    (rdet0 (lambda (a n) (rdet-rmat* a b n)))))))

(defthm rdet-multiplicative
  (implies (and (rmatp a n n) (rmatp b n n) (posp n))
    (equal (rdet (rmat* a b) n)
      (r* (rdet a n) (rdet b n))))
```

COFACTOR EXPANSION OF DETERMINANT

Expansion with respect to column j :

```
(defund minor (i j a)
  (delete-col j (delete-row i a)))

(defund rdet-cofactor (i j a n)
  (if (evenp (+ i j))
      (rdet (minor i j a) (1- n))
      (r- (rdet (minor i j a) (1- n)))))

(defun expand-rdet-col-aux (a i j n)
  (if (zp i)
      (r0)
      (r+ (r* (entry (1- i) j a)
                (rdet-cofactor (1- i) j a n))
            (expand-rdet-col-aux a (1- i) j n))))

(defund expand-rdet-col (a j n)
  (expand-rdet-col-aux a n j n))
```

CORRECTNESS OF COFACTOR EXPANSION

For fixed k , $(\text{expand-rdet-col } a \ k \ n)$ may be shown to satisfy the constraints on $(\text{rdet0 } a \ n)$. By functional instantiation of rdet-unique ,

```
(defthmd expand-rdet-col-val
  (implies (and (rmatp a n n) (posp n) (> n 1)
                (natp k) (< k n))
    (equal (expand-rdet-col a k n)
      (r* (rdet a n)
        (expand-rdet-col (id-rmat n) k n)))))
```

But $(\text{expand-rdet-col } (\text{id-rmat } n) \ k \ n) = 1$. Thus,

```
(defthmd expand-rdet-col-rdet
  (implies (and (rmatp a n n) (posp n) (> n 1)
                (natp k) (< k n))
    (equal (expand-rdet-col a k n)
      (rdet a n))))
```

RECURSIVE FORMULATION OF DETERMINANT

```
(mutual-recursion
  (defund rdet-rec-cofactor (j a n)
    (if (zp n) ()
        (if (evenp j)
            (rdet-rec (minor 0 j a) (1- n))
            (r- (rdet-rec (minor 0 j a) (1- n))))))
  (defun expand-rdet-rec-aux (a j n)
    (if (zp j) (r0)
        (r+ (r* (entry 0 (1- j) a)
                 (rdet-rec-cofactor (1- j) a n))
              (expand-rdet-rec-aux a (1- j) n))))
  (defund expand-rdet-rec (a n)
    (expand-rdet-rec-aux a n n))
  (defun rdet-rec (a n)
    (if (zp n) (r0)
        (if (= n 1) (entry 0 0 a)
            (expand-rdet-rec a n))))

  (defthmd rdet-rec-rdet
    (implies (and (rmatp a n n) (posp n))
              (equal (rdet-rec a n) (rdet a n))))
```

CLASSICAL ADJOINT

The *adjoint* of an $n \times n$ matrix is defined to satisfy

```
(defthmd adjoint-rmat-entry
  (implies (and (rmatp a n n) (natp n) (> n 1)
                (natp i) (< i n) (natp j) (< j n))
    (and (rmatp (adjoint-rmat a n) n n)
          (equal (entry i j (adjoint-rmat a n))
                  (rdet-cofactor j i a n)))))
```

A consequence of `expand-rdet-col-rdet`:

```
(defthmd rmat*-adjoint-rmat
  (implies (and (rmatp a n n) (natp n) (> n 1))
    (equal (rmat* a (adjoint-rmat a n))
            (rmat-scalar-mul (rdet a n) (id-rmat n)))))
```

THE FIELD F

```
(encapsulate ((fp *) => *) ;element recognizer
              ((f+ * *) => *) ((f* * *) => *) ;operations
              ((f0) => *) ((f1) => *) ;identities
              ((f- *) => *) ((f/ *) => *)) ;inverses
(local (defun fp (x) (rationalp x)))
(local (defun f+ (x y) (+ x y)))
(local (defun f* (x y) (* x y)))
(local (defun f0 () 0))
(local (defun f1 () 1))
(local (defun f- (x) (- x)))
(local (defun f/ (x) (/ x)))
;; Closure:
...
;; Multiplicative inverse:
(defthm fpf/
  (implies (and (fp x) (not (equal x (f0)))))
    (fp (f/ x))))
(defthm f*inv
  (implies (and (fp x) (not (equal x (f0)))))
    (equal (f* x (f/ x)) (f1)))))
```

REDUCED ROW-ECHELON FORM

A *reduced row-echelon* matrix, recognized by the predicate `row-echelon-p`, satisfies the following:

- (1) Every all-zero row is preceded by every nonzero row;
- (2) The first nonzero entry of each nonzero row is 1, and every other entry in the same column is 0;
- (3) The column of the leading 1 in the i th nonzero row is an increasing function of i .

ELEMENTARY ROW OPERATIONS

(1) Multiply row k by a scalar c :

```
(defund ero1 (a c k)
  (replace-row a k (flist-scalar-mul c (nth k a))))
```

(2) Add a scalar multiple of row j to row k , where $j \neq k$:

```
(defund ero2 (a c j k)
  (replace-row a k
    (flist-add (flist-scalar-mul c (nth j a))
      (nth k a))))
```

(3) Interchange rows j and k , where $j \neq k$:

```
(defund ero3 (a j k)
  (replace-row (replace-row a k (nth j a)) j (nth k a)))
```

CONVERSION TO REDUCED ROW-ECHELON FORM

The function `row-reduce` applies a sequence of elementary row operations to convert a matrix to RRE form:

```
(defthmd row-echelon-p-row-reduce
  (implies (and (natp m) (natp n) (fmatp a m n))
    (row-echelon-p (row-reduce a))))
```

Row rank of a matrix:

```
(defun row-rank (a) (num-nonzero-rows (row-reduce a)))
```

ENCODING ELEMENTARY ROW OPERATIONS

Each elementary row op is encoded as a list, e.g., multiplication of row 3 by -2 is represented by $(1 \ -2 \ 3)$.

`(apply-row-op op a)` applies an encoded op to matrix a

`(apply-row-ops ops a)` applies a list ops to matrix a

`(row-reduce-ops a)` is the list of ops required to reduce a

```
(defthmd apply-row-reduce-ops
  (implies (and (fmatp a m n) (posp m) (posp n))
    (equal (apply-row-ops (row-reduce-ops a) a)
      (row-reduce a))))
```

ROW REDUCTION AS MATRIX MULTIPLICATION

The *elementary matrix* corresponding to a row op:

```
(defund elem-mat (op m) (apply-row-op op (id-fmat m)))
```

Row reduction is equivalent to multiplication by a product of elementary matrices:

```
(defund row-ops-mat (ops m)
  (if (consp ops)
      (fmat* (row-ops-mat (cdr ops) m) (elem-mat (car ops) m))
      (id-fmat m)))
```

```
(defund row-reduce-mat (a)
  (row-ops-mat (row-reduce-ops a) (len a)))
```

```
(defthmd row-ops-mat-row-reduce
  (implies (and (fmatp a m n) (posp m) (posp n))
            (equal (fmat* (row-reduce-mat a) a)
                    (row-reduce a))))
```

INVERTIBILITY OF AN $n \times n$ MATRIX

```
(defund invertiblep (a n) (= (row-rank a) n))

(defund inverse-mat (a) (row-reduce-mat a))

(defthmd invertiblep-sufficient
  (implies (and (fmatp a n n) (posp n) (invertiblep a n))
    (let ((p (inverse-mat a)))
      (and (equal (fmat* a p) (id-fmat n))
        (equal (fmat* p a) (id-fmat n))))))

(defthmd invertiblep-necessary
  (implies (and (fmatp a n n) (fmatp b n n) (posp n)
    (= (fmat* a b) (id-fmat n)))
    (invertiblep a n)))

(defthmd invertiblep-fdet-not-zero
  (implies (and (fmatp a n n) (posp n))
    (iff (invertiblep a n)
      (not (equal (fdet a n) (f0))))))
```

SIMULTANEOUS SYSTEMS OF LINEAR EQUATIONS

Given an $m \times n$ matrix a with (entry $i \ j \ a$) $= a_{i,j}$
and a vector $b = (b_0 \dots b_{m-1})$, find a vector $x = (x_0 \dots x_{n-1})$
such that for $0 \leq i < m$,

$$a_{i,0}x_0 + \dots + a_{i,n-1}x_{n-1} = b_i.$$

As a matrix equation,

```
(defund solutionp (x a b)
  (equal (fmat* a (col-mat x))
    (col-mat b)))
```

SOLVABILITY

Let

```
bc = (col-mat b)
xc = (col-mat x)
p = (row-reduce-mat a)
ar = (fmat* p a)
br = (fmat* p bc).
```

Then $(\text{fmat} * a \text{ xc}) = bc \Leftrightarrow (\text{fmat} * ar \text{ xc}) = br.$

Let $q = (\text{row-rank } a).$

A solution exists $\Leftrightarrow$ the last $m - q$ entries of br are 0.

SOLVABILITY

Algorithmic definition of solvability:

```
(defun find-nonzero (br q m)
  (if (and (natp q) (natp m) (< q m))
      (if (= (entry (1- m) 0 br) (f0))
          (find-nonzero br q (1- m))
          (1- m))
      ()))

(defun solvablep (a b)
  (null (find-nonzero (fmat* (row-reduce-mat a) (col-mat b))
                      (row-rank a)
                      (len a)))))
```

UNIQUELY SOLVABLE CASE

If $(\text{solvablep } a \ b)$ and $(\text{row-rank } a) = n$,
then the solution is unique:

```
(defthmd linear-equations-unique-solution-case
  (let* ((br (fmat* (row-reduce-mat a) (col-mat b)))
        (bq (first-rows n br)))
    (implies (and (fmatp a m n) (posp m) (posp n)
                  (flistnp b m) (flistnp x n)
                  (solvablep a b) (= (row-rank a) n))
              (iff (solutionp x a b)
                    (equal x (col 0 bq))))))
```

CRAMER'S RULE

An alternative method of solution in the case of an invertible square coefficient matrix:

```
(defthmd cramer
  (implies (and (fmatp a n n) (natp n) (> n 1)
                (invertiblep a n)
                (flistnp b n) (flistnp x n)
                (solutionp x a b)
                (natp i) (< i n))
    (equal (nth i x)
            (f* (fdet (replace-col a i b) n)
                (f/ (fdet a n))))))
```

FUTURE WORK

Continuation of linear algebra:

- ▶ Vector spaces and linear transformations
- ▶ Polynomials and factorization
- ▶ Eigenvectors and diagonalization

Algebraic number theory:

- ▶ Algebraic extensions and number fields
- ▶ Galois groups
- ▶ Quadratic fields and Hilbert's 10th problem