CS311H

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Good Morning, Colleagues



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Are there any questions?





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- Until Thanksgiving: Big O, Master Theorem, Proving program correctness
 - This week may have been review consider it vacation after exam
- No discussion Wed. before Thanksgiving



- How does O, Ω, Θ relate to limits?
- f(x) being of "order" g(x) is a way of saying f(x) is $\Theta(g(x))$



• $f(n) = 4n^2 - 5n + 3$ is $O(n^2)$.



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4. $< (x^2 + x^2)/x$ {because $x > 1$ }
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7. $= 2|x|$
8. Therefore $C = 2$ and $\forall x > K$, $|(x^2 + 1)/(x + 1)| \le C|x|$.







Proof by Contradiction:



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Proof by Contradiction: Suppose n³ is O(7n²) Then there are C and k such that $n^3 \le C7n^2$, $\forall n \ge k$ But $n^3 \le C7n^2$ implies that $n \le 7C$



Proof by Contradiction: Suppose n³ is O(7n²) Then there are C and k such that $n^3 \leq C7n^2$, $\forall n \geq k$ But $n^3 \leq C7n^2$ implies that $n \leq 7C$ But this fails for values of n that are greater than 7C. So we have a contradiction.



• Suppose f(x) is O(g(x)) and g(x) is O(h(x)). Prove f(x) is O(h(x)).



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• Prove if f(x) is O(g(x)), then g(x) is $\Omega(f(x))$



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• Suppose f(x) is O(g(x)) and g(x) is O(h(x)). Prove f(x) is O(h(x)). 1. f(x) is $O(g(x)) \Rightarrow \forall x > K_1, |f(x)| \le C_1|g(x)|$ for some K_1, C_1 .



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 - 3. Let $K = max(K_1, K_2)$ and $C = C_1C_2$.



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 - 3. Let $K = max(K_1, K_2)$ and $C = C_1C_2$.
 - 4. Then $\forall x > K$, $|f(x)| \le C_1|g(x)| \le C_1(C_2|h(x)|) = C|h(x)|$.



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 - 4. Then $\forall x > K$, $|f(x)| \le C_1 |g(x)| \le C_1 (C_2 |h(x)|) = C |h(x)|$. 5. Therefore f(x) is O(h(x)).
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 - 4. Then $\forall x > K$, $|f(x)| \le C_1 |g(x)| \le C_1 (C_2 |h(x)|) = C |h(x)|$. 5. Therefore f(x) is O(h(x)).
- Prove if f(x) is O(g(x)), then g(x) is $\Omega(f(x))$
 - (Try on piazza)



• Consider $f(n) = n(\sin n)$



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- Show that f(n) is $\Omega(\sin n)$.
- Is $f(n) O(\sin n)$?
- Show that f(n) is neither O(1) nor $\Omega(1)$
- Find a function g(n) such that f(n) is $\Theta(g(n))$.



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- Show that f(n) is O(n).
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- Is $f(n) O(\sin n)$?
- Show that f(n) is neither O(1) nor $\Omega(1)$
- Find a function g(n) such that f(n) is $\Theta(g(n))$.
 - $-g(n)=n(\sin n)$
 - Every function is Θ of itself!