CS311H

Prof: Peter Stone

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Good Morning, Colleagues



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Are there any questions?





• Extra big-O problems on piazza





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- Midterm stats:





- Extra big-O problems on piazza
- Midterm stats:
 - Mean: 46 (out of 60)
 - Median: 49.2
 - High score: 64





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- log:remember that $\log_b x = \frac{\log_d x}{\log_d b}$
- Examples of usage, coming up with recurrences
- Why is Master theorem true? (intuition, proof)





• For each function f(x) find a function g(x) such that f(x) is $\theta(g(x)).$





For each function f(x) find a function g(x) such that f(x) is θ(g(x)).
1. f(x) = 10
2. f(x) = 3x + 7
3. f(x) = x² + x + 1
4. f(x) = 5 log x
5. f(x) = floor(x)
6. f(x) = ceiling(x/2)





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Big Theta

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• Prove if f(x) is O(g(x)), then g(x) is $\Omega(f(x))$



Peter Stone

f increasing function with $f(n) = af(n/b) + cn^d$, $a \ge 1, b \in \mathbb{N}, c, d > 0$.

• Show that if $a = b^d$ and n a power of b, then $f(n) = f(1)n^d + cn^d \log_b n$.



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- Show that if $a = b^d$ and n a power of b, then $f(n) = f(1)n^d + cn^d \log_b n$.
- Show that if $a = b^d$, then f(n) is $O(n^d \log n)$.



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Answer: $7 \log n = O(n) \Rightarrow 7 \log n \le Cn$
So, $T(n) = 8T(n/2) + 6n + 7 \log n \le 8T(n/2) + 6n + Cn =$
 $8T(n/2) + (6 + C)n$
Use Master Theorem: $8 > 2^1 \Rightarrow T(n) = O(n^{\log_2 8}) = O(n^3)$