CS313H Logic, Sets, and Functions: Honors Fall 2012

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Department of Computer Science The University of Texas at Austin

Good Morning, Colleagues



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Are there any questions?





• Office hours delayed





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- Grades





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- Extra big-O problems on piazza





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- Extra big-O problems on piazza
- Class Tuesday next week important
- No discussion Wed. before Thanksgiving



Questions

• log:





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- Examples of usage, coming up with recurrences
- Is Master theorem tight?
- Can d be 0?
- Why is Master theorem true? (intuition, proof)





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For each function f(x) find a function g(x) such that f(x) is θ(g(x)).
1. f(x) = 10
2. f(x) = 3x + 7
3. f(x) = x² + x + 1
4. f(x) = 5 log x
5. f(x) = floor(x)
6. f(x) = ceiling(x/2)





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Big Theta

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• Prove if f(x) is O(g(x)), then g(x) is $\Omega(f(x))$



Peter Stone

f increasing function with $f(n) = af(n/b) + cn^d$, $a \ge 1, b \in \mathbb{N}, c, d > 0$.

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- Show that if $a = b^d$ and n a power of b, then $f(n) = f(1)n^d + cn^d \log_b n$.
- Show that if $a = b^d$, then f(n) is $O(n^d \log n)$.



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 $8T(n/2) + (6 + C)n$
Use Master Theorem: $8 > 2^1 \Rightarrow T(n) = O(n^{\log_2 8}) = O(n^3)$