CS344M Autonomous Multiagent Systems Spring 2008

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- Auction time constraints (sniping)?



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- Nash eq. in voting systems?
- Bargaining why take arbitraily little money?



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- Voting: maximize social good
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What about Clarke tax algorithm?



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Is computational complexity of strategizing important?



Arrow's Theorem

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Criterion of independence of irrelevant alternatives. If one set of preference ballots would lead to an an overall ranking of alternative X above alternative Y and if some preference ballots are changed without changing the relative rank of X and Y, then the method should still rank X above Y.



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Non-dictatorship. There should not be one specific voter whose preference ballot is always adopted.



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Citizen Sovereignty. Any ranking possible

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Independence of irrelevant alternatives. Removing or adding a non-winner doesn't change winner



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Not all possible!



 Strategy proof under weaker irrelevant alternatives criterion



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- Smith set: smallest set of candidates such that each candidate in the set preferred over each candidate not in the set
- Every candidate in the Smith set is relevant



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• Does that solve everything? What about cycles?



small market, both can come out favorably



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 - If rejects, both get nothing
- Another version
 - One person makes an offer
 - Other accepts, rejects, or counters
 - If counters, \$.05 lost
 - Game ends with an accept or reject



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Unique solution that satisfies:



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Invariance: only preference orders matter
Anonymity: no discrimination
Pareto efficiency: if one does better, other does worse
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Maximize $u_1(o) * u_2(o)$



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Consumers: utilities, endowments Producers: production possibility sets Variables: prices on goods Equilibrium: allocation (prices) such that consumers maximize preferences, producers maximize profits

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 - Braess' paradox



• Contract nets: task allocation among agents



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 - Contingencies
 - Leveled commitment (price)



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- Contract nets: task allocation among agents
 - Contingencies
 - Leveled commitment (price)
- Coalitions
 - Formation
 - Optimization within
 - Payoff division



Task allocation among agents



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 - Contingency (future events)
 - Leveled commitment (price)
 - What are some of the tradeoffs?



Contingency problems:



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Leveled commitment problems:

- 1. Breacher's gain may be smaller than victim's loss
- 2. May decommit insincerely (wait for other) inefficent contracts executed.





- Formation
- Optimization within
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For many agents: voting, general equilibrium, auctions

For fewer agents: auctions, contract nets, bargaining

Possible in all: coalitions



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All self-interested, rational agents

