Parallel Systems Welcome to cs380p

Chris Rossbach + Calvin Lin

CS380p

Outline for Today

- Course Overview
- Course Details and Logistics
- Concurrency & Parallelism Basics
 - Motivation
 - Problem Decomposition

Acknowledgments: some materials in this lecture borrowed from or built on materials from:

- Emmett Witchel, who borrowed them from: Kathryn McKinley, Ron Rockhold, Tom Anderson, John Carter, Mike Dahlin, Jim Kurose, Hank Levy, Harrick Vin, Thomas Narten, and Emery Berger
- Mark Silberstein, who borrowed them from: Blaise Barney, Kunle Olukoton, Gupta

Course Details

Course Name:	CS380P – Parallel Systems	
Lectures:	Online	
Class Web Page: http://www.cs.utexas.edu/users/rossbach/cs380p		
Instructors:	<u>Chris Rossbach + Calvin Lin</u>	
Text:	Principles of Parallel Programming (ISBN-10: 0321487907)	

PRINCIPLES OF PARALLEL PROGRAMMING



Please read the syllabus!



• Parallelism is super-cool and super-important

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- Have *fun* programming cool systems
 - GPUs! Multi-core!
 - Modern infrastructure and programming languages
 - Interesting synchronization primitives (not just about locks!)

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Two perspectives:

- The "just eat your kale and quinoa" argument
- The "it's going to be fun" argument



















(also good for playing heavy metal music)









My current computer



My current computer



Too boring...













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Applications















	ice API Me	ice API MiC	ice API me	ice API								
p Vendor-specific												
	driver		driver									
	vCPU	VGP U	vDIS K	vASI C								
	vNV M	vFP GA	vDS P	VCR PT								
HYPERVISOR												
	CPU	GPU	DISK	ASIC								
	NVM	FPG A	DSP	CRY PT								

	ice API Me	ice API	ice API me	ice API	ice API me	ice API	ice API TTTE		
d v(driver P Vendor-specific								
	vCPU vNV M	Ч (vGP U vFP GA	۲ ۱R	vDIS K vDS P		VASI C VCR PT		
	HYPERVISOR								
N	CPU NVM	F	GPU FPG		DISK DSP		ASIC CRY PT		












Cluster OS



Introduction



Cluster OS



Introduction



Introduction











Introduction



Key concerns:

- Concurrency/parallelism can't be avoided anymore (want a job?)
- A program or two playing with locks and threads isn't enough

Course goal is to expose you to lots of ways of programming systems like these

CPU(s) GPU Image DSP Crypto ...



<u>Goal</u>: Make Parallelism Your Close Friend <u>Method</u>: Use Many Different Approaches

Abstract	Concrete
Locks and Shared Memory Synchronization	Basic Locking
	Prefix sum – pthreads
Language Support	Go lab: condition variables, channels, go routines Rust lab: type-safety, 2PC
Parallel Architectures	GPU Programming Lab
HPC	MPI: Barnes-Hut lab
Modern/Advanced Topics	 Specialized Runtimes / Programming Models Auto-parallelization Race Detection



One instruction at a time (apparently)



One instruction at a time (apparently)



Multiple instructions in parallel







Multiple instructions in parallel



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Technology Trends

35 YEARS OF MICROPROCESSOR TREND DATA



Original data collected and plotted by M. Horowitz, F. Labonte, O. Shacham, K. Olukotun, L. Hammond and C. Batten Dotted line extrapolations by C. Moore

Free lunch – is over Θ



Original data collected and plotted by M. Horowitz, F. Labonte, O. Shacham, K. Olukotun, L. Hammond and C. Batten Dotted line extrapolations by C. Moore

SISD	SIMD		
Single Instruction stream	Single Instruction stream		
Single Data stream	Multiple Data stream		
MISD	MIMD		
Multiple Instruction stream	Multiple Instruction stream		
Single Data stream	Multiple Data stream		





Uncommon architecture: Fault – tolerance Pipeline parallelism



prev instruct	prev instruct	prev instruct	h.
load A(1)	load A(2)	load A(n)	
load B(1)	load B(2)	load B(n)	=
C(1)=A(1)*B(1)	C(2)=A(2)*B(2)	C(n)=A(n)*B(n)	me
store C(1)	store C(2)	store C(n)	
next instruct	next instruct	next instruct	+
P1	P2	Pn	



• Example: vector operations (e.g., Intel SSE/AVX, GPU)



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MIMD

MIMD

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 - SPMD
 - Input domain
 - Output Domain
 - Both
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Game of Life

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- Given a 2D Grid:
- $v_t(i,j) = F(v_{t-1}(of \ all \ its \ neighbors))$



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What decomposition fits "best"?

- Domain (data parallel)
- Functional (task parallel)

Each CPU gets part of the input

Each CPU gets part of the input



Each CPU gets part of the input



For next time:

Each CPU gets part of the input



For next time:

• What issues/challenges might arise with this solution?

Each CPU gets part of the input



For next time:

- What issues/challenges might arise with this solution?
- How could we do a functional decomposition?