Hardware Success Stories

Chris Rossbach and Calvin Lin

Computer Architecture

The Instruction Set Architecture (ISA)

The ISA is an example of a successful parallel abstraction

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The Instruction Set Architecture (ISA) The ISA is an example of a successful parallel abstraction

Today we'll look at microprocessor trends over the years

What can we learn from this success story?

Parallelism in Hardware

Microprocessors are highly parall Consider a block diagram of Pentium processor



Figure 1. Pentium block diagram.

Bit-serial ALUs







Divide an instruction into stages

Stage 1		Stage 2		Stage 3		Stage 4		Stage 5
Instruction Fetch	\Rightarrow	Instruction Decode & Register Fetch	⇒	Execute	\Rightarrow	Memory Access	\Rightarrow	Register Write-back
IF	\Rightarrow	ID/RF	\Rightarrow	EX	\Rightarrow	MEM	\Rightarrow	WB

time —



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IF	\Rightarrow	ID/R	F	=	⇒	EX	\Rightarrow	Ν	1EM		>	WI	3
				time									
A form of tas	A form of task parallelism			IF	ID	EX	MM	WB		_			
					IF	ID	EX	MM	WB				
Increases la	itenc	v of a	instructions			IF	ID	EX	MM	WB		_	
		•	SU				IF	ID	EX	MM	WB		
single inst	ructi	on	1					IF	ID	EX	MM	WB	
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				time										
A form of task parallelism			ins	IF	ID	EX	MM	WB						
			tru		IF	ID	EX	MM	WB					
Increases latency of a			instructions			IF	ID	EX	MM	WB				
		•	SU				IF	ID	EX	MM	WB			
single inst	.ructi	ON	1					IF	ID	EX	MM	WB		
Increases throughput-									IF	ID	EX	ММ	WB	
			•											

ideally completes one instruction per cycle

Program order

Bit-serial ALUs Bit-parallel ALUs Pipelined microprocessors Out-of-order execution

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Pro	Program order										
lss	Issue order										
	time		-		-						
ins	IF	ID	EX	MM	WB						
tru		IF	ID	EX	ММ						
						WB				1	
ctio			IF	ID	EX	VV B		MM	WB		
instructions						EX	MM	MM WB	WB		
ctions				ID	EX		MM EX		WB WB		



How does out-of-order execution help?

Program order

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Issue order

How does out-of-order execution help? Initiate a slow instruction early



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Bit-parallel ALUs

Pipelined microprocessors

Out-of-order execution

Superscalar execution



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Bubbles in the pipeline represent performance loss



Bubbles in the pipeline represent performance loss



Fill bubbles with instructions from other threads

Bubbles in the pipeline represent performance loss



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Bubbles in the pipeline represent performance loss



What is the cost of multithreading? Is it ever not profitable?



roblem – Still have lots of unfilled issue slots

Improving Hardware Utilization

Limitation of Multithreading

At each cycle, issue instructions from any one thread



Improving Hardware Utilization

Limitation of Multithreading At each cycle, issue instructions from any one thread

Simultaneous Multithreading (SMT)

[Tulsen, Eggers, Levy, ISCA95] At each cycle, issue instructions from multiple threads



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What is the cost of SMT?

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Out-of-order execution

Superscalar execution

Multithreading

SMT

VLIW

VLIW

Very Long Instruction Word

Wide instructions for controlling multiple functional units
ISA explicitly encodes parallelism
Statically scheduled
Reduces power by devoting less
hardware to control



Transmeta TM8000 Core

Vector Processors

Goal:

Reduce instruction fetch bandwidth

Apply instructions to all elements of vectors simultaneously

eg. VectorAdd

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Have compilers convert existing programs to vector programs

Vectorizing compilers have been quite successful

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Can we follow a similar path for parallelizing compilers?

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Can we follow a similar path for parallelizing compilers?

Parallelization is more difficult.

The question is not just "is a loop parallelizable?" There are many more tradeoffs to consider.

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Multi-cores

Increasing granularity of parallelism

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Which of these expose parallelism to through the ISA?



For years, the ISA was able to hide hardware parallelism, so programmers could just write sequential code

As the granularity of parallelism has increased, it's become difficult to hide the parallelism

Why is the granularity increasing?



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Why is the granularity increasing?

Is hidden complexity always good?