## Topic 15 Implementing and Using Stacks

"stack n.

The set of things a person has to do in the future. "I haven't done it yet because every time I pop my stack something new gets pushed." If you are interrupted several times in the middle of a conversation, "My stack overflowed" means "I forget what we were talking about."

#### -The Hacker's Dictionary

Friedrich L. Bauer German computer scientist who proposed "stack method of expression evaluation" in 1955.

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#### Stack Overflow



#### **Sharper Tools**







#### **Stacks**

- Access is allowed only at one point of the structure, normally termed the *top* of the stack
  - access to the most recently added item only
- Operations are limited:
  - push (add item to stack)
  - pop (remove top item from stack)
  - top (get top item without removing it)
  - clear
  - isEmpty
  - size?
- Described as a "Last In First Out" (LIFO) data structure

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## **Stack Operations**

Assume a simple stack for integers. Stack s = new Stack(); s.push(12); s.push(4);s.push(s.top() + 2);s.pop() s.push( s.top() ); //what are contents of stack?

### **Stack Operations**

## Write a method to print out contents of stack in reverse order.

#### **Common Stack Error**

Stack s = new Stack();
// put stuff in stack
for(int i = 0; i < 5; i++)
 s.push( i );
// print out contents of stack
// while emptying it. (??)
for(int i = 0; i < s.size(); i++)
 System.out.print( s.pop() + " ");</pre>

#### // What is output?

#### **Attendance Question 1**

- What is output of code on previous slide?
- A 0 1 2 3 4
- **B** 4 3 2 1 0
- **C** 4 3 2
- **D** 2 3 4
- E No output due to runtime error.

#### **Corrected Version**

Stack s = new Stack(); // put stuff in stack for(int i = 0; i < 5; i++)s.push(i); // print out contents of stack // while emptying it int limit = s.size(); for (int i = 0; i < limit; i++) System.out.print( s.pop() + " "); //or // while( !s.isEmpty() ) System.out.println( s.pop() ); CS 307 Fundamentals of 9 Stacks **Computer Science** 

## Implementing a stack

- need an underlying collection to hold the elements of the stack
- 2 basic choices
  - array (native or ArrayList)
  - linked list
- array implementation
- Inked list implementation
- Some of the uses for a stack are much more interesting than the implementation of a stack

#### **Applications of Stacks**

#### **Problems that Use Stacks**

- The runtime stack used by a factorial process (running program) to keep track of methods in factorial progress
- Search problems
- Undo, redo, back, forward







factorial

main



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#### **Mathematical Calculations**

What is 3 + 2 \* 4? 2 \* 4 + 3? 3 \* 2 + 4?

The precedence of operators affects the order of operations. A mathematical expression cannot simply be evaluated left to right.

A challenge when evaluating a program.

*Lexical analysis* is the process of interpreting a program. Involves Tokenization

#### What about 1 - 2 - 4 ^ 5 \* 3 \* 6 / 7 ^ 2 ^ 3

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## Infix and Postfix Expressions

- The way we are use to writing expressions is known as infix notation
- Postfix expression does not
- require any precedence rules
- 32\*1 + is postfix of 3 \* 2 + 1
- evaluate the following postfix expressions and write out a corresponding infix expression:



1234 \* +

25^1-

#### **Attendance Question 2**

What does the following postfix expression evaluate to?

632+\*

- A. 18
- B. 36
- C. 24

D. 11

#### E. 30

#### **Evaluation of Postfix Expressions**

- Easy to do with a stack
- given a proper postfix expression:
  - get the next token
  - if it is an operand push it onto the stack
  - else if it is an operator
    - pop the stack for the right hand operand
    - pop the stack for the left hand operand
    - apply the operator to the two operands
    - push the result onto the stack
  - when the expression has been exhausted the result is the top (and only element) of the stack

#### Infix to Postfix

- Convert the following equations from infix to postfix:
  - 2 ^ 3 ^ 3 + 5 \* 1
  - 11 + 2 1 \* 3 / 3 + 2 ^ 2 / 3

Problems:

Negative numbers?

parentheses in expression

#### Infix to Postfix Conversion

- Requires operator precedence parsing algorithm
  - parse v. To determine the syntactic structure of a sentence or other utterance
- Operands: add to expression
- Close parenthesis: pop stack symbols until an open parenthesis appears
- Operators:
  - Have an on stack and off stack precedence
  - Pop all stack symbols until a symbol of lower precedence appears. Then push the operator
- End of input: Pop all remaining stack symbols and add to the expression

Infix Expression:

3 + 2 \* 4

PostFix Expression:

Operator Stack:

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| /      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

#### Infix Expression: + 2 \* 4

PostFix Expression: 3

Operator Stack:

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| 1      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

# Infix Expression:2 \* 4PostFix Expression:3

Operator Stack: +

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| 1      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

# Infix Expression:\* 4PostFix Expression:3 2

Operator Stack: +

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| /      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

# Infix Expression:4PostFix Expression:3 2Operator Stack:+ \*

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| /      | 2          | 2          |
| Λ      | 10         | 9          |
| (      | 20         | 0          |

Infix Expression:

PostFix Expression: 324

Operator Stack: + \*

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| /      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

Infix Expression:

PostFix Expression: 324\*

Operator Stack: +

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| 1      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

Infix Expression:

PostFix Expression: 324\*+

Operator Stack:

| Symbol | Off Stack  | On Stack   |
|--------|------------|------------|
|        | Precedence | Precedence |
| +      | 1          | 1          |
| -      | 1          | 1          |
| *      | 2          | 2          |
| 1      | 2          | 2          |
| ٨      | 10         | 9          |
| (      | 20         | 0          |

#### Example

1 - 2 ^ 3 ^ 3 - ( 4 + 5 \* 6 ) \* 7

Show algorithm in action on above equation

## **Balanced Symbol Checking**

In processing programs and working with computer languages there are many instances when symbols must be balanced {},[],()

A stack is useful for checking symbol balance. When a closing symbol is found it must match the most recent opening symbol of the same type.

Algorithm?

## Algorithm for Balanced Symbol Checking

- Make an empty stack
- read symbols until end of file
  - if the symbol is an opening symbol push it onto the stack
  - if it is a closing symbol do the following
    - if the stack is empty report an error
    - otherwise pop the stack. If the symbol popped does not match the closing symbol report an error
- At the end of the file if the stack is not empty report an error

## Algorithm in practice

- Iist[i] = 3 \* (44 method( foo( list[ 2 \* (i + 1) + foo( list[i 1])) / 2 \* ) list[ method(list[0])];
- Complications
  - when is it not an error to have non matching symbols?
- Processing a file
  - *Tokenization*: the process of scanning an input stream.
     Each independent chunk is a token.
- Tokens may be made up of 1 or more characters