CS 345

Functions

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Reading Assignment

Mitchell, Chapter 7

C Reference Manual, Chapters 4 and 9

Procedural Abstraction

Contains local variable declarations Can be overloaded and statements and statements

(e.g., binary +) 🔪 into the scope Procedure is a named parameterized scope

 Allows programmer to focus on a function interface, ignoring the details of how it is computed

Value-returning functions

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• Example: x = (b*b - sqrt(4*a*c))/2*a

Non-value returning functions

- Called "procedures" (Ada), "subroutines" (Fortran), "void functions/methods" (C, C++, Java)
- Have a <u>visible side effect</u>: change the state of some data value not defined in the function definition
- Example: strcpy(s1,s2)

System Calls

OS procedures often return status codes

 Not the result of computing some function, but an indicator of whether the procedure succeeded or failed to cause a certain side effect

```
int open(const char* file, int mode)
{
    if (file == NULL) {
        return -1; // invalid file name
    if (open(file, mode) < 0)
        return -2; // system open failed
    ...</pre>
```

Arguments and Parameters

int h, i; Argument: expression that void B(int w) { appears in a function call Parameter: identifier that $i = 2^*w;$ W = W + 1;appears in function declaration Parameter-argument matching d A(int x, int y) { by number and position ooli,j; B(h) Exception: Perl. Instead of being declared in a function header, main() { int parameters are available as nt a, b; elements of special array @_ = 5; a = 3; b = 2; A(a, b);

Parameter Passing Mechanisms

By value
By reference
By value-result
By result
By name

Pass by Value

Caller passes r-value of the argument to function

- Compute the <u>value</u> of the argument at the time of the call and assign that value to the parameter
- Reduces "aliasing"
 - Aliasing: two names refer to the same memory location
- Function cannot change value of caller's variable

All arguments in C and Java are passed by value

- To allow caller's variables to be modified, pointers can be passed as arguments
 - Example: void swap(int *a, int *b) { ... }
 - Is there a contradiction here?

Pass by Reference

- Caller passes I-value of the argument to function
 - Compute the <u>address</u> of the argument and assign that address to the parameter
 - Increases aliasing (why?)
- Function can modify caller's variable via the address it received as argument

int h, i; void B(int* w) { int j, k; i = 2*(*w); *W = *W+1; }

ML Example

pseudo-code



function f(x) =

{ x = x+1; return x; } var y = 0;

print (f(y)+y);



Standard ML

fun f (x : int ref) =
 (x := !x+1; !x);
y = ref 0 : int ref;
f(y) + !y;

fun f (z : int) =
 let x = ref z in
 x := !x+1; !x
 end;
y = ref 0 : int ref;
f(!y) + !y;

Pass by Reference in C++

- Special "reference type" indicates that I-value is passed as argument
 - Recall that in C, only r-values can be arguments
 void swap ((int& a, (int& b))
 int temp = a;
 a = b;
 b = temp;
 }
- operator is overloaded in C++
 - When applied to a variable, gives its I-value
 - When applied to type name in parameter list, means pass the argument by reference

Two Ways To Pass By Reference

C or C++

```
void swap (int *a, int *b) {
    int temp = *a;
    *a = *b;
    *b = temp;
}
```

int x=3, y=4; swap(&x, &y); C++ only

int x=3, y=4;

swap(x, y);

void swap (int& a, int& b) {
 int temp = a;
 a = b;
 b = temp;
}

Which one is better? Why?

Pass by Value-Result

- Pass by value at the time of the call and/or copy the result back to the argument at the end of the call (copy-in-copy-out)
 - Example: "in out" parameters in Ada
- Reference and value-result are the same, except when aliasing occurs
 - Same variable is passed for two different parameters
 - Same variable is both passed and globally referenced from the called function

Pass by Name

- Textually substitute the argument for every instance of its corresponding parameter in the function body
 - Originated with Algol 60 but dropped by Algol's successors -- Pascal, Ada, Modula

Example of late binding

- Evaluation of the argument is delayed until its occurrence in the function body is actually executed
- Associated with lazy evaluation in functional languages (e.g., Haskell)

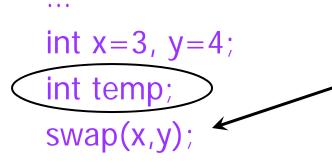
Jensen's Device

```
• Computes \sum_{i=1}^{100} \frac{1}{i} in Algol 60
                                          passed by name
           begin
             integer i;
             real procedure sum (i, lo, hi, term);
                  value lo, hi;
                  integer i, lo, hi;
                  real term;
             begin
                  real temp;
                  temp := 0;
                  for i := lo step 1 until bi do
                     temp := temp +(term;)
                                             becomes 1/i when
                  sum := temp
                                             sum is executed
             end:
             print (sum (i, 1, 100, 1/i))
           end
```

Macro

Textual substitution

#define swap(a,b) temp=a; a=b; b=temp;



Textually expands to temp=x; x=y; y=temp;



- Does not obey the lexical scope rules (i.e., visibility of variable declarations)
- No type information for arguments or result

Problems with Macro Expansion

#define swap(a,b) temp=a; a=b; b=temp;

 $\begin{array}{ccc} & & & & \\ \text{if } (x < y) & & \\ & & &$

Why not #define swap(a,b) { int temp=a; a=b; b=temp; }?

Instead #define swap(a,b) do { Fixes type of swapped variables
 inf temp=a; a=b; b=temp;
 } while(false);

Variable Arguments

In C, can define a function with a variable number of arguments

- Example: void printf(const char* format,(...)
- Examples of usage:

printf("hello, world"); printf("length of (%s) = %d)n", str, str.length()); printf("unable to open file descriptor %d)n", fd);

Format specification encoded by special %-encoded characters

- %d,%i,%o,%u,%x,%X integer argument
- %s string argument
- %p pointer argument (void *)
- Several others (see C Reference Manual!)

Part of

syntax!

Implementation of Variable Args

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Special functions va_start, va_arg, va_end compute arguments at run-time (how?)

```
void printf(const char* format, ...)
ł
     int i; char c; char* s; double d;
     va list ap; /* declare an "argument pointer" to a variable arg list */
    va start(ap, format); /* initialize arg pointer using last known arg */
     for (char* p = format; *p != \ 0'; p++) {
       if (*p == `%') {
          switch (*++p) {
            case 'd':
               i = va arg(ap, int); break;
            case 's':
               s = va arg(ap, char*); break;
            case 'c':
               c = va arg(ap, char); break;
            ... /* etc. for each % specification */
          }
    va end(ap); /* restore any special stack manipulations */
```