

# Aggregating Inter-App Traffic to Optimize Cellular Radio Energy Consumption on Smartphones

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# Outline

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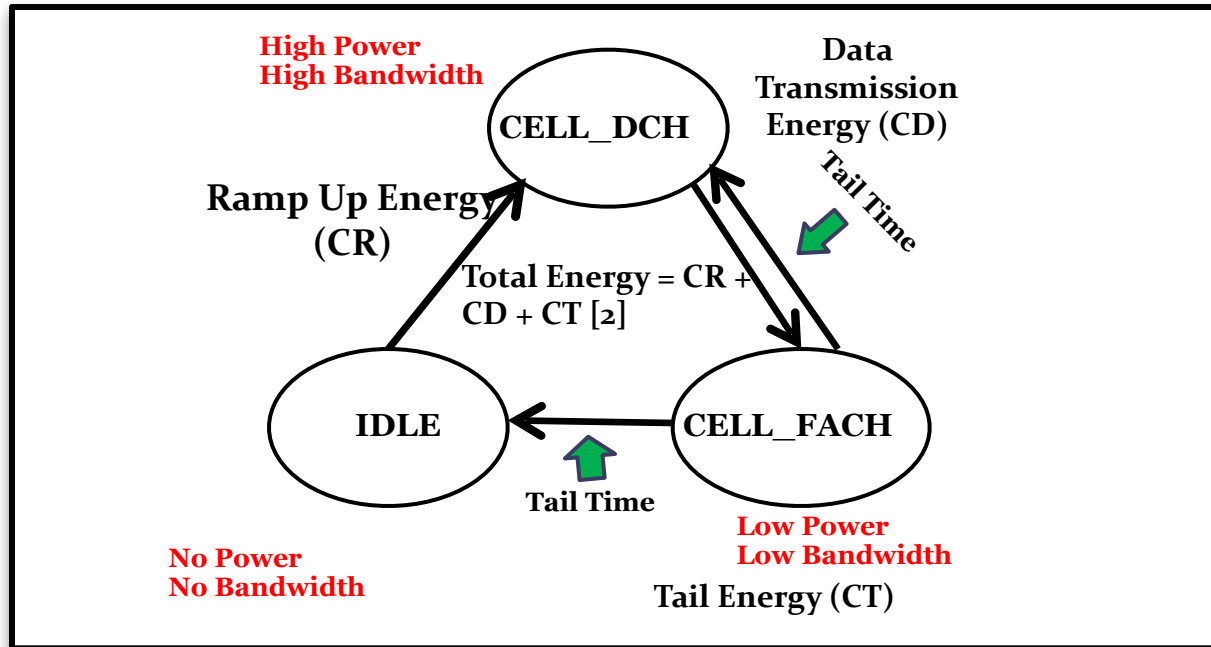
- Introduction
- Balanced Scheduling Protocol
- Evaluation Setup and Traces
- Results and Takeaways
- Conclusion and Future Works

# Introduction

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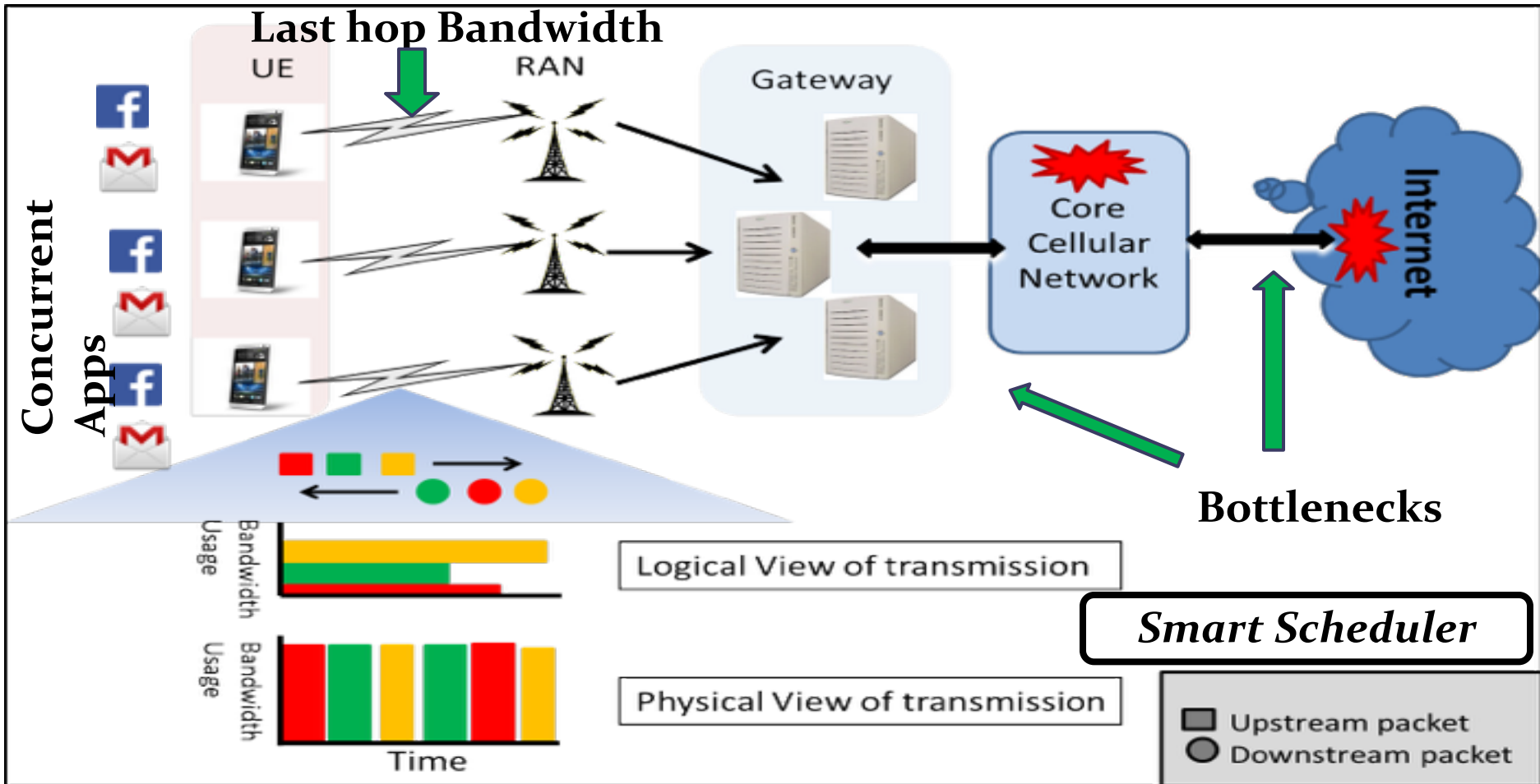
- ❑ Display, Network, and CPU are main components of Energy Drain. [ **Mittal et. al., MobiCom '13** ]
- ❑ Poorly written apps can sap 30% to 40% of a phone's battery. [ **Mahajan et. al., IMC '09** ]
- ❑ Network intensive applications are increasing ( ~69% of the apps are cloud based ).
- ❑ Different background services running intermittently and waking up the network card for a small duration. [ **Qian et. al., WWW '12** ]

# Cellular Radio Energy Model

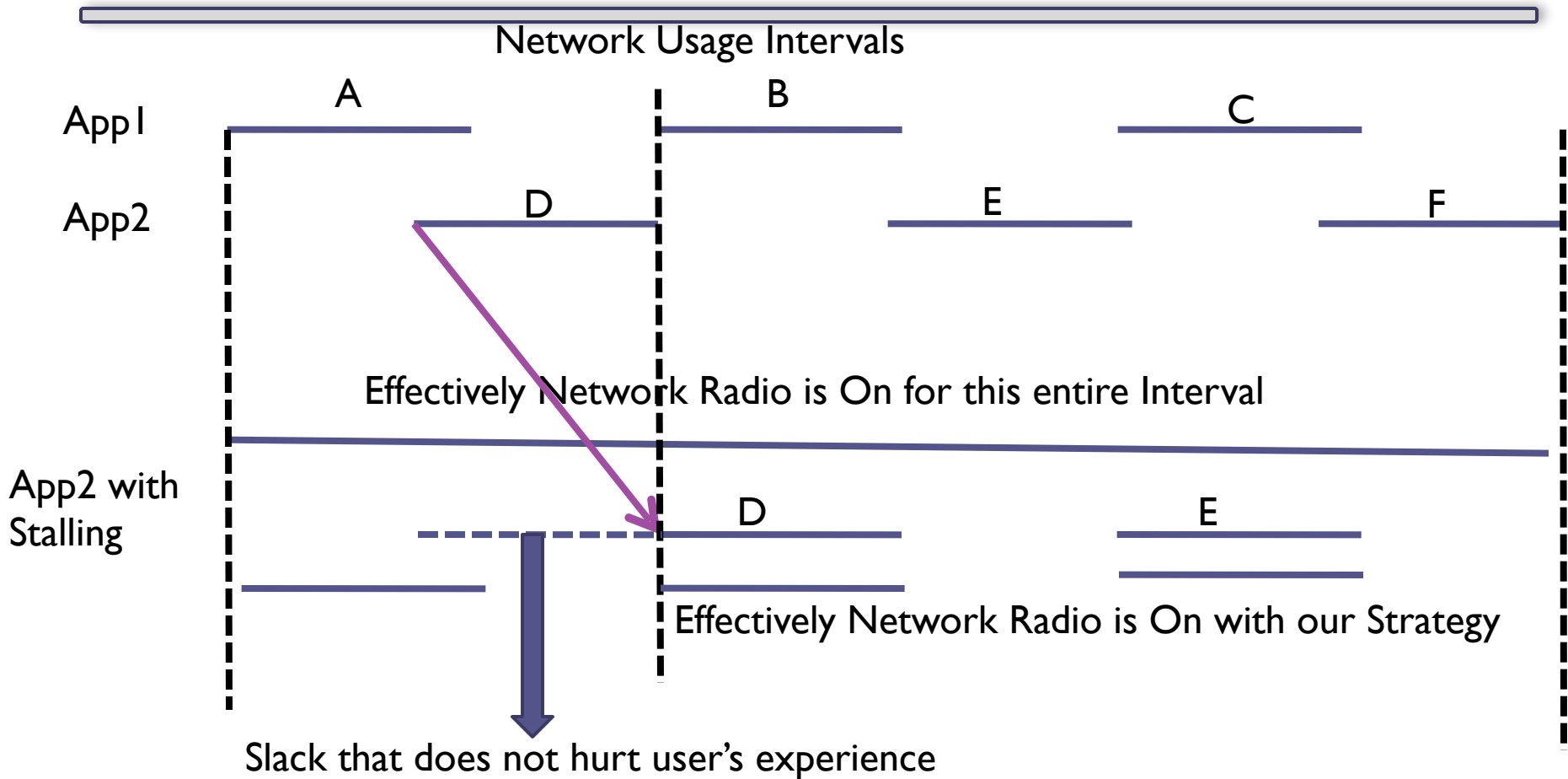


**Total Energy Consumption = CR + CD + CT** ,  
where CR is the ramp up energy (**IDLE to CELL DCH**),  
CD is the data transmission energy,  
and CT is the tail energy (in **CELL FACH**).

# Total View of App Connectivity



# Idea : Cross Application Traffic Aggregation



# Problem Objective

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- ❑ Scheduling all network requests using minimum energy without hurting user's experience.
  
- ❑ Multi-objective optimization problem
  - Minimum Energy => **Best utilization of bandwidth** (Side-effect : Lower Switching Frequency).
  - User's Experience => Request should be served within deadline.

# Problem Constraints

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- Depending on expected response time of application, a **flexibility or slack time** is allowed to schedule each packet.
- Requests from the **same application cannot be triggered simultaneously**.
- **Total bandwidth consumption** by all the scheduled requests should be **less than the available channel bandwidth** (We consider *constant last hop bandwidth*).



# Approach Intuition : Deciding Function

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- ❑ If a request is delayed then there is potentially more opportunity of batching.
- ❑ If a request is delayed much, it may miss deadline.
- ❑ So, we need to *develop a function* to decide at certain time ***if a request should be scheduled or should be delayed further.***

# Terminology

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- $A_i$  =  $i^{\text{th}}$  application
- $A_{ij}$  =  $j^{\text{th}}$  request of  $i^{\text{th}}$  application
- $r_{ij}$  = Arrival time of  $A_{ij}$
- $x_{ij}$  = Scheduling time of  $A_{ij}$
- $f_{ij}$  = slack time time of  $A_{ij}$
- $d_{ij}$  = service duration of  $A_{ij}$
- $ft$  = finish time of all requests in run queue

# Deciding Function (F)

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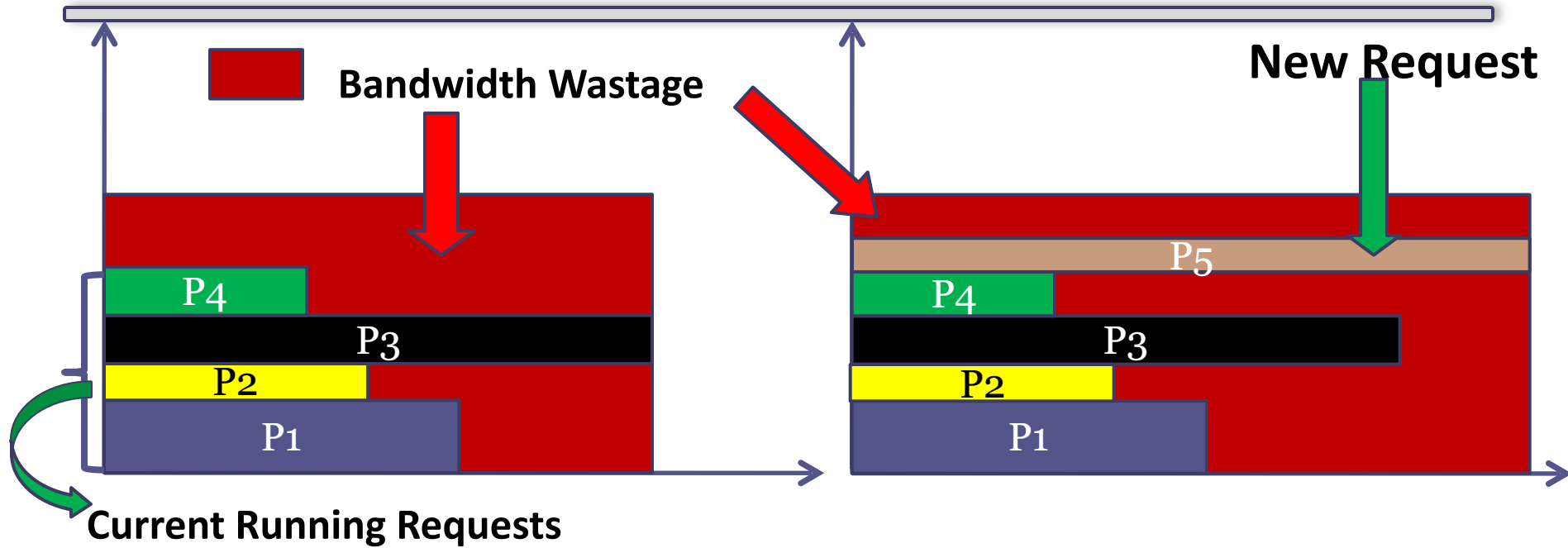
$$F = \beta \cdot \text{Bandwidth\_wastage} + (1 - \beta) \cdot \text{Experience\_user}$$

Where  $\beta$  is normalizing constant.

$$F = \alpha \cdot \beta \cdot \text{Bandwidth\_wastage} + (1 - \alpha)(1 - \beta) \cdot \text{Experience\_user}$$

Where  $\alpha$  is factor to give priority over other

# F : Bandwidth Wastage Component

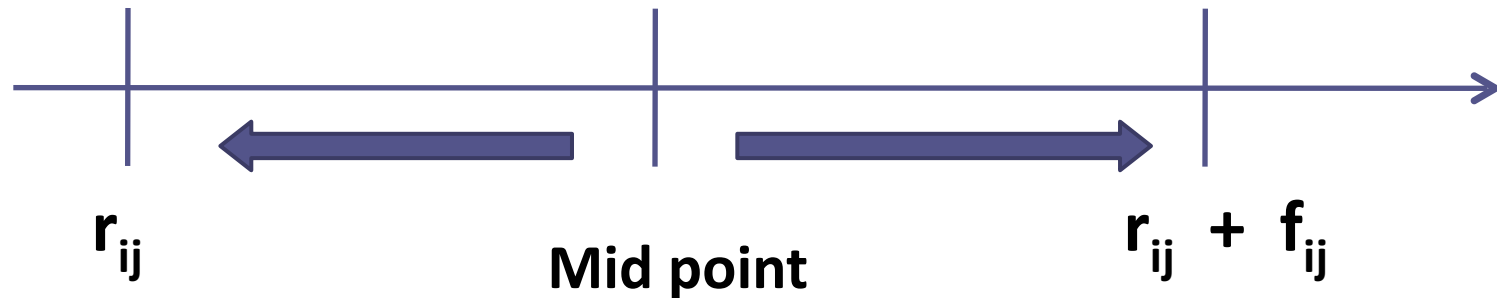


$$\text{Bandwidth\_wastage} = \frac{BW1 - BW2}{\text{Max}(BW1, BW2)}$$

**BW1 : Bandwidth Wastage Before Scheduling**

**BW2 : Bandwidth Wastage After Scheduling**

# F : User Experience Component



Do not schedule if **ft** lies here Wait for more requests

Schedule if **ft** lies here

$$\text{Experience\_user} = \frac{ft - \frac{r_{ij} + (r_{ij} + f_{ij})}{2}}{\text{Max}(ft, \frac{r_{ij} + (r_{ij} + f_{ij})}{2})}$$

# Balanced Scheduling Protocol

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- ❑ There are two queues:
  - *Running Queue* has all the running requests served by Cellular Radio.
  - *Waiting Queue* has all the pending requests.
  
- ❑ Requests are put into wait queue as soon as they arrive.
  
- ❑ Pushed to run queue when Deciding Function (F) is positive.

# Experimental Setup

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## □ Application Types (MobiSys '12)

- Gaming (Short Bursts)
- Browsing (Medium Bursts)
- Streaming (Large Bursts)

# Experimental Setup

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## □ Synthetic Trace Generation tuning parameters

- User Interaction Timing (Power Law )
- Data Transmission Size ( Power Law with set of sizes)
- Bandwidth Demand ( Fixed Demand per App)
- Slack Duration (Fixed per application type)

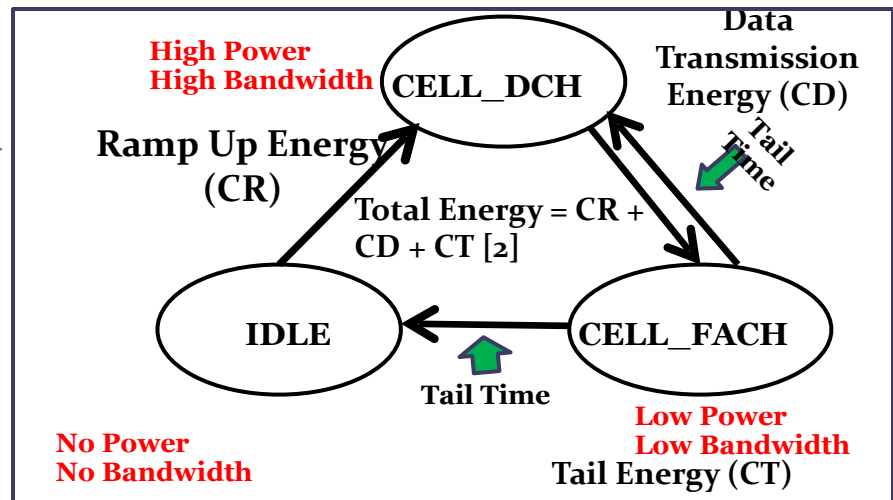
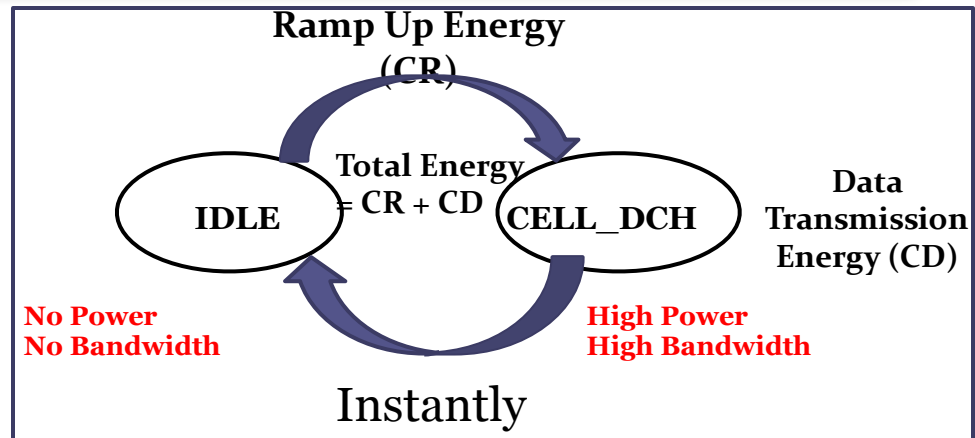


# Experimental Setup

## Switching Strategies

➤ Fast Dormancy.

➤ Fast Dormancy with Fixed Tail Timer.



# Real Trace Collection

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- Collected using ARO (AT&T) tool, tcpdump, and ps.
- Samsung Galaxy S3 GTI9300 (Rooted).
- One hour Browsing Trace from a user.
- Applications are differentiated through port mapping.

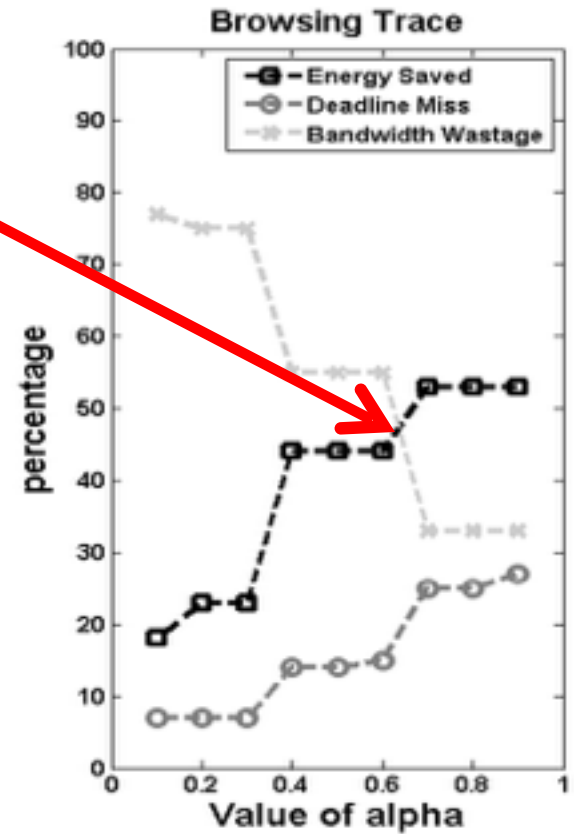
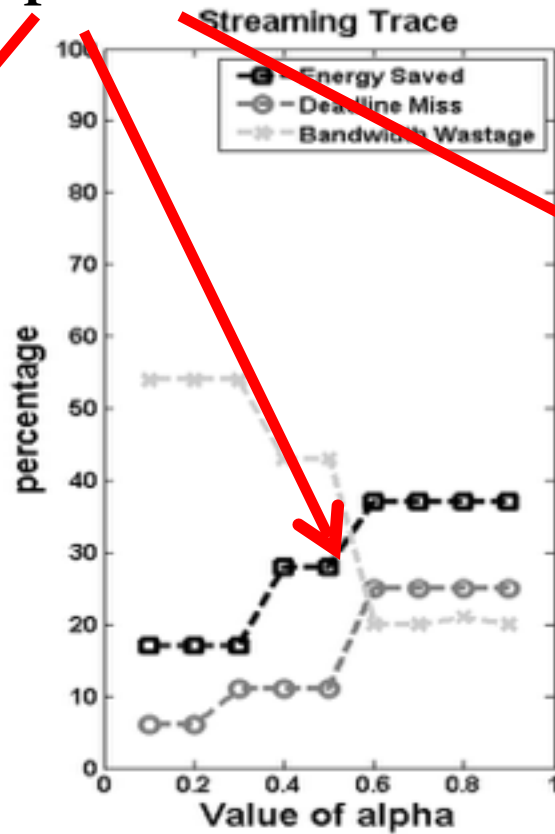
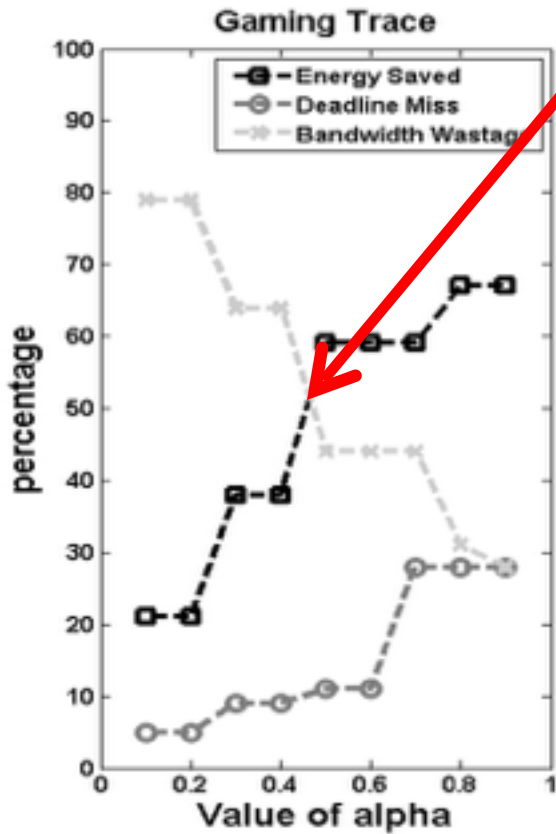
# Evaluation Metrics

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- ❑ **Energy Consumption per KB:** Total energy spent to transmit one KiloByte of data.
- ❑ **Deadline Miss:** Proportion of requests which have missed their transmission deadline.
- ❑ **State Switch Rate:** Number of times per unit time the radio changes state - from IDLE to DCH, and DCH/FACH to IDLE.
- ❑ **Radio On Time:** Radio on time as a fraction of total data transmission duration.

# Alpha Value Tuning

**Good Trade-off between  
Energy and User  
Experience**

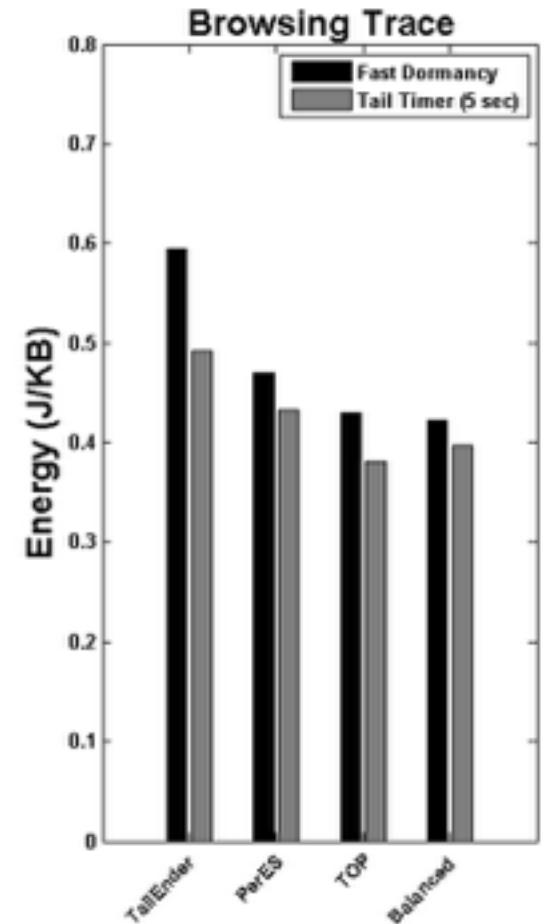
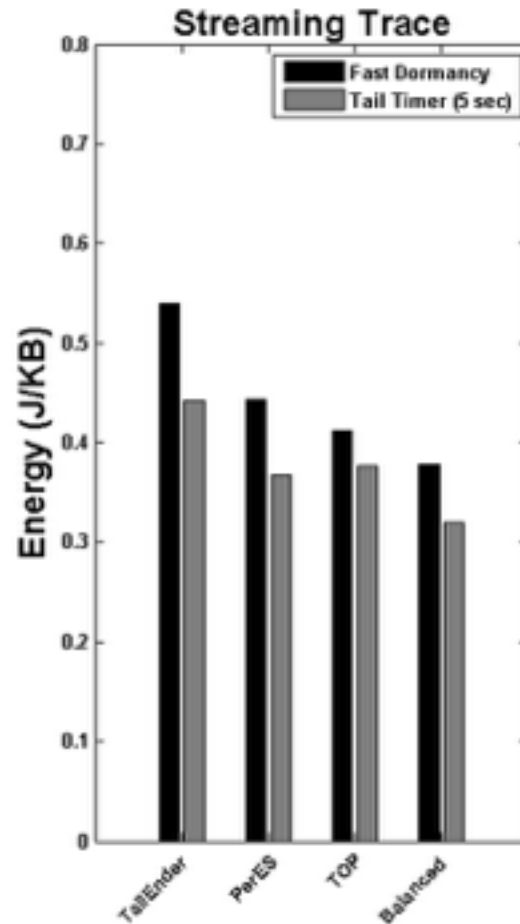
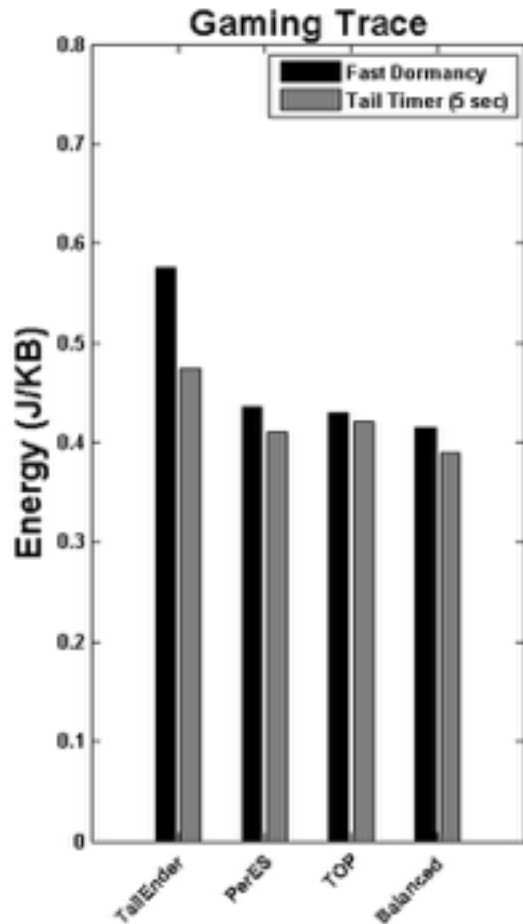


# Competing Scheduling Techniques

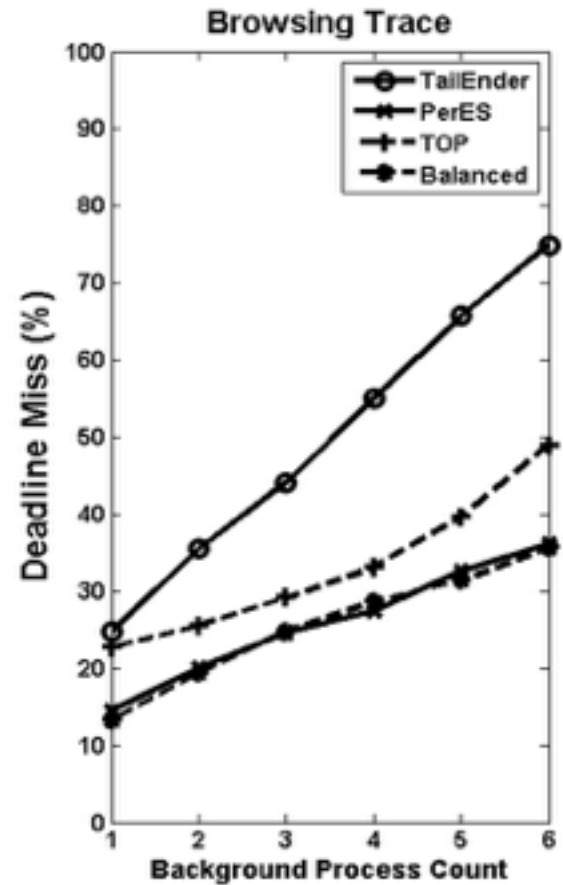
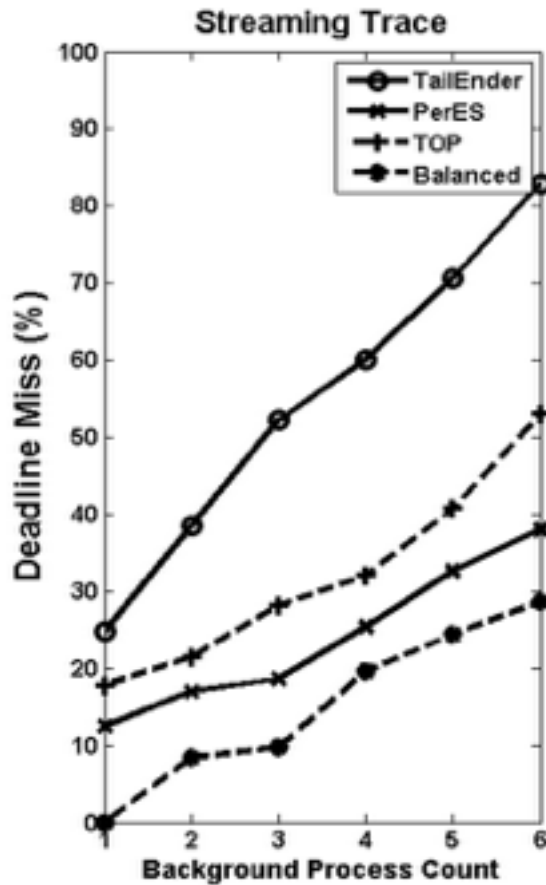
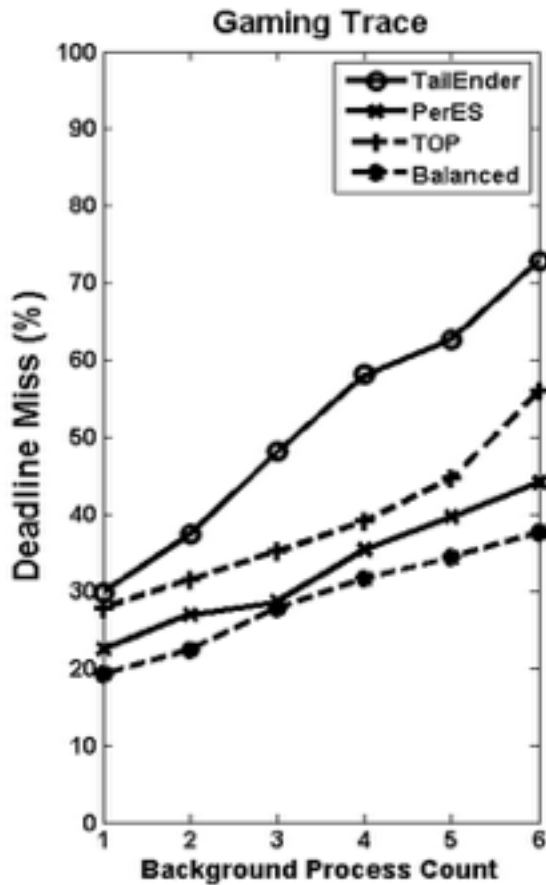
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- ❑ **TailEndeR**: Uses threshold based tail time prediction by considering deadlines of packets of an application.
- ❑ **PerES**: Performance-aware Energy Scheduler or PerES models cross application energy-delay tradeoff as an optimization problem and applies Lyapunov optimization framework.
- ❑ **TOP**: Tail Optimization Protocol reduces tail energy wastage by predicting the application behavior.

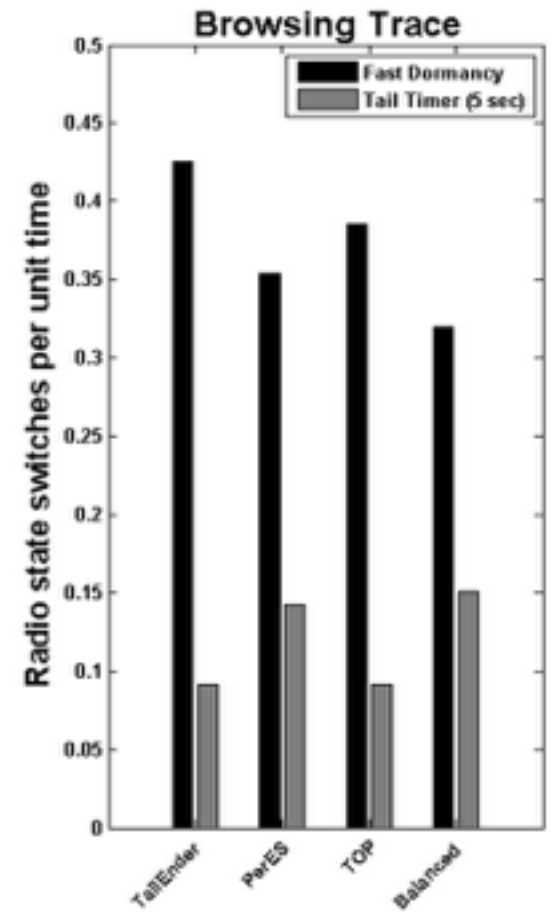
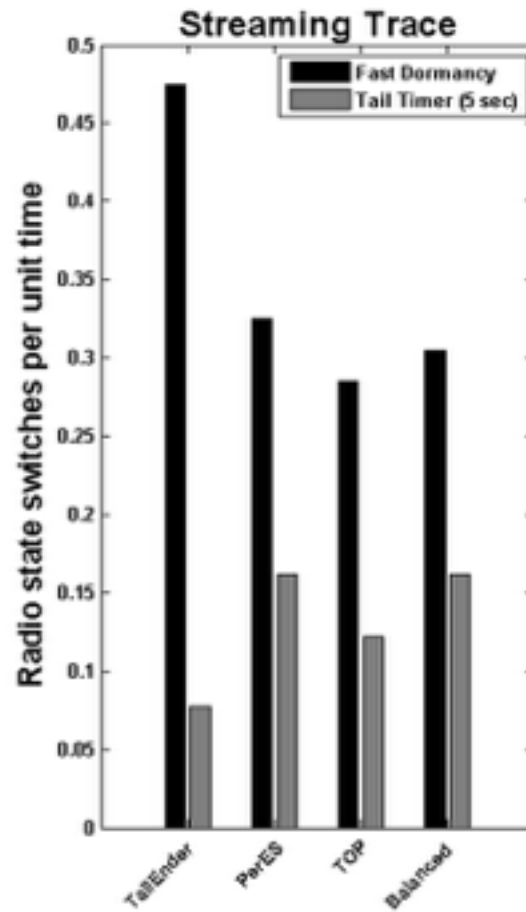
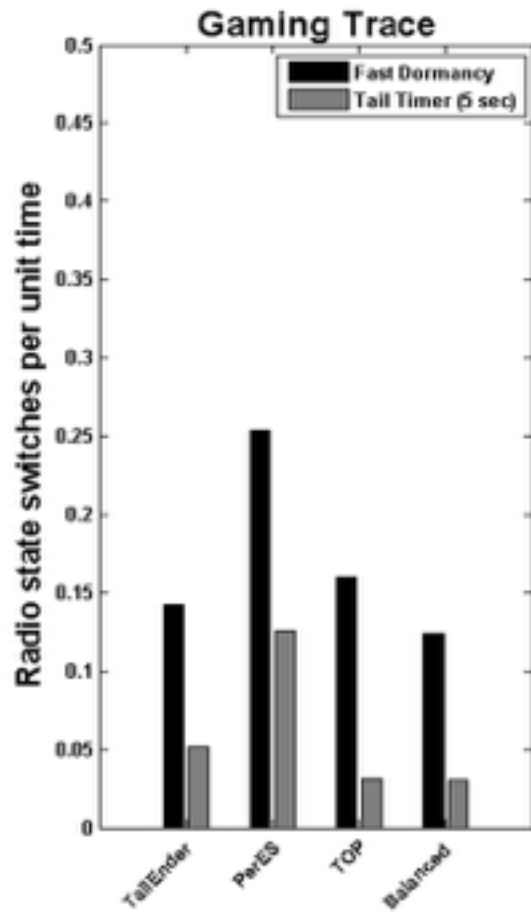
# Energy Consumption per KB (~10%)



# Deadline Miss

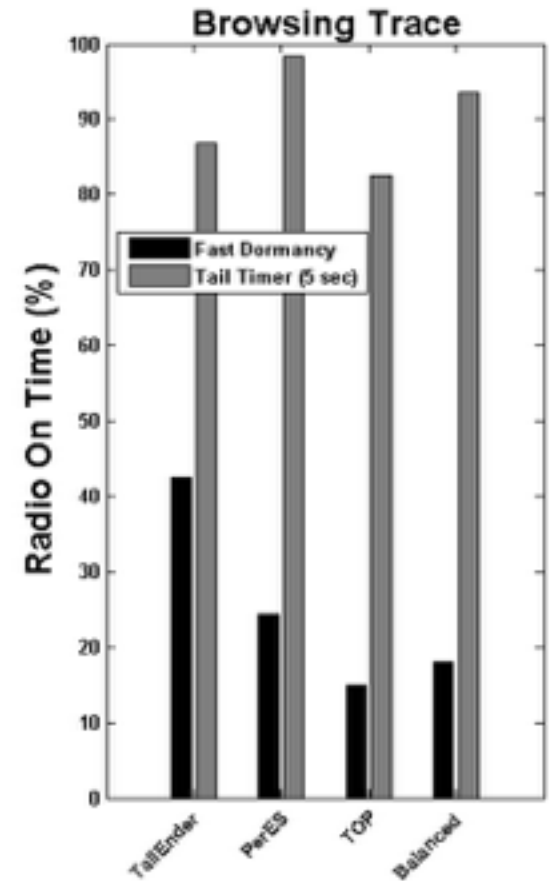
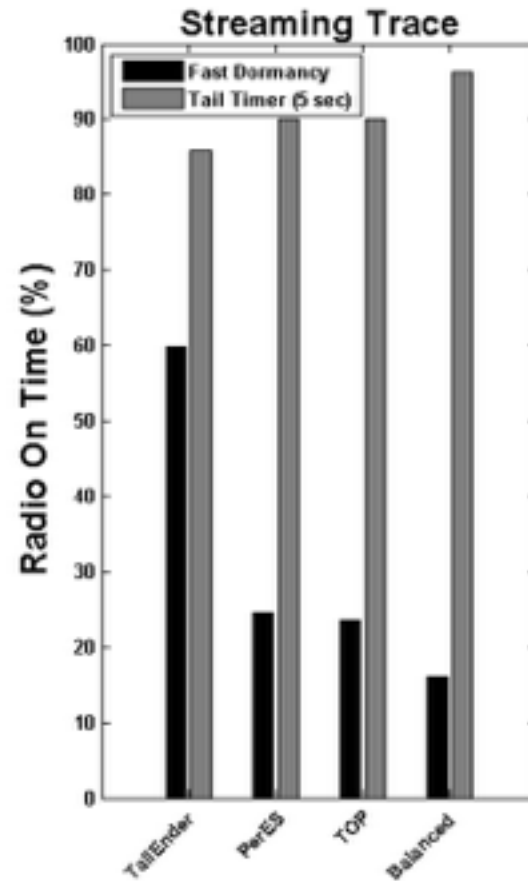
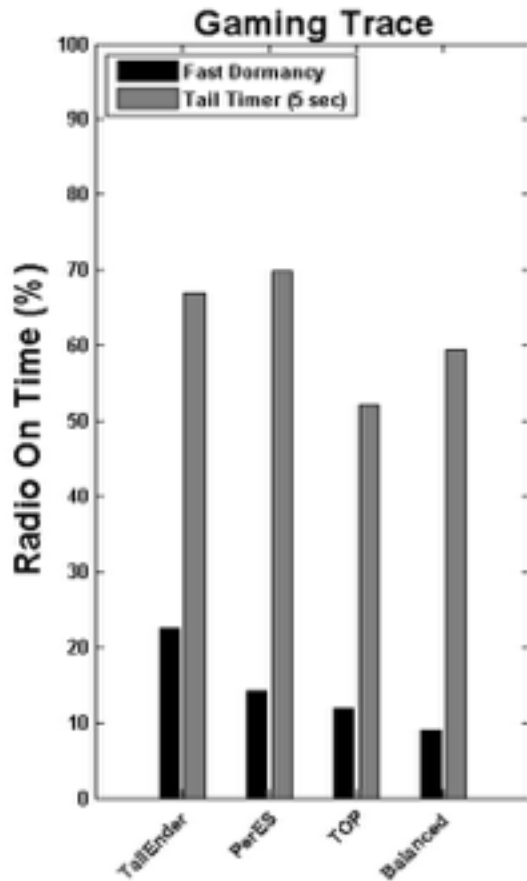


# Switching

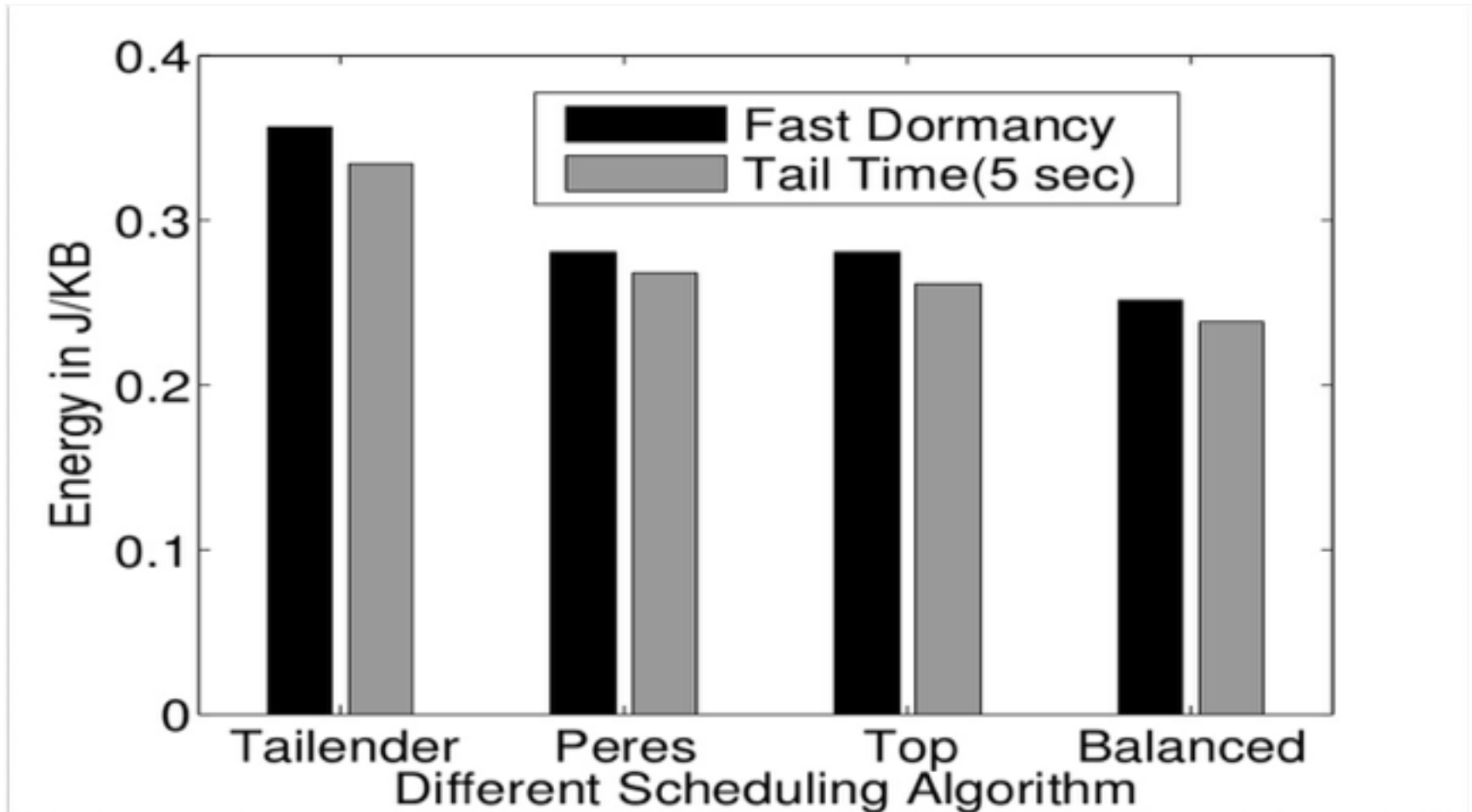




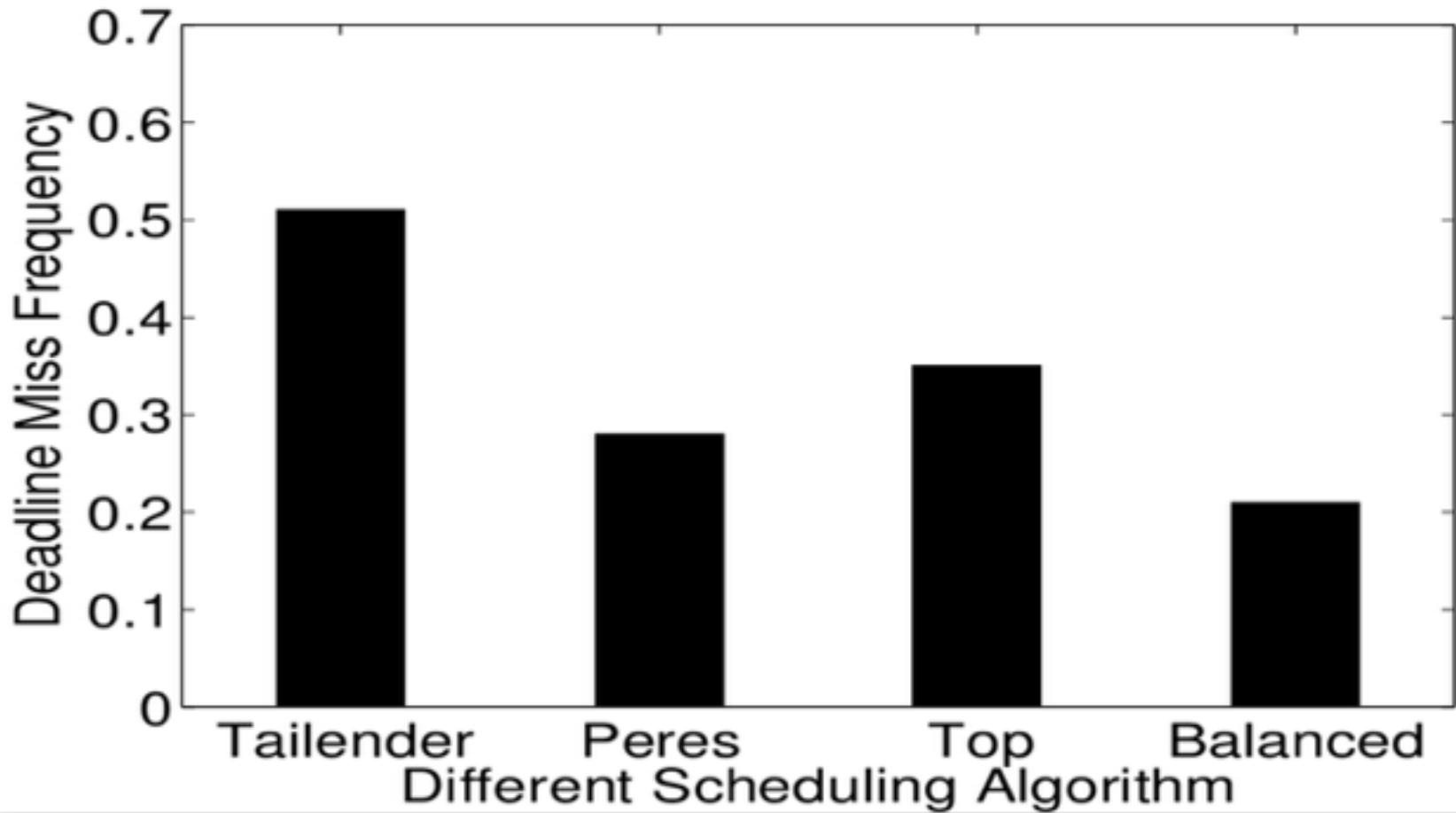
# Radio On Time



# Energy Consumption in Real Trace



# Deadline Miss in Real Trace



# Takeaways

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- ❑ Around 10% better than PerEs and TOP in Energy Gain wise, but far better from TailEnder.
- ❑ Percentage of Deadline Misses for Foreground App remains satisfactory.
- ❑ Reducing number of state transitions of the network interface can save more energy than optimizing utilization of the tail period of the card.

# Future Works

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- ❑ Extensive and large scale real trace based evaluation to validate the simulation based results.
- ❑ Building middleware which will run our aggregation strategy across applications.
- ❑ Extension and implementation of in other elements like sensors, GPS etc.
- ❑ Building a Application network activity recorder tool which can be installed without rooting.

# Thank you

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