

# CS345H: Programming Languages

## Lecture 17: Introduction to Object-Oriented Languages

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- ▶ Not in the sense of how to use it, but to understand the fundamental aspects of this paradigm
- ▶ We will also explore how to formalize some aspects of it

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- ▶ We have seen this before: **Polymorphism**

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- ▶ Here, anyone who expects something of type  $X$  can work just as well with something of type  $Y$
- ▶ **Why?** Because the fields of  $Y$  are a superset of the fields of  $X$

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- ▶ But the exact definition of subtype depends on the language!
- ▶ Observe that subtyping is really another kind of polymorphism as it allows us to write code that works with more than one type

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- ▶ This is also known as “duck typing”

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- ▶ We create an **instance** of a class to use it
- ▶ Every instance has its own data

# Classes Example

- Consider the following program:

```
struct point {  
    int x;  
    int y;  
};  
void inc_x(point* p) {  
    p->x++;  
}  
  
point *p = new point;  
inc_x(p);
```

## Classes Example Cont.

- Here is the same program rewritten with classes:

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class point {  
    int x;  
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- ▶ Here, p is an **object**
- ▶ Object = instantiated class

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- ▶ **Important:** Every instance of a class has its own set of instance variables!
- ▶ **Constructor:** A special method that is run on instance creation

## Classes Example Extended

- Here is a slightly extended version of the program

```
class point {  
    int x;  
    int y;  
    point() { this->x = 0; this->y = 0; }  
    void inc_x() {  
        this->x++;  
    }  
};
```

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point *p = new point();  
point *q = new point();  
p->inc_x();
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- What is the value of y->x? 0

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- ▶ But subtyping now also needs to take methods into account!
- ▶ Specifically, any subtype must also have (at least) the same methods as the original type
- ▶ This way, we can use the subtype in any context that expects the original type



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- ▶ Allowing redefinition of methods is known as **virtual methods**
- ▶ This is another key OO feature

# Virtual Method Example

- Consider the following code:

```
class point {
    int x;
    int y;
    point() { this->x = 0; this->y = 0; }
    virtual void inc_x() {
        this->x++;
    }
};

class bigpoint:public point {
    int x;
    int y;
    virtual void inc_x() {
        this->x+=2;
    }
};

point *p = new bigpoint();
p->inc_x();
```

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- ▶ The dynamic type decides which virtual method is called!
- ▶ This is sometimes called the essence of OO

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- ▶ This is a new use of typing: So far, we only used types to prevent run-time errors
- ▶ In OOP, we also use types **at run-time** to decide which method to invoke!
- ▶ You can think of this operationally as an implicit run-time check on a type tag that decides which version of a method is called

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# Encapsulation

- ▶ The last major OOP feature we have not yet discussed in **encapsulation**
- ▶ This means that class data can be made non-accessible to clients of the class
- ▶ **Example:** Declaring an instance variable `private`
- ▶ Fortunately, this only rejects some programs at compile time but does not change semantics

# Why Encapsulation?

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- ▶ We are therefore free to change the implementation of classes without affecting its clients
- ▶ And we are free to pass any subtype with the same public interface
- ▶ **Actually, old idea:** This is also known as **abstract data types (ADT)** and predates OOP

# Essential OOP Features

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- ▶ This often allows for much cleaner and more extensible code
- ▶ If the problem you are solving fits into the OO model!

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- ▶ In fact, OOP fits so many problems reasonably well that it has become the **default** paradigm used in most software
- ▶ In fact, many modern languages, such as Java, **force** an OOP style
- ▶ However, not every problem maps well into objects!

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- ▶ The idea of objects and subtyping originated at MIT in the 1950s and 1960s
- ▶ This was in the context of AI research in LISP
- ▶ Over the years, various features that we would call object-oriented today made their way into various LISP dialects

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# Object Orientation: History Cont.

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  - ▶ Developed by Ole-Johan Dahl and Kristen Nygaard at the Norwegian Computing Center in Oslo
  - ▶ Simula was designed as an **special-purpose** language for **discrete event simulations**
  - ▶ But it as certainly not designed as a general-purpose programming language

# Features Combined in Simula

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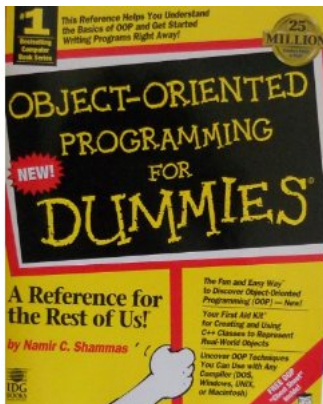
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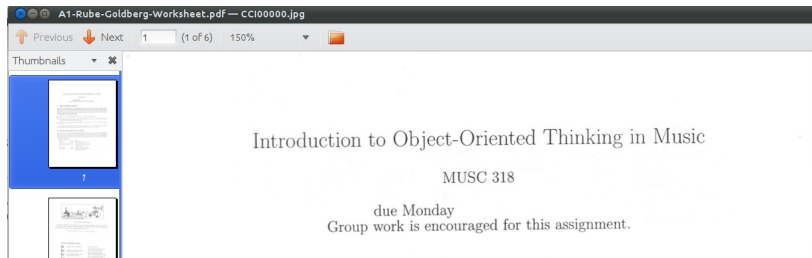
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- ▶ Smalltalk coined the term “object oriented”
- ▶ And lead to a huge wave of OO languages
- ▶ This was an huge fad in the late 90's

# OO in the Real World



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# Other OO languages

- ▶ C++

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- ▶ C++
- ▶ Java
- ▶ C#
- ▶ Pretty much any newer imperative language (and plenty functional ones as well)

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- ▶ Object oriented programming is one of the very few techniques that actually seems to make it easier for humans to build software
- ▶ But it does not work well for every problem
- ▶ However, since it is so pervasive, this tends to be forgotten

# Summary

- ▶ We have looked at the four aspects that define object-oriented programming



# Summay

- ▶ We have looked at the four aspects that define object-oriented programming
- ▶ **Next time:** Some issues with semantics and typing in OO languages