

# CS345H: Programming Languages

## Lecture 18: Java

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## Overview

- ▶ Java: Our model OO language
- ▶ Arrays
- ▶ Exceptions
- ▶ Interfaces
- ▶ Coercions
- ▶ Threads
- ▶ Dynamic Loading and Initialization

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## Java History

- ▶ Began as **Oak** at SUN
  - ▶ Originally targeted at set-top devices
  - ▶ Initial development from 91-94
- ▶ Retargeted as the Internet Language 94-95
- ▶ **Lesson:** Every new language needs a "killer app"
- ▶ Java beat out TCL, Python

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## The People

- ▶ James Gosling, principal designer. CMU PhD
- ▶ Bill Joy. ABD from Berkeley (UNIX)
- ▶ Guy Steele, MIT PhD and famous language designer

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## Influences

- ▶ Modula 3: Types
- ▶ Eiffel, Objective C, C++: Object orientation, Interfaces
- ▶ LISP: Java's dynamic flavors (lots of them)

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## Java Design

- ▶ From our perspective, what we discussed last time plus
  - ▶ Exceptions
  - ▶ Interfaces
  - ▶ Threads
  - ▶ Dynamic Loading
  - ▶ Other less important ones ...
- ▶ Java is a **big** language
- ▶ Lots of features and feature interactions
- ▶ And language is still growing

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## Arrays

- ▶ Let's look at arrays in Java
- ▶ Assume  $B$  is a subtype of  $A$  ( $B < A$ )
- ▶ What happens in the following?  

```
B[] b = new B[10];  
a[] a = b;  
  
a[0] = new A();  
b[0].aMethodNotDeclaredInA();
```

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## Subtyping in Java

- ▶  $B < A$  if  $B$  inherits from  $A$
- ▶  $C < A$  if  $C < B$  and  $B < A$
- ▶  $B[] < A[]$  if  $B < A$
- ▶ Last rule is unsound!

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## The Right Solution

- ▶ Disallow subtyping through arrays
- ▶  $B < A$  if  $B$  inherits from  $A$
- ▶  $C < A$  if  $C < B$  and  $B < A$
- ▶  $B[] < A[]$  if  $B = A$

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## The Java Solution

- ▶ Java "fixes" this problem by checking each array assignment at **run-time** for type correctness
- ▶ This means that we can now get type errors at run-time!
- ▶ Also, huge overhead on array computations
- ▶ **Note:** Primitive types unaffected since they are not classes in Java

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## A Common Problem

- ▶ Deep in a section of code, you encounter an unexpected error
  - ▶ Out of memory
  - ▶ A list that is supposed to be sorted is not
  - ▶ ...
- ▶ What do you do?

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## Exceptions

- ▶ Add a new type (class) of **exceptions**
- ▶ Add new syntactic forms:  

```
try { something } catch(x) { cleanup }  
throw exception
```

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## Exceptions Example

- ▶ Add a new type (class) of **exceptions**
- ▶ Add new syntactic forms:

```
class Foo {
    public static void main(String[] args) {
        try { X(); } catch (Exception e) {
            System.out.println("Error!") } }

    public void X() throws MyException {
        throw new MyException();
    }
}
```

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## Operational Semantics with Exceptions

- ▶  $T(o)$  an exception has been thrown
- ▶  $o$  = an ordinary object
- ▶ Here is a (pseudo) rule for try:

$$\frac{E \vdash e_1 : o}{E \vdash \text{try}\{e_1\}\text{catch}(x)\{e_2\} : o}$$

$$\frac{\begin{array}{l} E \vdash e_1 : T(o_1) \\ E[x \leftarrow o_1] \vdash e_2 : o_2 \end{array}}{E \vdash \text{try}\{e_1\}\text{catch}(x)\{e_2\} : o_2}$$

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## Semantics (Cont.)

$$\frac{E \vdash e : o}{E \vdash \text{throw } e : T(o)}$$

$$\frac{E \vdash e_1 : T(o)}{E \vdash e_1 + e_2 : T(o)}$$

- ▶ All forms except catch propagate thrown exceptions

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## Simple Implementation

- ▶ When we encounter a **try**, mark current location in stack
- ▶ When we throw an exception, unwind the stack to the first **try** and execute corresponding **catch**
- ▶ More complex techniques reduce the cost of try and throw

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## Type Checking and Exceptions

- ▶ Methods must declare types of exceptions they may raise

```
public void X() throws MyException
```

  - ▶ Checked at compile time
  - ▶ Some exception need not be declared, such as dereferencing null pointers
- ▶ Other expected rules for exception constructs, such as throw must be applied to something of type Exception

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## Interfaces

- ▶ Specify relationships between classes without inheritance

```
interface PointInterface {
    void move(int dx, int dy);}

class Point implements PointInterface {
    void move(int dx, int dy) { ... }
}
```

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## Interfaces

- ▶ “ Java programs can use interfaces to make it unnecessary for related classes to share a common abstract super class or to add methods to it.”
- ▶ In other words, interfaces play the same role as multiple inheritance in C++, because classes can implement multiple interfaces

## Why is this Useful?

- ▶ A graduate student may be both a University employee and a student  
`class GraduateStudent implements Employee, Student {...}`
- ▶ No good way to incorporate Employee, Student methods for grad students with single inheritance

## Coercions

- ▶ Java allows primitive types to be **coerced** in certain contexts
- ▶ In `1 + 2.0`, the int 1 is widened to a float 1.0
- ▶ A coercion is really just a primitive function the compiler inserts for you
- ▶ Most languages have extensive coercions between basic numeric types

## Coercions & Casts

- ▶ Java distinguished two kinds of coercions & casts:
  1. **Widening**: Always succeeds (int -> float)
  2. **Narrowing**: May fail or lose information (float -> int, downcasts)
- ▶ Narrowing casts must be explicit
- ▶ Widening casts/coercion can be implicit

## Coercions in PL/I

- ▶ Let A,B,C be strings of 3 characters:  
`B = '123'`  
`C = '456'`  
`A = B+C`
- ▶ What is A?

## Threads

- ▶ Java has concurrency built in through **threads**
- ▶ Threads have class Thread and start and stop methods
- ▶ Synchronization obtains a lock on the object:  
`synchronized { e }`
- ▶ In synchronized methods, this is locked

## Example

```
class Simple {
  int a = 1, b = 2;
  void to() { a = 3; b = 4; }
  void fro() {println("a= " + a + ", b=" + b); } }
```

- ▶ Two threads call to() and fro(). What is printed?

## Example (Cont.)

```
class Simple {
  int a = 1, b = 2;
  void synchronized to() { a = 3; b = 4; }
  void fro() {println("a= " + a + ", b=" + b); } }
```

- ▶ Two threads call to() and fro(). What is printed?

## Example (Cont.)

```
class Simple {
  int a = 1, b = 2;
  void synchronized to() { a = 3; b = 4; }
  void synchronized fro() {println("a= " + a + ", b=" +
b); } }
```

- ▶ Two threads call to() and fro(). What is printed?

## Semantics

- ▶ Even without synchronization, a variable should only hold values written by some thread
  - ▶ Writes of values are atomic
  - ▶ Violated for doubles, though
- ▶ Java's concurrency semantics are very difficult in details, but at least Java attempts to specify them!

## Dynamic Loading

- ▶ Java allows classes to be loaded at run time
  - ▶ Type checking of source code takes place at compile time
  - ▶ Bytecode verification takes place at run-time
- ▶ This introduces many additional complications, but can be extremely flexible

## Features and Feature Interactions

- ▶ In any system with  $N$  features, there are potentially  $N^2$  feature interactions
- ▶ Big featureful systems are hard to understand, including programming languages

## Summary

- ▶ Java is pretty well done. By production language standards, very well done.
- ▶ Java brought many important ideas to the mainstream, such as strong static typing and garbage collection
- ▶ But Java also has many features that are hard to understand and lots of features