

## Introduction to Mathematical Logic, Handout 9

### Logical Validity and Entailment

In predicate logic, we say about a sentence that it is *logically valid* if it is satisfied by all interpretations.

**Problem 9.1** Determine whether the sentences

$$\begin{aligned}\exists x(P(x) \wedge Q(x)) &\rightarrow (\exists xP(x) \wedge \exists xQ(x)), \\ (\exists xP(x) \wedge \exists xQ(x)) &\rightarrow \exists x(P(x) \wedge Q(x))\end{aligned}$$

are logically valid.

The *universal closure* of a formula  $F$  is the sentence  $\forall v_1 \dots v_n F$ , where  $v_1, \dots, v_n$  are all free variables of  $F$ . About a formula with free variables we say that it is *logically valid* if its universal closure is logically valid.

**Problem 9.2** For each of the formulas

$$\begin{aligned}P(x) &\rightarrow \exists xP(x), \\ P(x) &\rightarrow \forall xP(x)\end{aligned}$$

determine whether it is logically valid.

A formula  $F$  is *equivalent* to a formula  $G$  if the formula  $F \leftrightarrow G$  is logically valid.

**Problem 9.3** Determine whether the formula

$$\forall x \exists y P(x, y)$$

is equivalent to

$$\exists y \forall x P(x, y).$$

We say that a set  $\Gamma$  of sentences *entails* a sentence  $F$ , or that  $F$  is a *logical consequence* of  $\Gamma$ , if every interpretation that satisfies all sentences in  $\Gamma$  satisfies  $F$ .

**Problem 9.4** Determine whether

$$\exists x P(x, x)$$

is a logical consequence of the sentences

$$\forall x \exists y P(x, y), \forall x y z ((P(x, y) \wedge P(y, z)) \rightarrow P(x, z)).$$