

Workflow Patterns in Orc

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Overview of Orc

- **Orchestration language**
 - Invoke services
 - Manage time-outs, priorities, and failures
- **Structured concurrent programming**
 - Implicit “result” channel, also explicit channels
 - Easy to create and terminate processes
- **Simple calculus with formal semantics**
 - Labeled transition system & traces
 - Needs to be extended for time, etc.
- **Prototype implementation available**

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Pipe: $f > x > g$

CNN $> n >$ Email(user, n)

- call CNN, bind result (if any) to n
- then call **Email** to send news to user

- **CNN and Email are sites**
 - Perform basic computations
 - Return at most one value, or never return

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Parallel: $f | g$

(CNN | BBC) $> n >$ Email(user, n)

- CNN | BBC may produce 0, 1, or 2 values

- **Pipe also acts as “for each”**
 - Email is sent for *each* value
 - Instances of **g** are executed in parallel
- **Use $f \gg g$ if no variable is needed**
 - if f returns 1 value, acts as **f; g**

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Where: $f \text{ where } x : \in g$

Email(user, n)
where n : \in (CNN | BBC)

- Binds n to *first* value of CNN | BBC
- then *terminates* CNN | BBC

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Some Basic Sites

| | |
|-------------------|--|
| let(x, ..., z) | Returns argument values as a tuple |
| if(b) | Returns a signal if b is true; it does not respond if b is false |
| or(a, b) a + b | Sites, for basic operations |
| Rtimer(t) | Relative Timer: returns a signal after exactly t time units |
| 0 | never returns |

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Definitions: $E(x_1, \dots, x_n) \Delta f$

- **Definitions**
First(g) Δ let(x) **where** x : \in g
- **Using a definition**
First(CNN | BBC) $> n >$ Email(user, n)

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Subtleties: $f \text{ where } x : \in g$

- f and g are executed in *parallel*
- **Site calls are strict**
 - not called until arguments are defined
- **Example**
(M | N(x)) **where** x : \in P
 - M and P called immediately
 - N is called after P returns

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A larger example: Priority

- If M responds in 10 time units, take its response, otherwise take first response

Delay(N) \triangleq

(Rtimer(10) >> let(u)) **where** u : \in N

Priority(M, N) \triangleq First(M | Delay(N))

- Concise notation for
 - communication, blocking, and termination

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Another example: Parallel Or

- Return true if either M or N return true

POr(M, N) \triangleq let(z)

where z : \in if(x) | if(y) | or(x, y)

where x : \in M

where y : \in N

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Orc Summary

| | |
|---|--------------------|
| e, f, g ::= c | constant |
| x | variable |
| x(e ₁ , ..., e _n) | call |
| f >x> g | pipe (also f >> g) |
| f g | parallel |
| f where x : \in g | asymmetric |
| x(x ₁ , ..., x _n) \triangleq f | definition |

- See web site for semantics, implementation
www.cs.utexas.edu/~wcook/projects/orc

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Workflow Patterns

- Workflow products
 - use pictures to define workflows
 - resemble Petri nets, Statecharts, concurrent flowcharts
- No formal model of workflow
- Alternative: identify common *Patterns*
 - 20 patterns have been proposed
- We show that Orc can implement the workflow patterns

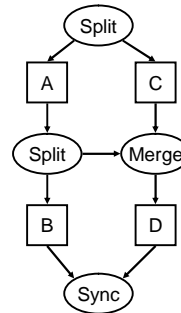
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Simple Workflow Patterns

1. Sequence
Seq(f, g) \triangleq f >> g
2. Parallel
Par(f*) \triangleq f₁ | ... | f_n
3. Synchronization (fork-join):
Sync(f, g) \triangleq let(x, y) **where** x : \in f
where y : \in g

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Unstructured Workflows



Condition >M>
 Sync(A >> M.set >> B,
 C >> M.wait >> D)

Condition creates a local object with two methods, set and wait, which blocks until set is called.

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Choices

4. Exclusive Choice (arbitration):
XOR(b, f, g) \triangleq (if(b) >> f) | (if(¬b) >> g)
6. Multi-choice:
Multi(b*, f*) \triangleq Par(XOR(b_i, f_i, 0))
7. Synchronizing Merge:
SyncMerge(b*, f*) \triangleq
 Sync(XOR(b_i, f_i, Signal))

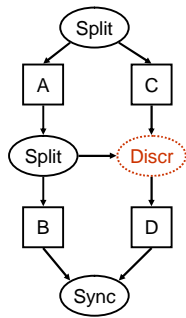
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9. Discriminator

- Produce result when first f_i completes
 - Continue other computations
- Discr(f*) \triangleq Buffer >S>
 S.get | (Par(f*) >x> S.put(x))
- Buffer creates a local object with two methods, put and get, which blocks until put is called.

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Modular Composition

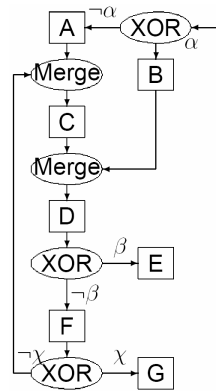


Condition $\langle M \rangle$
 $\text{Sync}(A \gg M.\text{set} \gg B,$
 $\text{Discr}(C, M.\text{wait}) \gg D)$

D runs after C or A completes

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10. Arbitrary Cycles



$P \triangle \text{XOR}(\alpha, PB, PA)$

$PA \triangle A \gg PC$

$PB \triangle B \gg PD$

$PC \triangle C \gg PD$

$PD \triangle D \gg \text{XOR}(\beta, E, PF)$

$PF \triangle F \gg \text{XOR}(\gamma, G, PC)$

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12-15: Instantiation

- **Creating multiple instances of workflows**

12. No synchronization
13. Design-time knowledge
14. Run-time knowledge
15. No knowledge

- **These can be hard with Petri nets, but easy with a process calculus**

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16 Deferred Choice

- **Let the environment make a choice by signaling an event**

$\text{DefChoiceTerm}(e^*, f^*) \triangle$

$\text{Which}(e^*) \triangleright k \triangleright \text{Select}(k, f^*)$

$\text{Which}(e^*) \triangle$

$\text{First}(e_1 \gg \text{let}(1) \mid \dots \mid e_n \gg \text{let}(n))$

$\text{Select}(k, f^*) \triangle$

$\text{if}(k = 1) \gg f_1 \mid \dots \mid \text{if}(k = n) \gg f_n$

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17: Interleaved Parallel Routing

- Order is decided at run-time, and no two activities are executed at the same moment

$\text{Interleave}(f^*) \triangle \text{Lock} \triangleright M \triangleright$

$(\text{wait}(M, f_1) \mid \dots \mid \text{wait}(M, f_n))$

$\text{wait}(M, f) \triangle$

$M.\text{acquire} \gg f \triangleright x \triangleright M.\text{release} \gg \text{let}(x)$

–Lock creates a local object with two methods, acquire and release. Only one process may acquire the lock at a time

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19/20: Cancel Activity/Case

- **Interrupt f when event e occurs**

$\text{Interrupt}(f, e) \triangle \text{First}(f \mid e)$

- **Interrupt f after t time units, if it is not complete**

$\text{Timeout}(f, t) \triangle \text{Interrupt}(f, \text{Rtimer}(t))$

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Evaluation

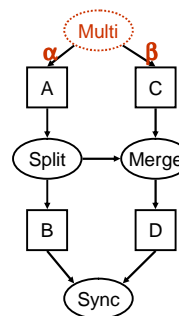
- **Van der Aalst's concern with encoding is legitimate**

- Encodings are fine for theory
- But don't help programmers
- Require manual, global program rewrites

- **Orc uses definitions, not encodings**

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Orc patterns are not fully compositional



Condition $\langle M \rangle$

$\text{SyncMerge}(\langle \alpha, A \gg M.\text{set} \gg B \rangle,$

$\langle \beta, C \gg M.\text{wait} \gg D \rangle)$

–Does not work, because **M.set** is not called if α is false

–Problem: interaction between non-structured control and synchronizing merge

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Related Work

- **van der Aalst**
 - Defined initial set of 20 patterns
 - not claimed to be complete
 - Yet Another Workflow Language (YAWL)
 - Add constructs to Petri nets until they can model the workflow patterns without “encoding”
- **Business Process Markup Language (BPML)**
 - Pattern solutions similar to Orc
 - More verbose, features missing
- **Workflow in π -Calculus (Puhmann, Weske)**
 - Uses explicit channels to signal start/end of workflows
 - Incomplete model: termination requires encoding

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Conclusions

- **Orc can implement all the workflow patterns**
- **Used definitions, avoids encoding**
 - Still work to do: not completely compositional
- **Need to formalize more patterns**
 - Parallel Or
 - Priority
 - Barrier synchronization

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